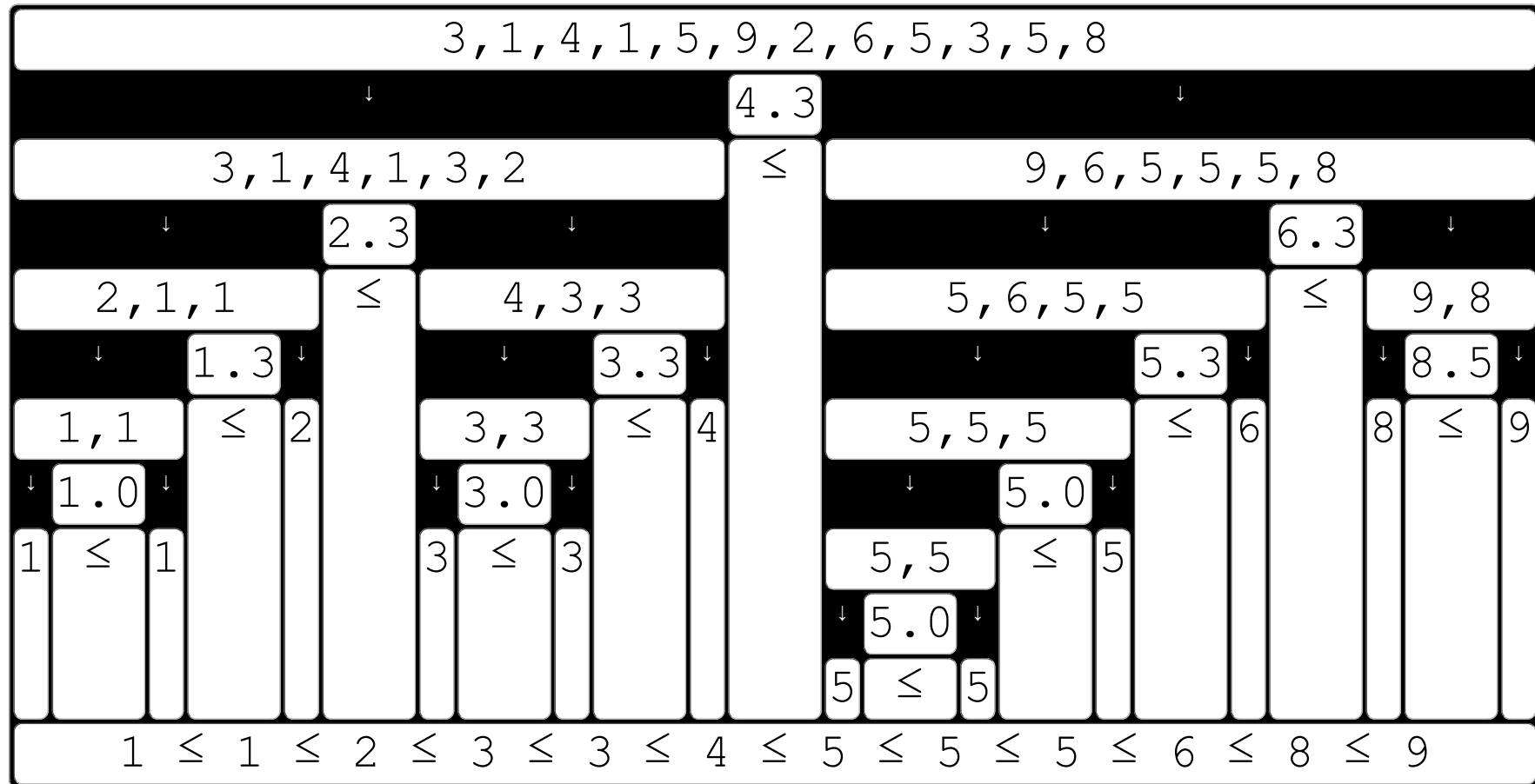


Quicksort

Tutorial Highlight

Julio M. Otuyama

otuyama@alumni.usp.br



previous

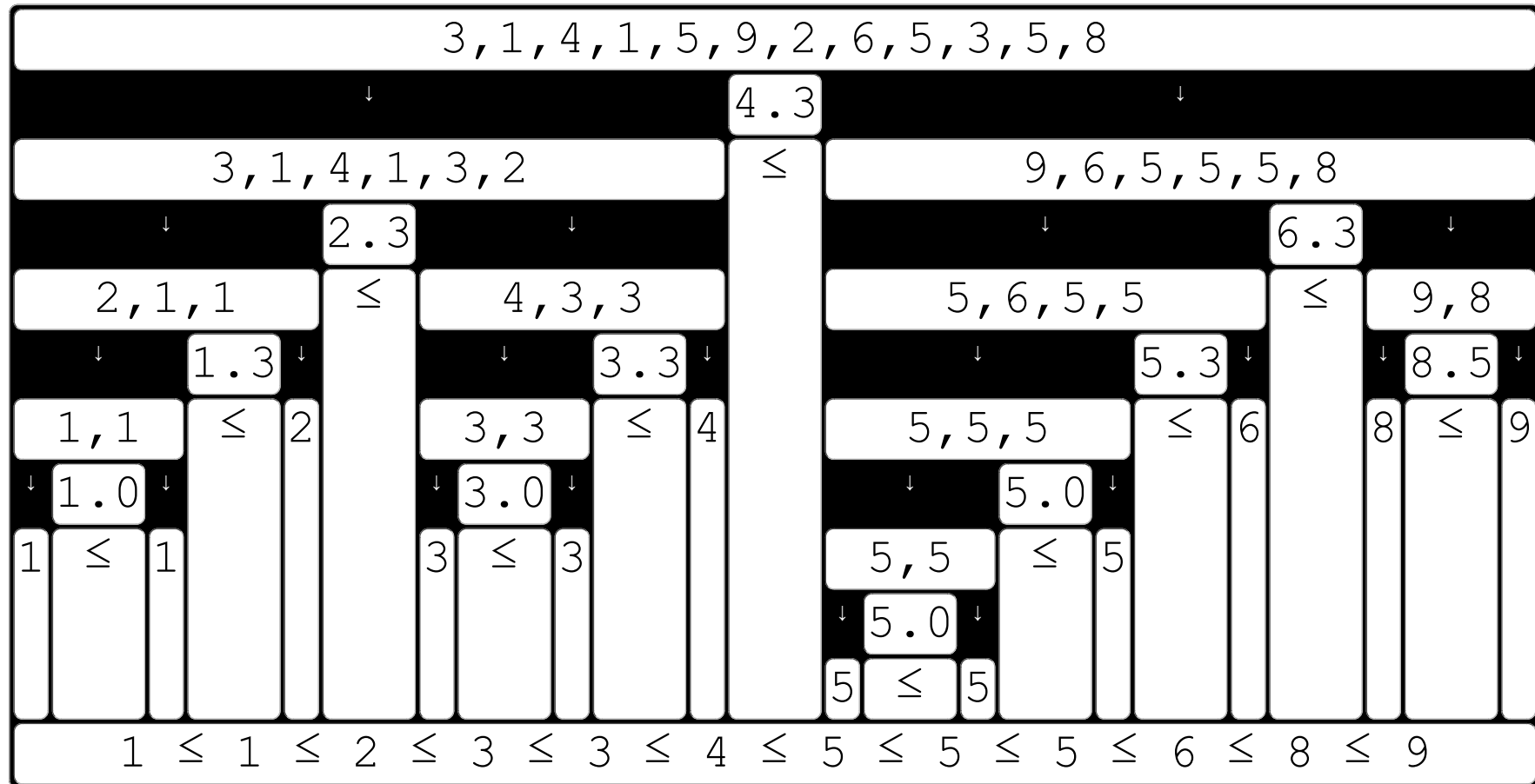
first

Welcome to Quicksort Interactive Tutorial.

Input your own values:

Sort

next

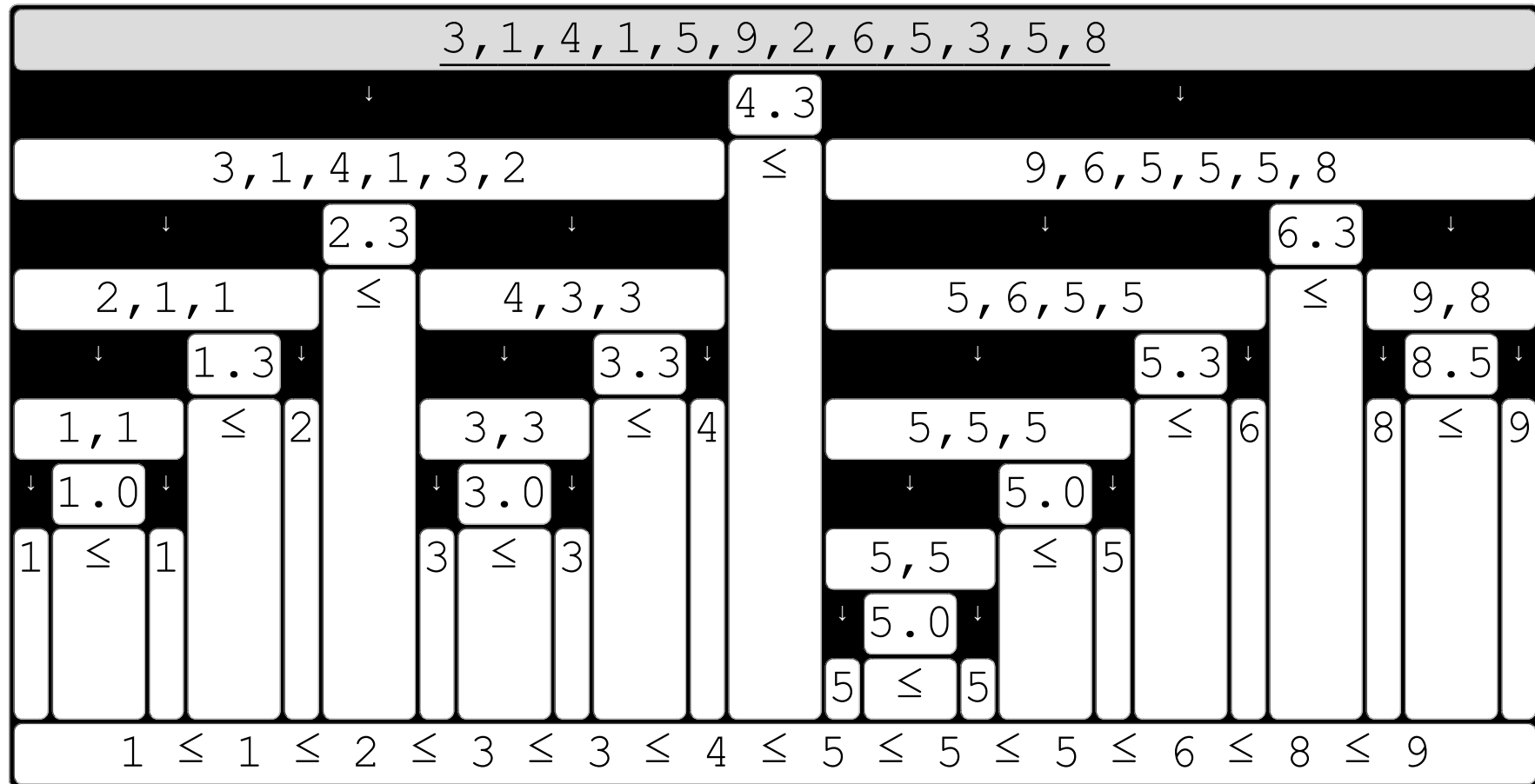


previous

first

(1) Introduction.

next

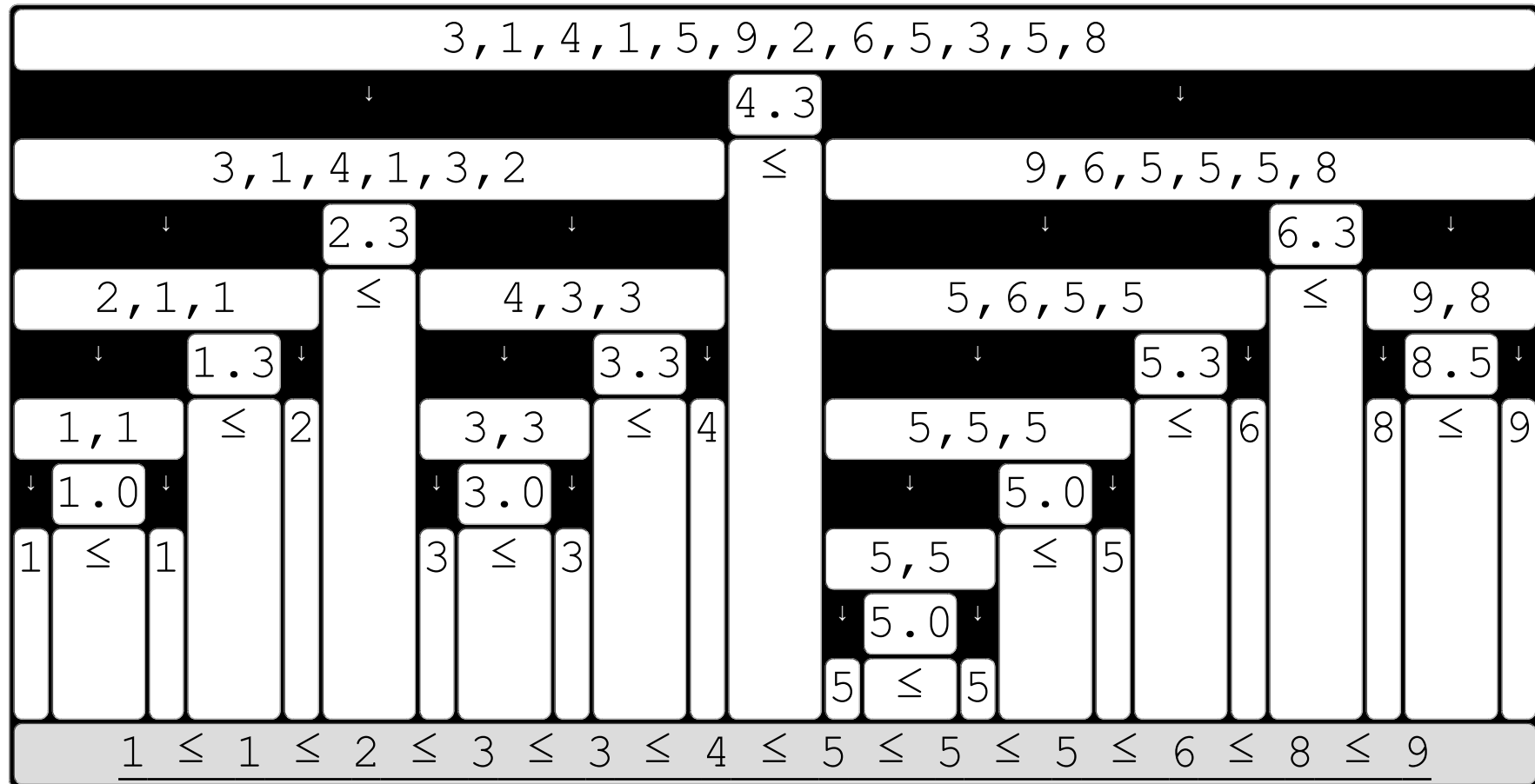


previous

At the top of the diagram, an [initial list](#) contains random values (not sorted).

first

next

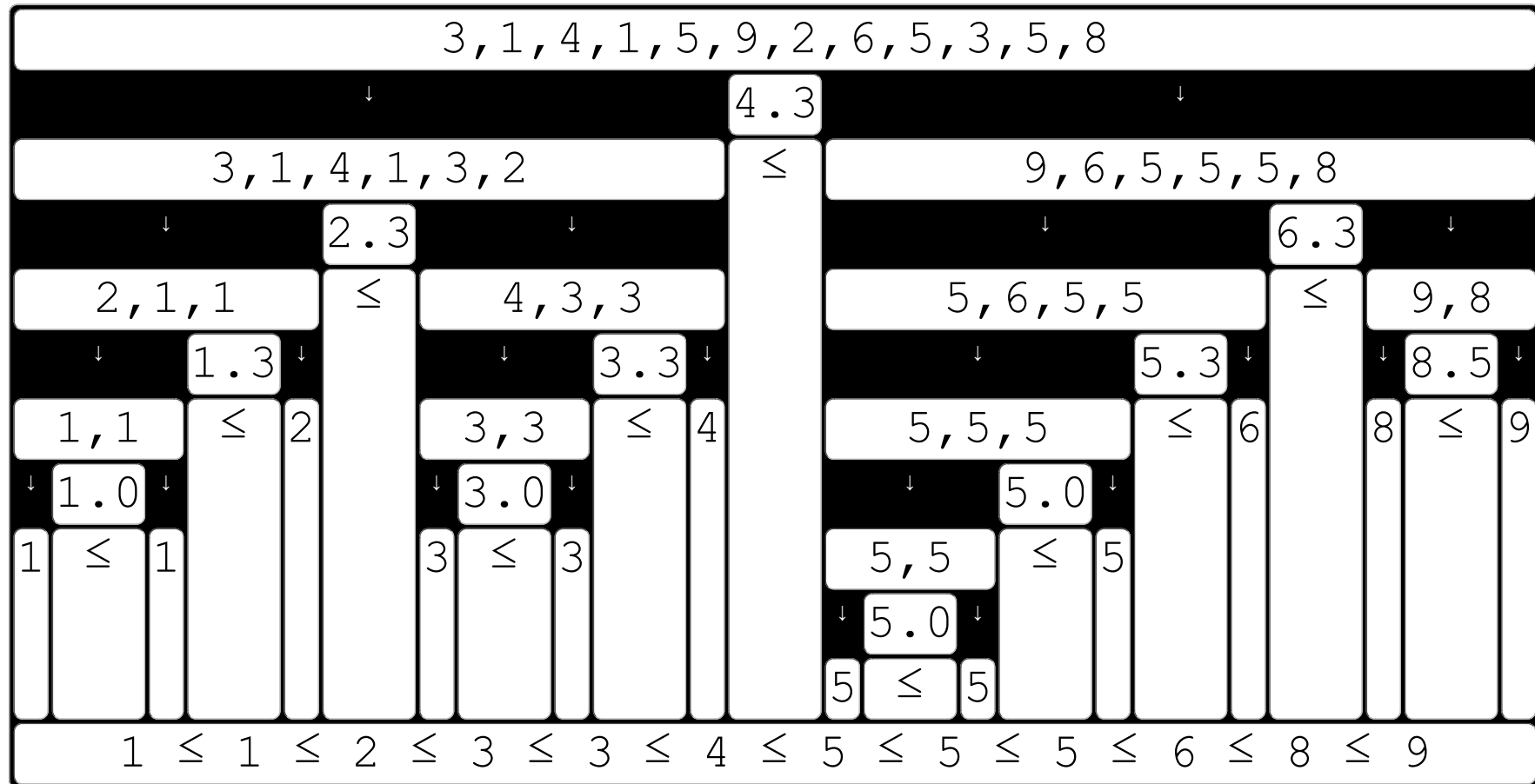


previous

first

At the bottom of the diagram, a [sorted list](#) is the result of the algorithm.

next

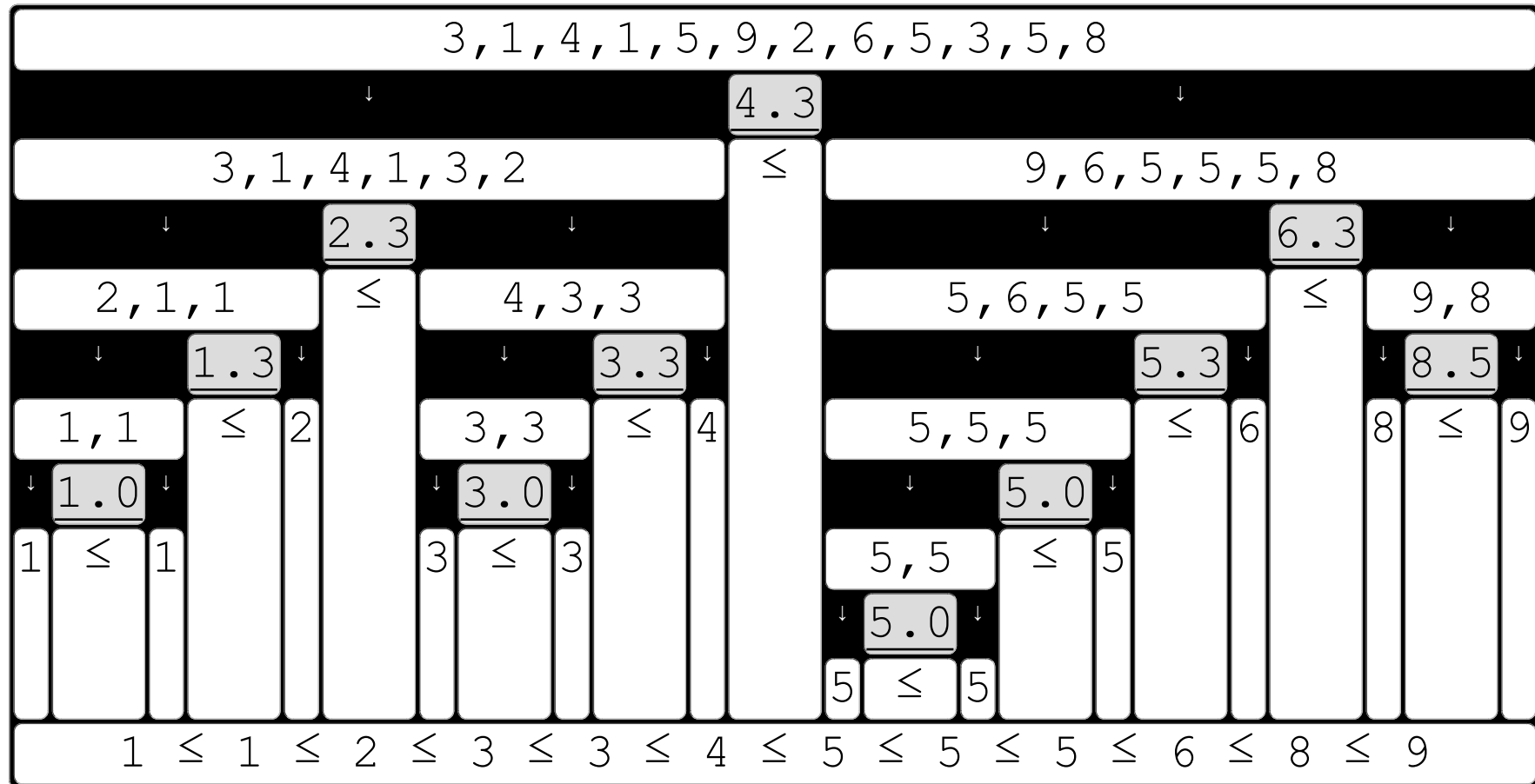


previous

first

Quicksort uses the concept of [pivots](#) to split [each list](#) near the middle.

next

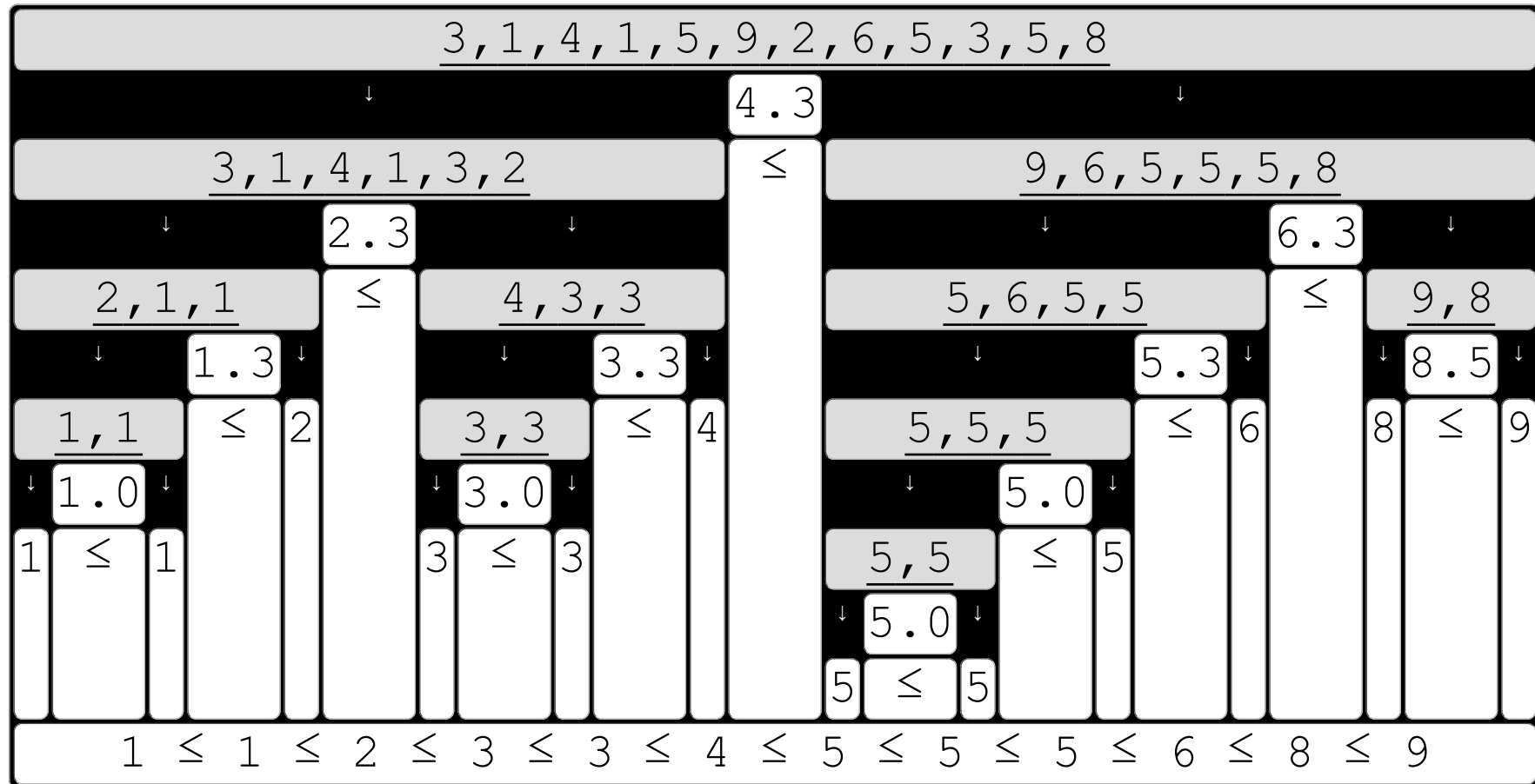


previous

Quicksort uses the concept of [pivots](#) to split each list near the middle.

first

next

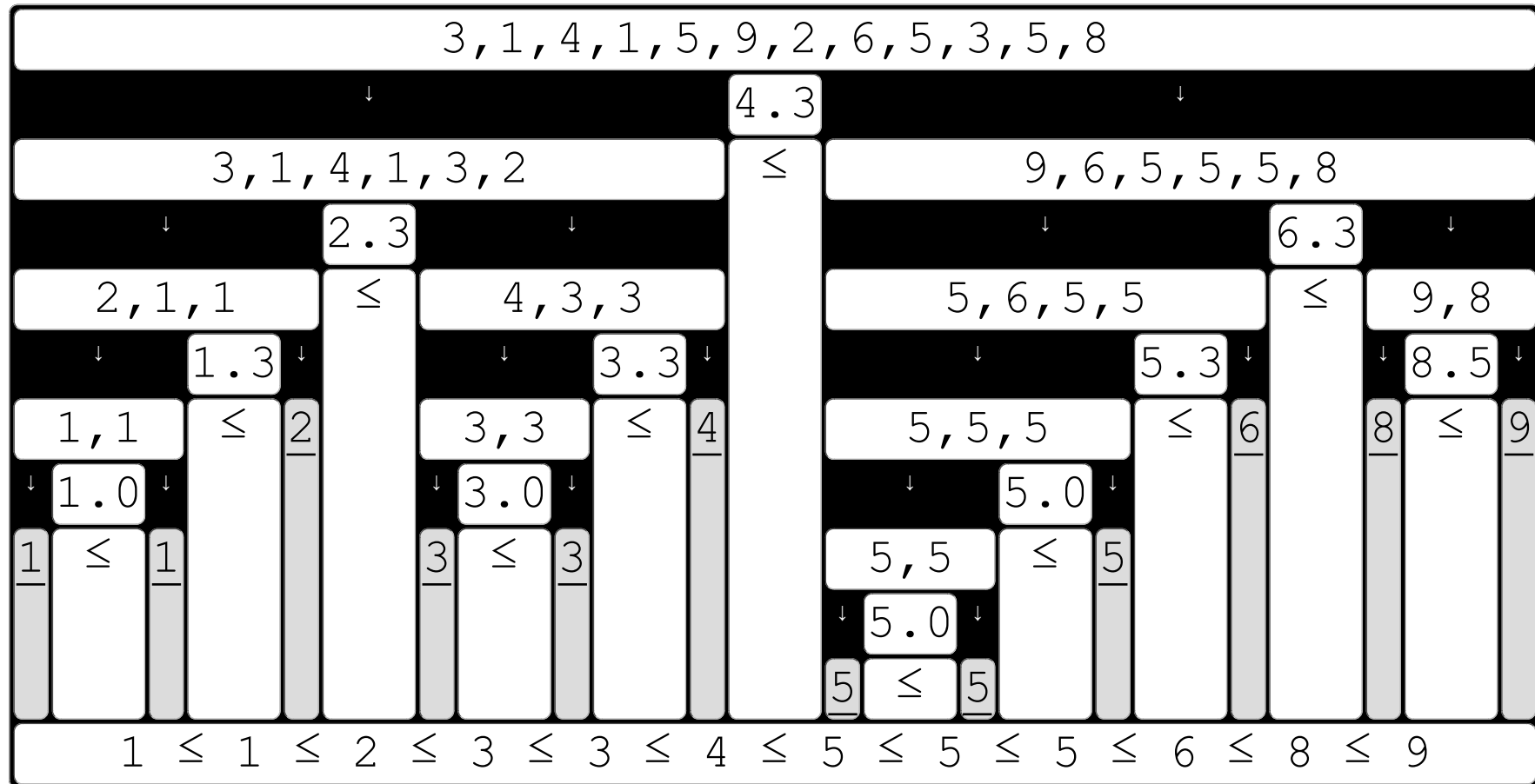


previous

first

Quicksort uses the concept of pivots to split [each list](#) near the middle.

next

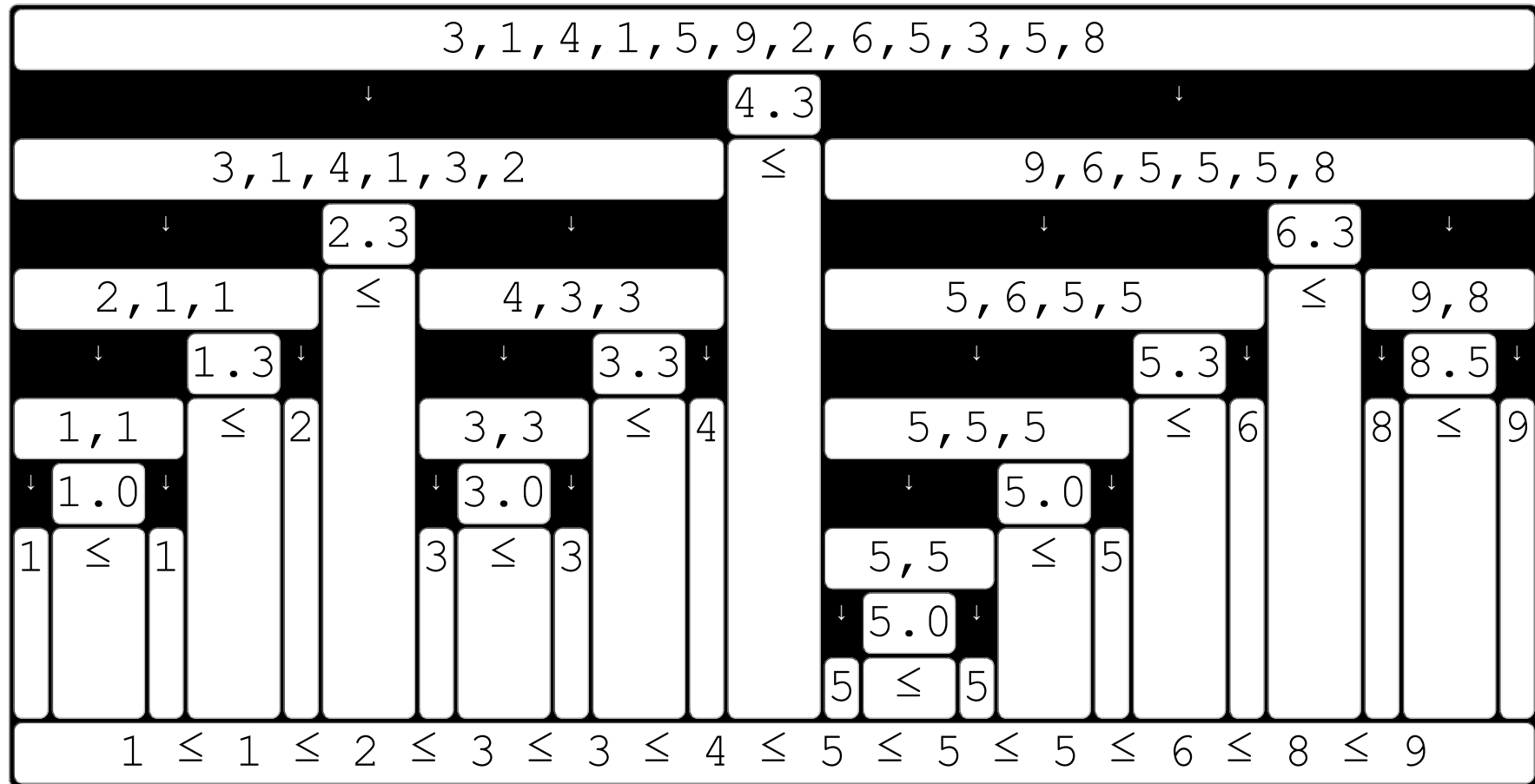


previous

first

This procedure is repeated until each list has only one value (as described in the next steps).

next

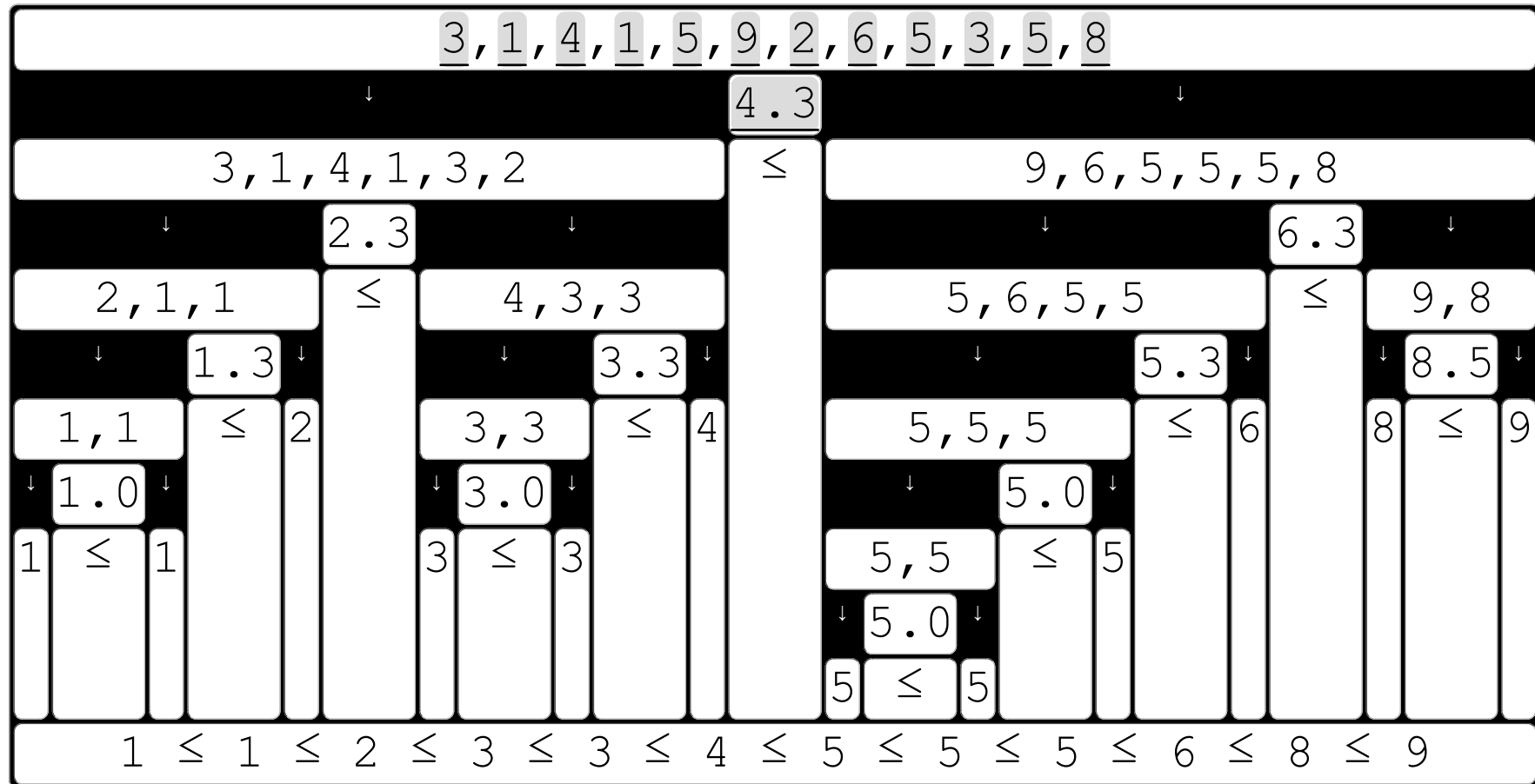


previous

first

(2) Details of Quicksort.

next



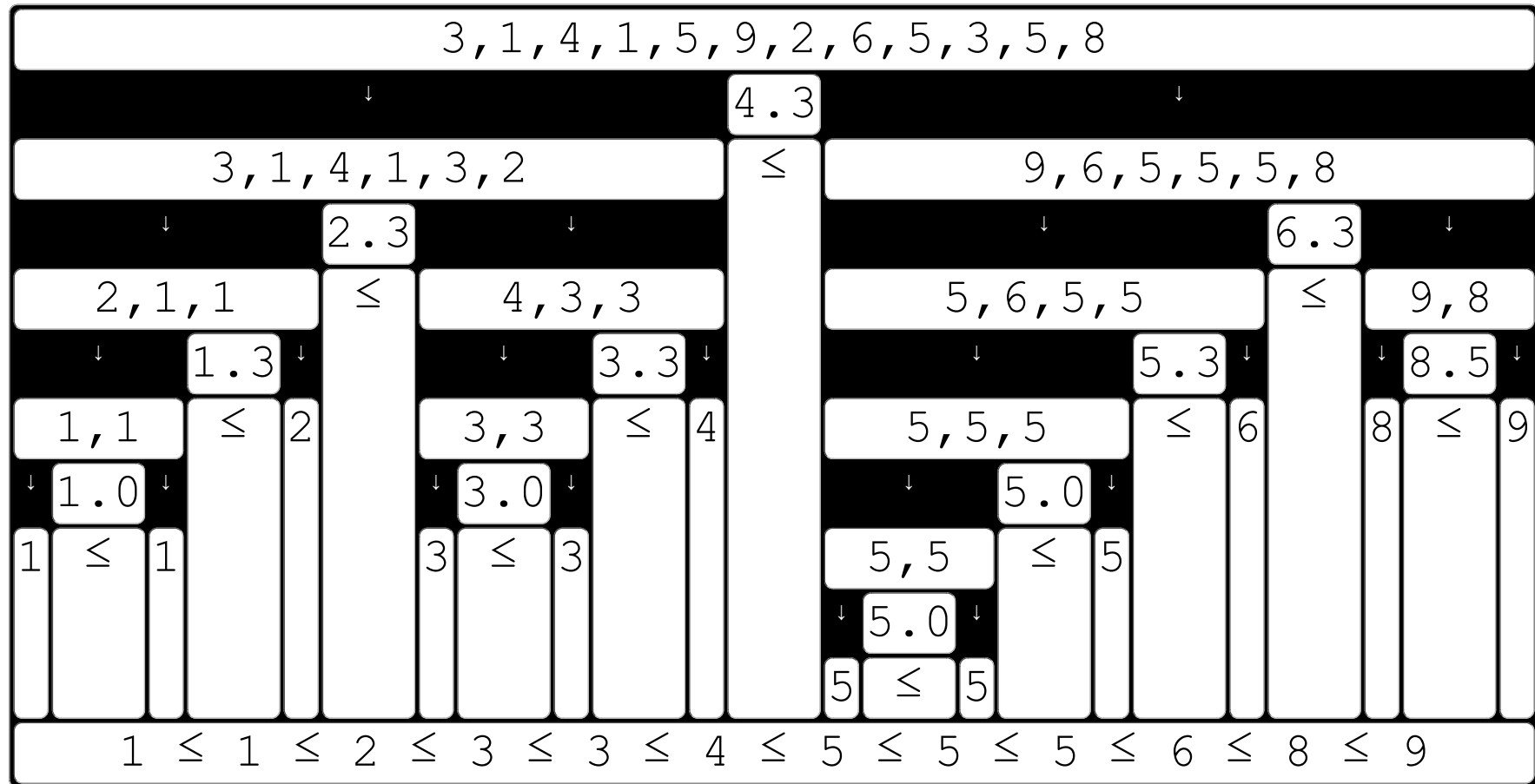
previous

For didactic purposes, the pivot is defined as average (of values of a list). For example:

first

[the first pivot is the average of the initial list.](#)

next

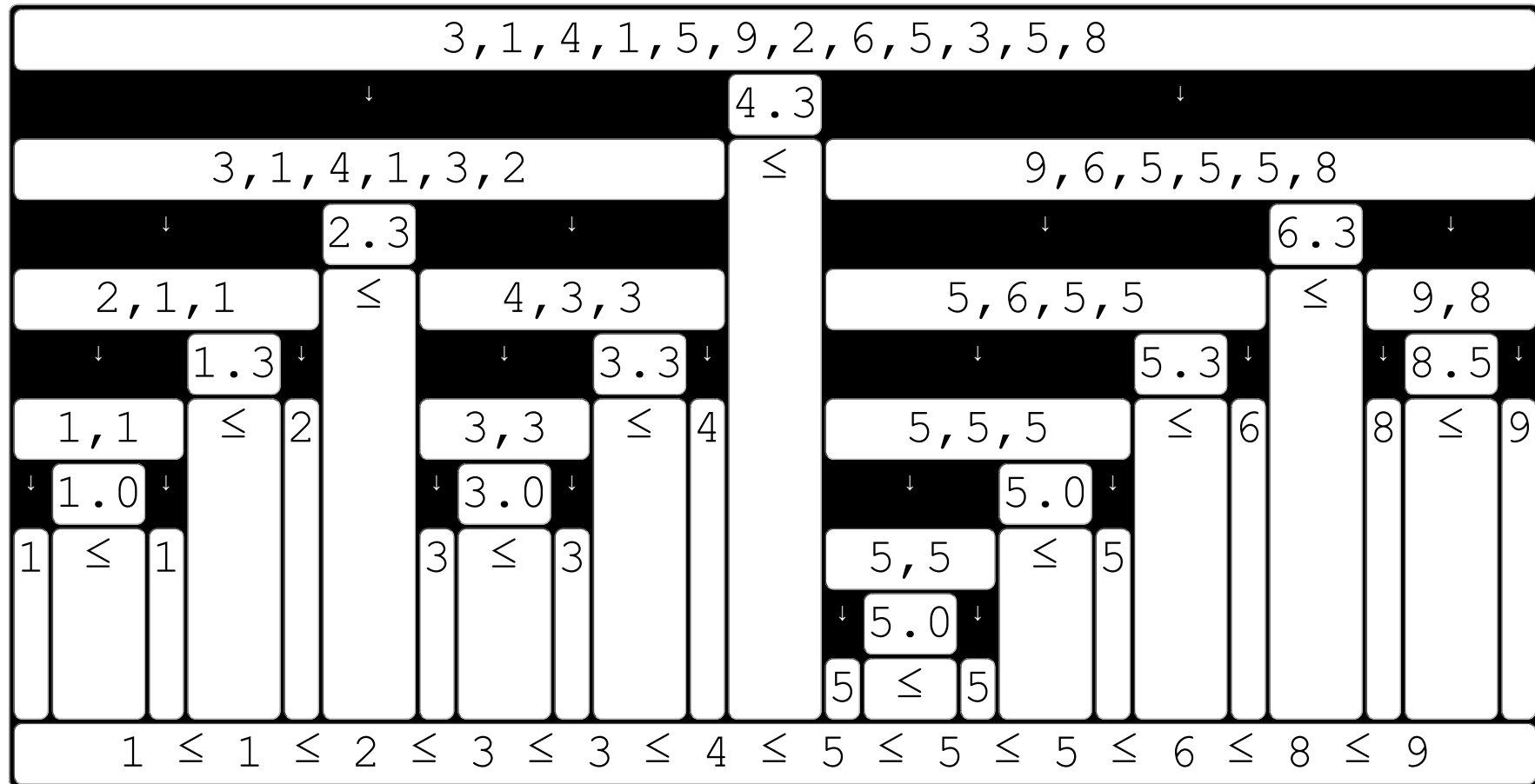


previous

first

Usually, the average is not able to split the list exactly in the middle, but it gets something near that.

next

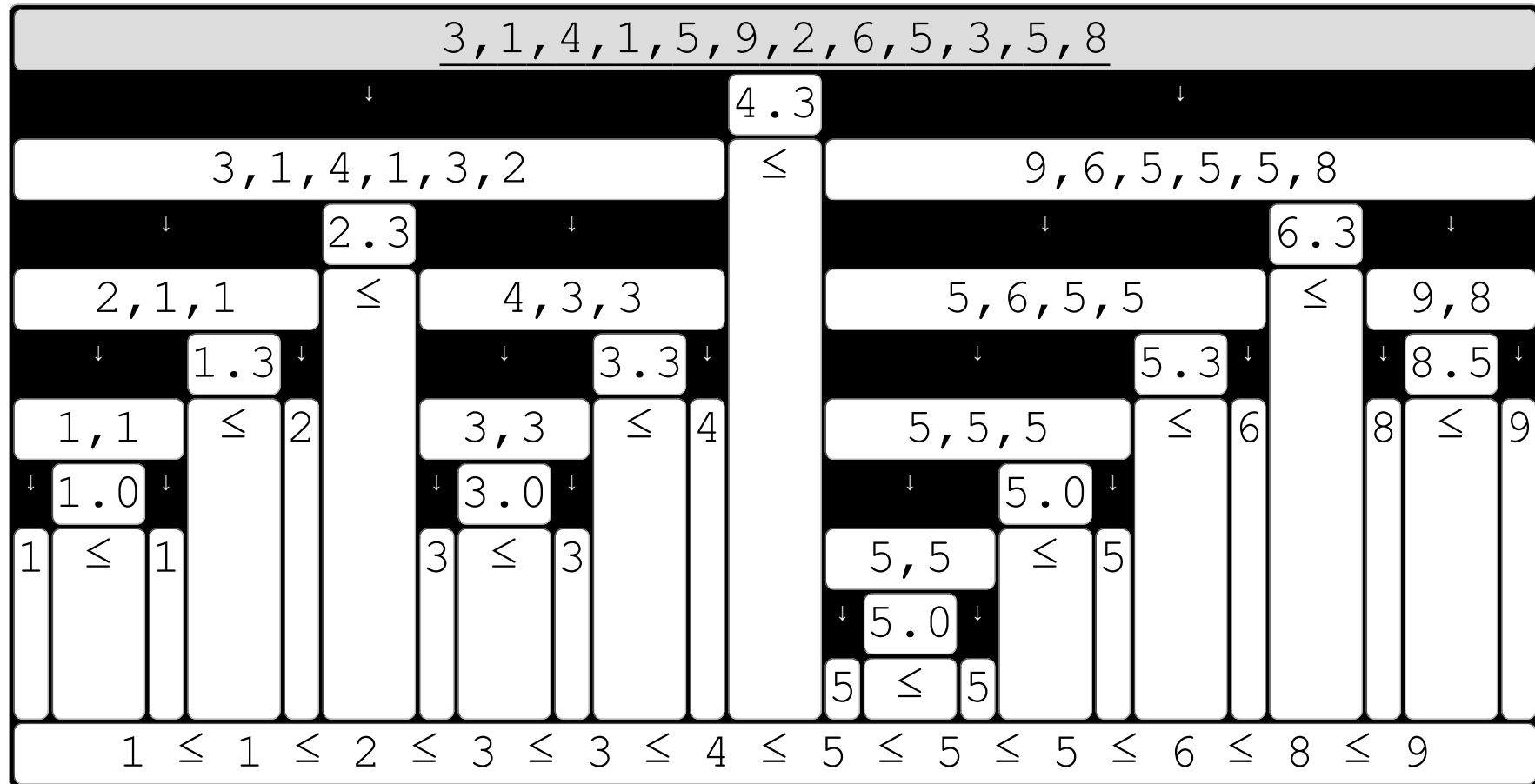


previous

At the first level, the [initial list](#) is splitted (around the [pivot](#)) which results in [two lists](#).

first

next

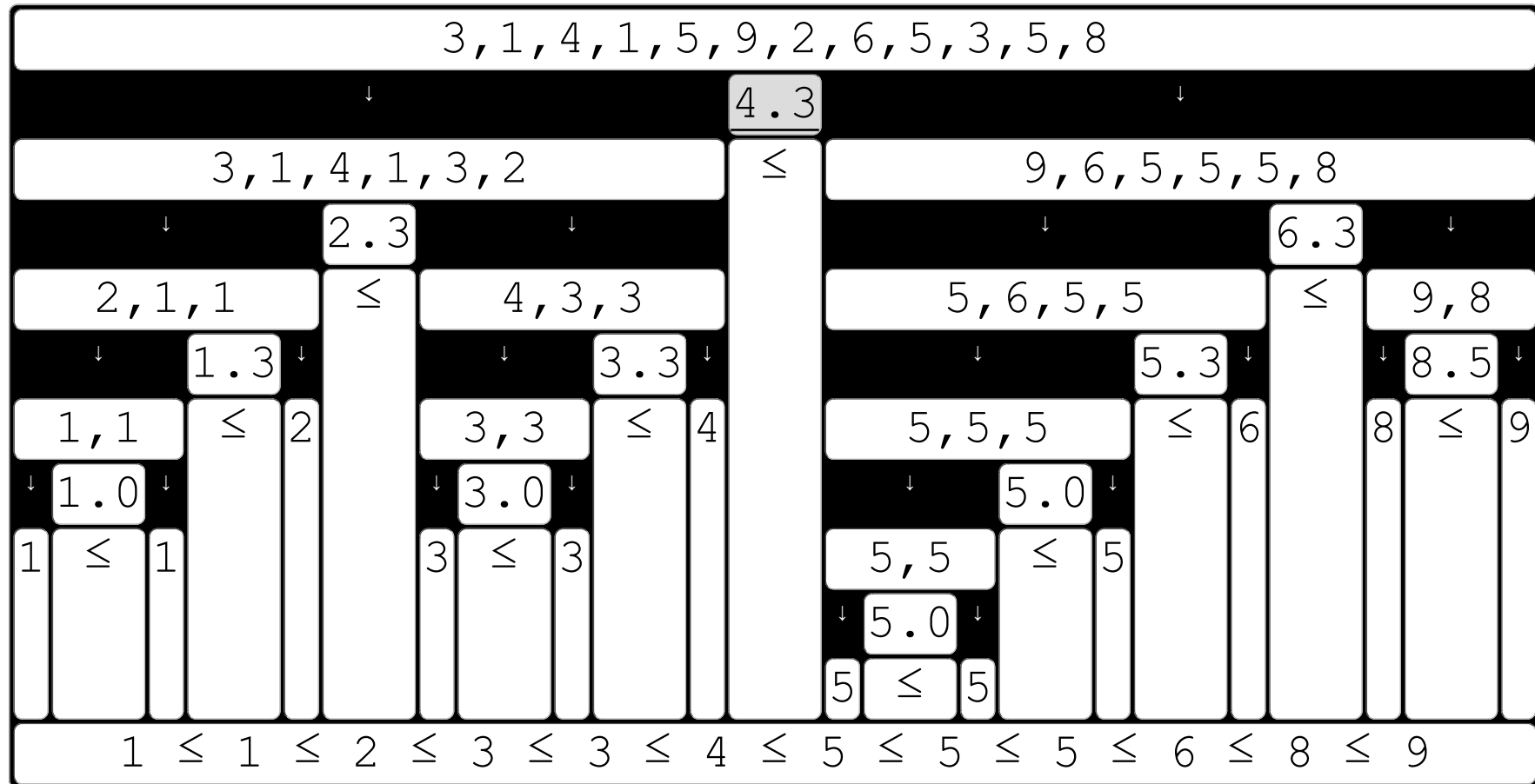


previous

At the first level, the [initial list](#) is splitted (around the pivot) which results in two lists.

first

next

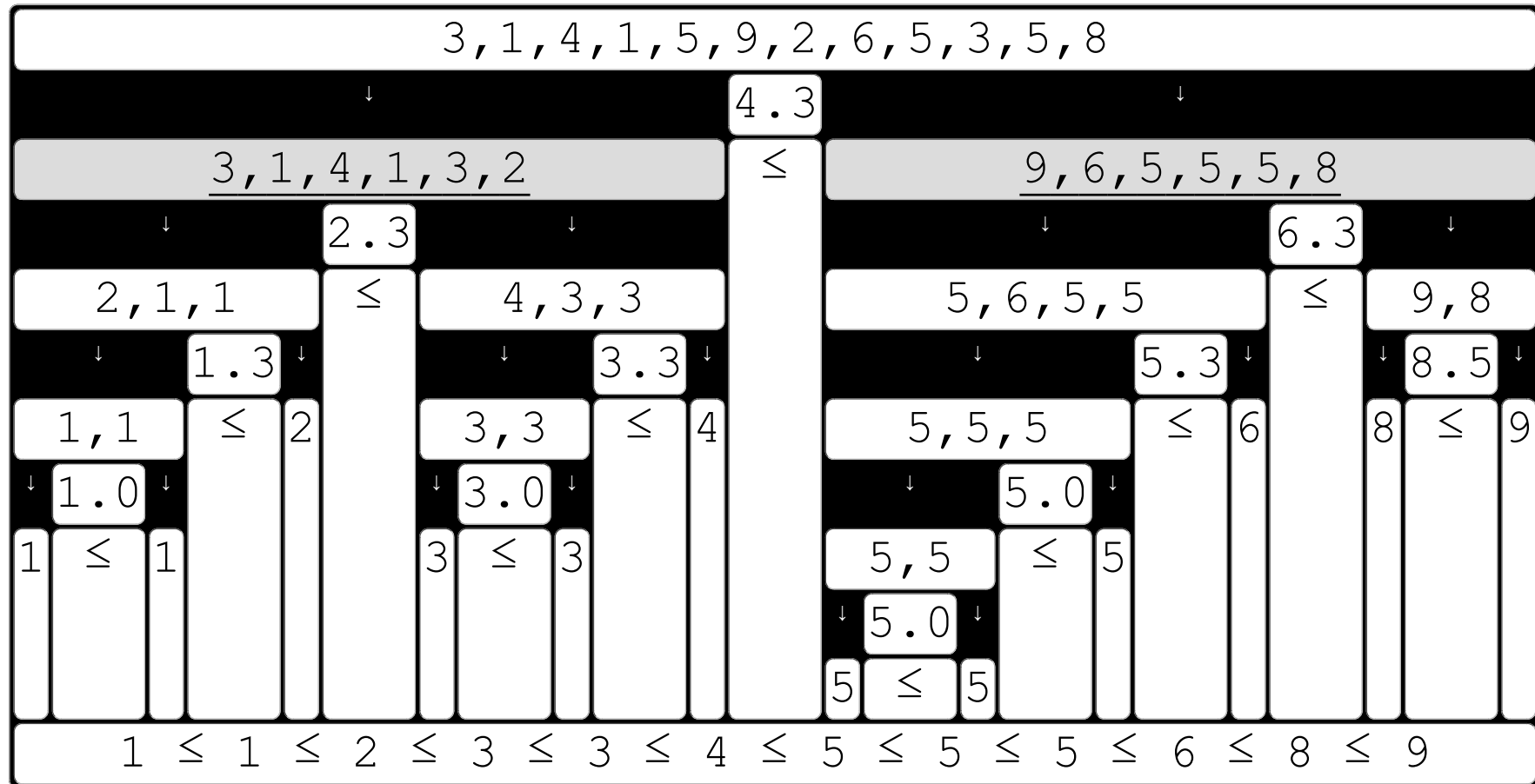


previous

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first

next

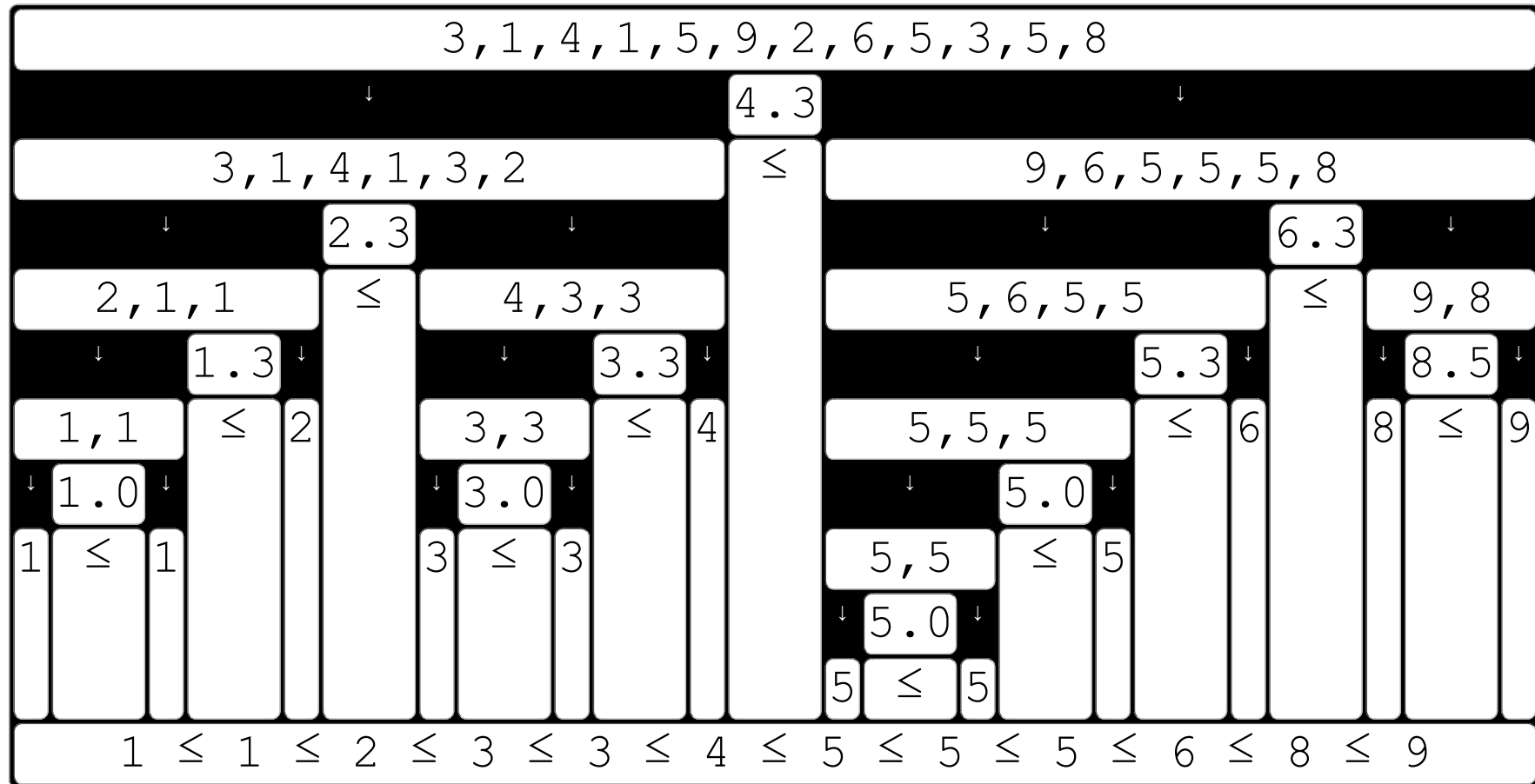


previous

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first

next

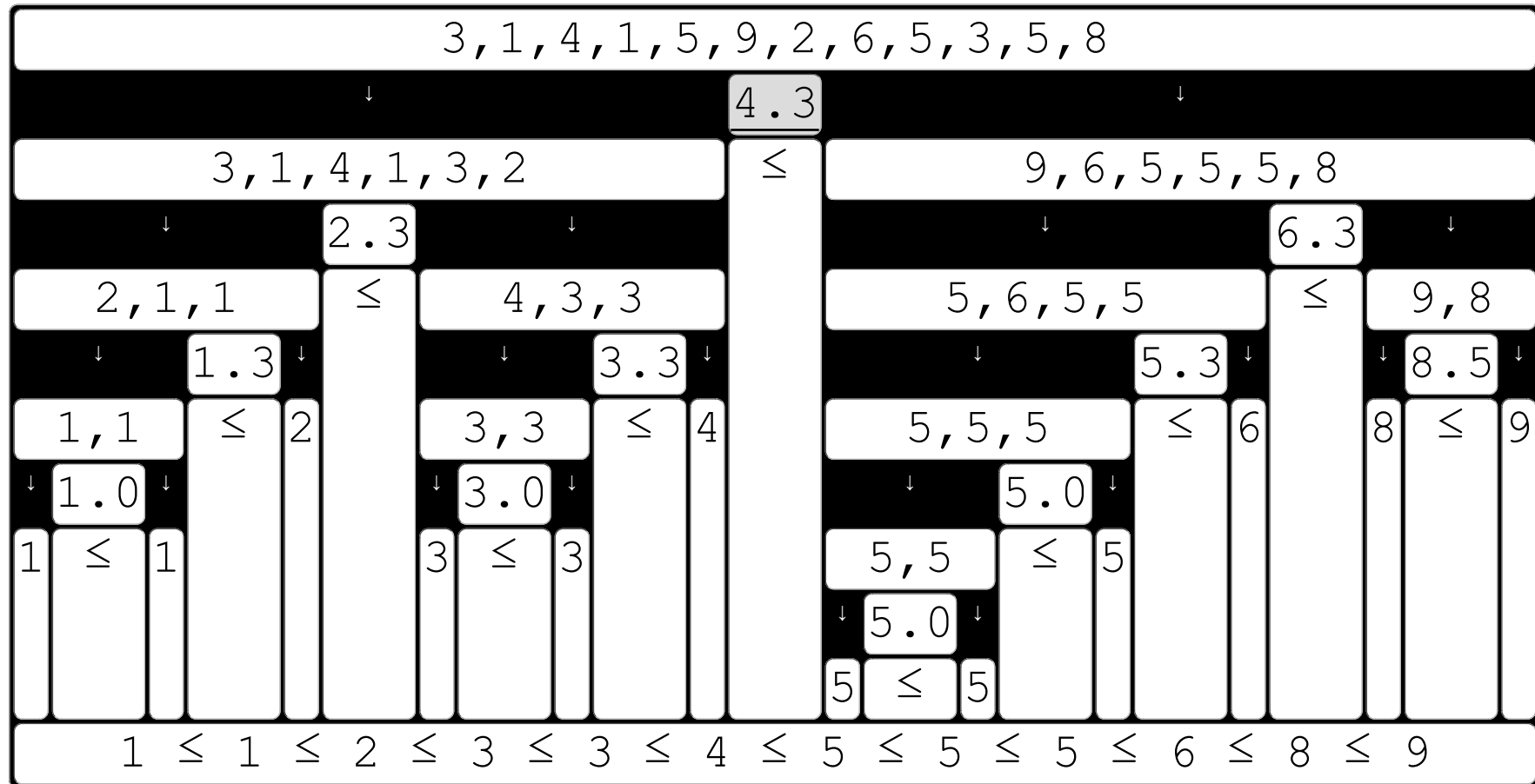


previous

Compared to [pivot](#) those two lists have: values [smaller than the pivot](#), and values [larger than the pivot](#).

first

next

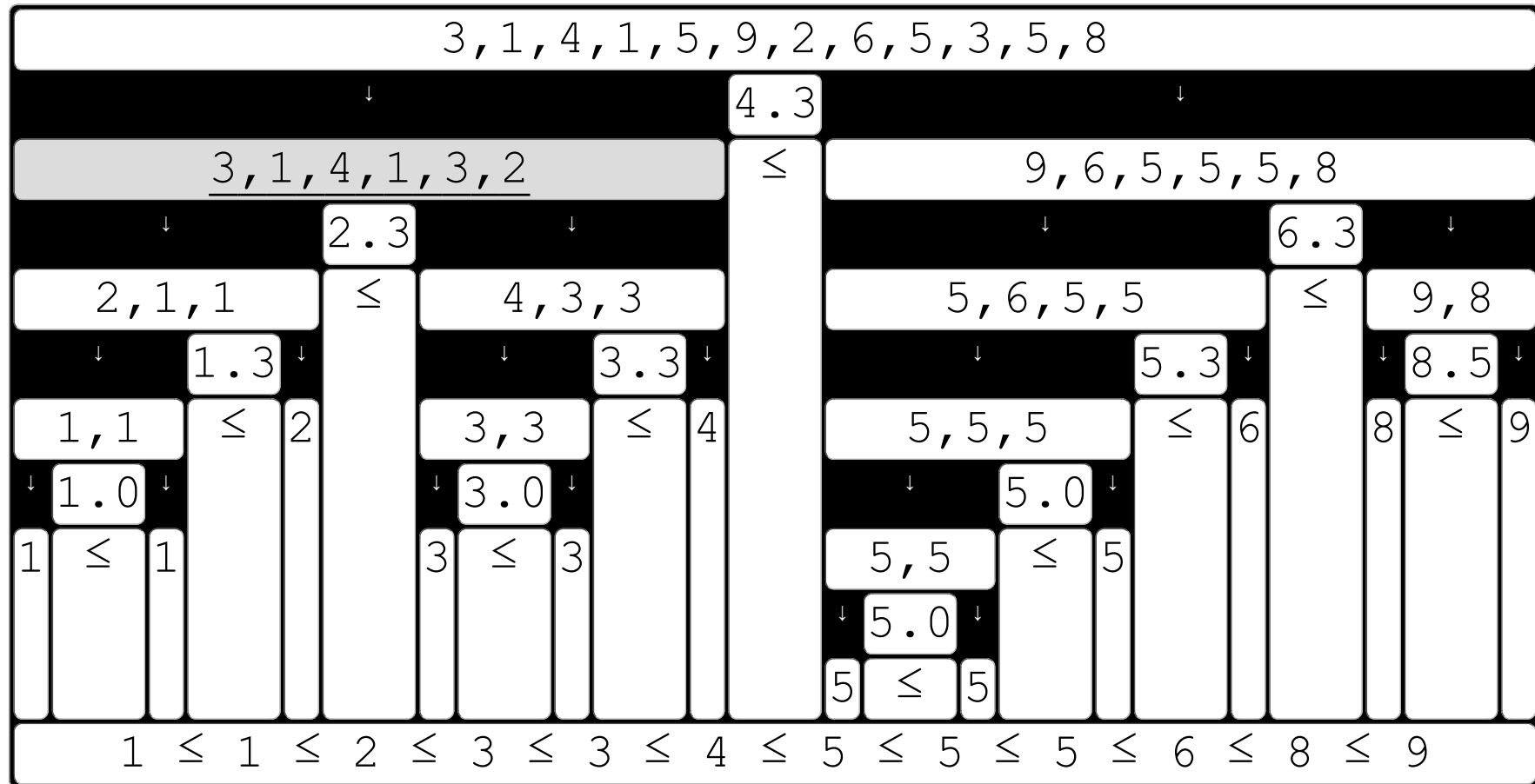


previous

Compared to [pivot](#) those two lists have: values smaller than the pivot, and values larger than the pivot.

first

next

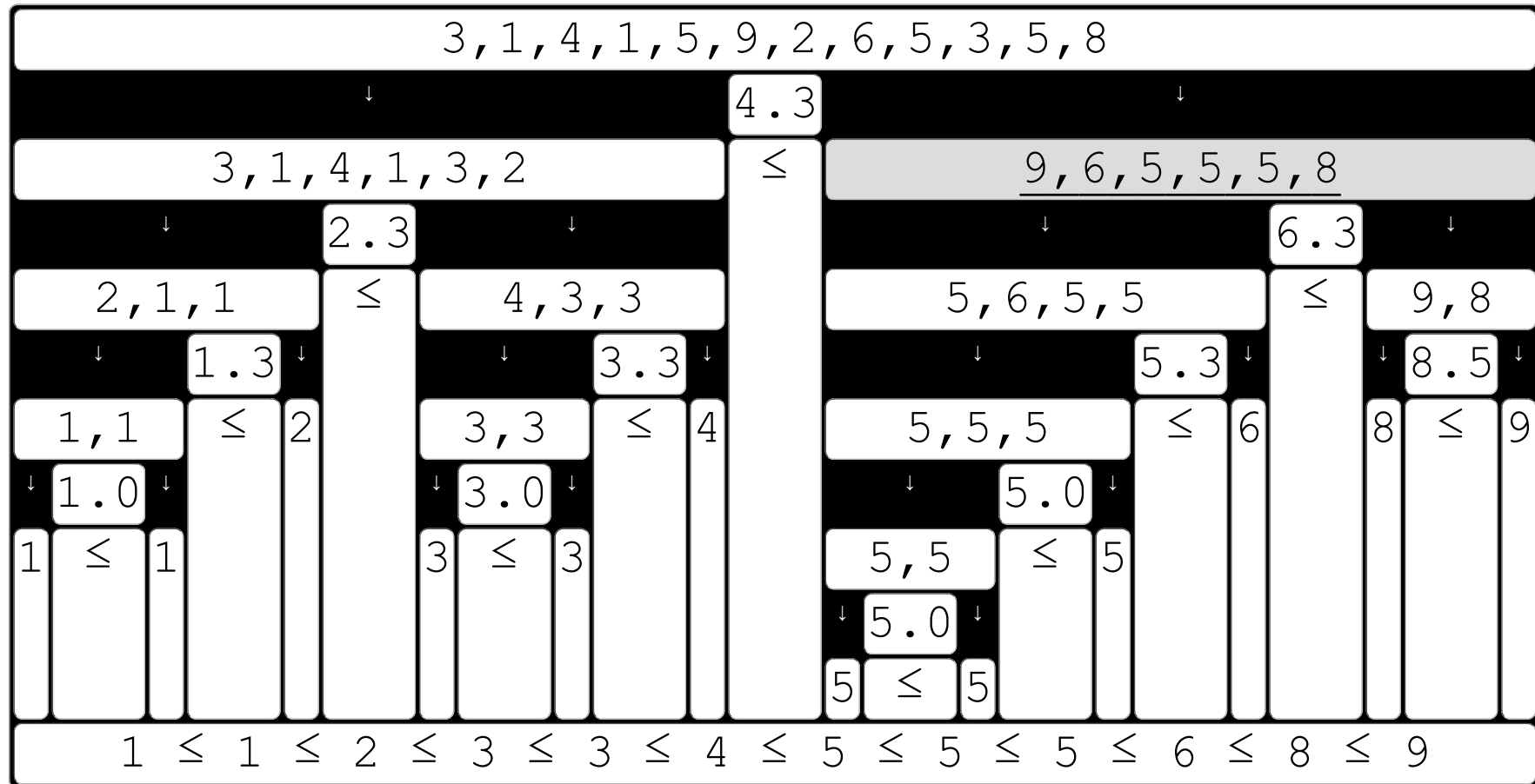


previous

Compared to pivot those two lists have: values [smaller than the pivot](#), and values larger than the pivot.

first

next

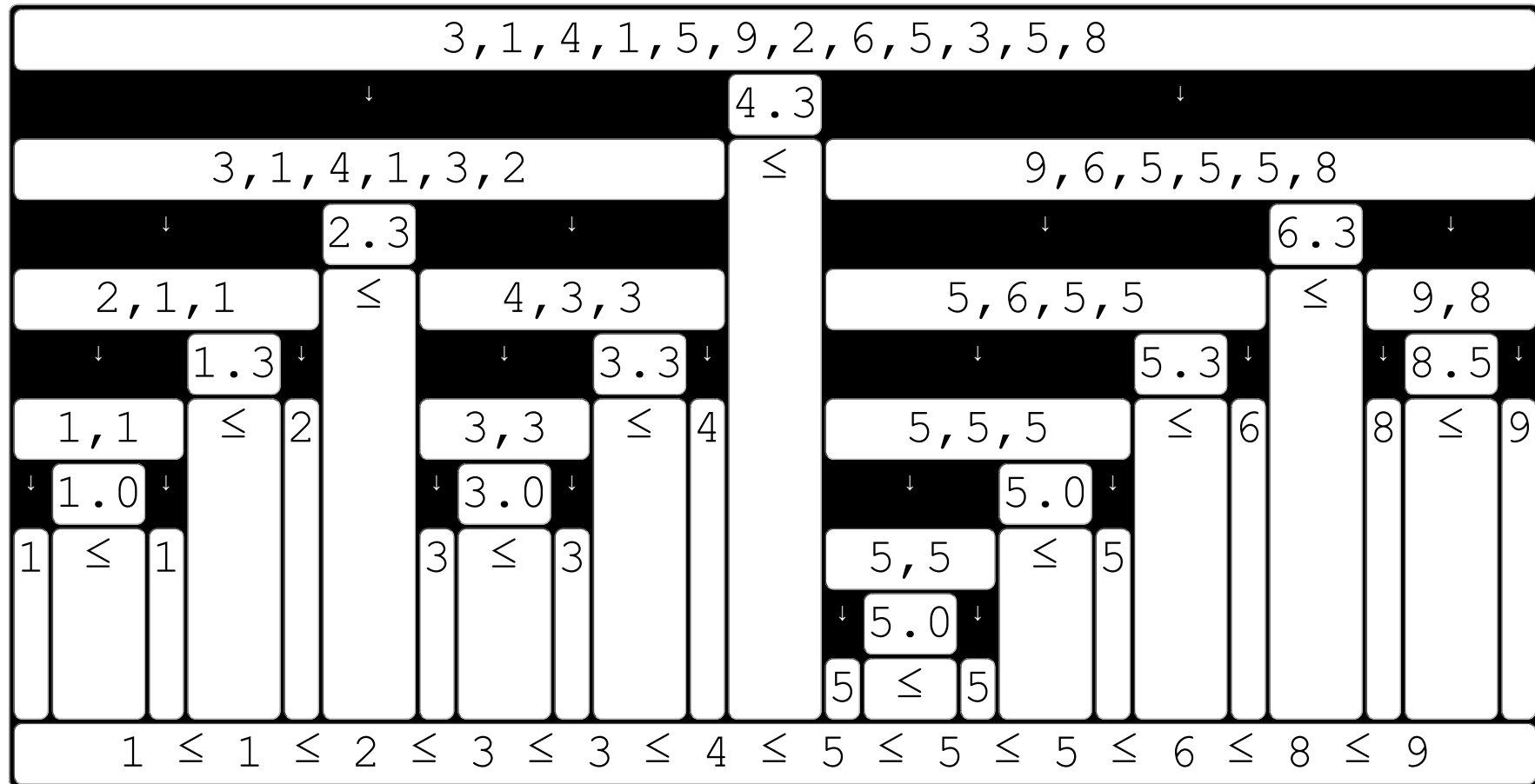


previous

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first

next

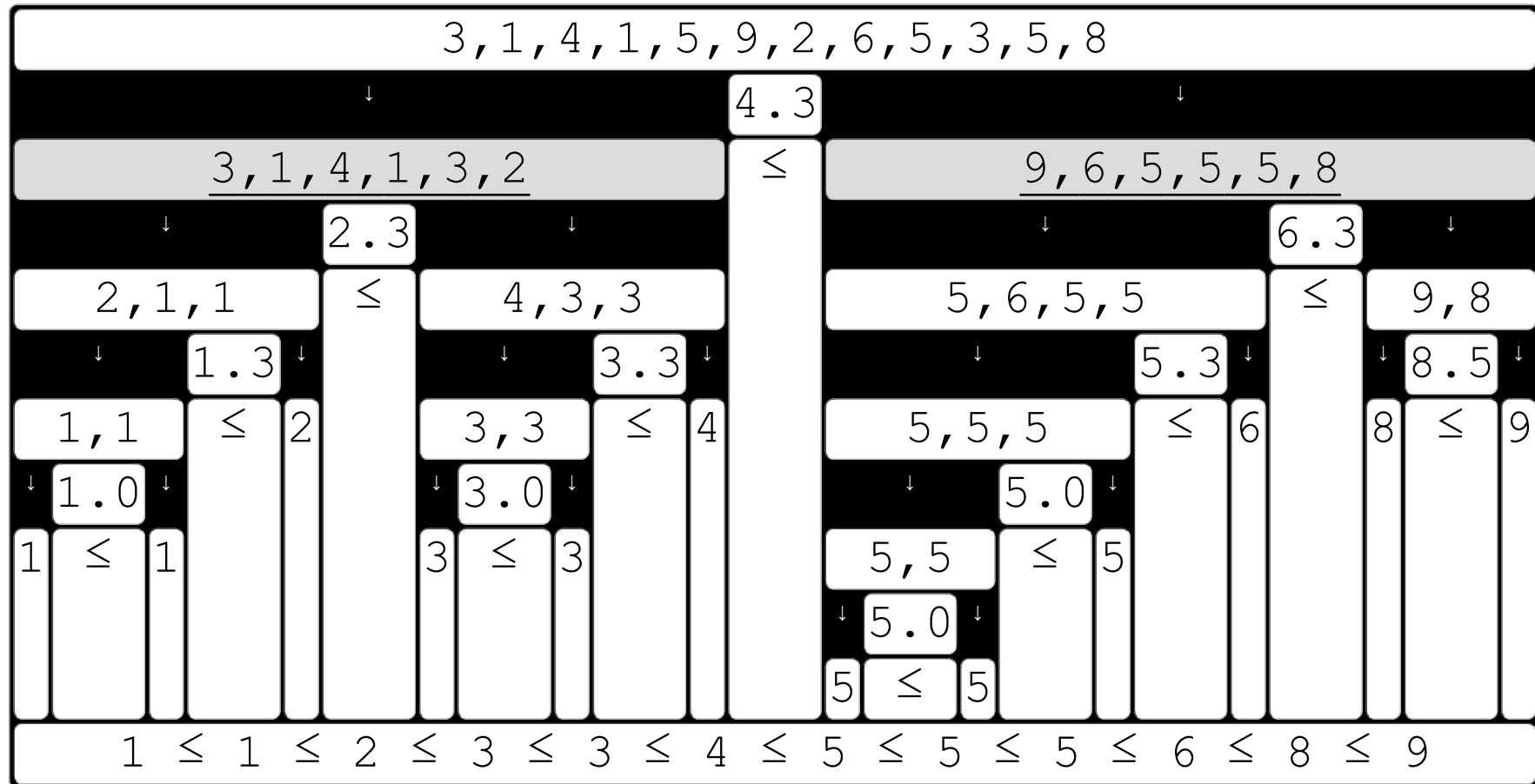


previous

At the second level, those [two lists](#) require their own pivots (that means: [two pivots](#)). The split procedure results in [four lists](#).

first

next

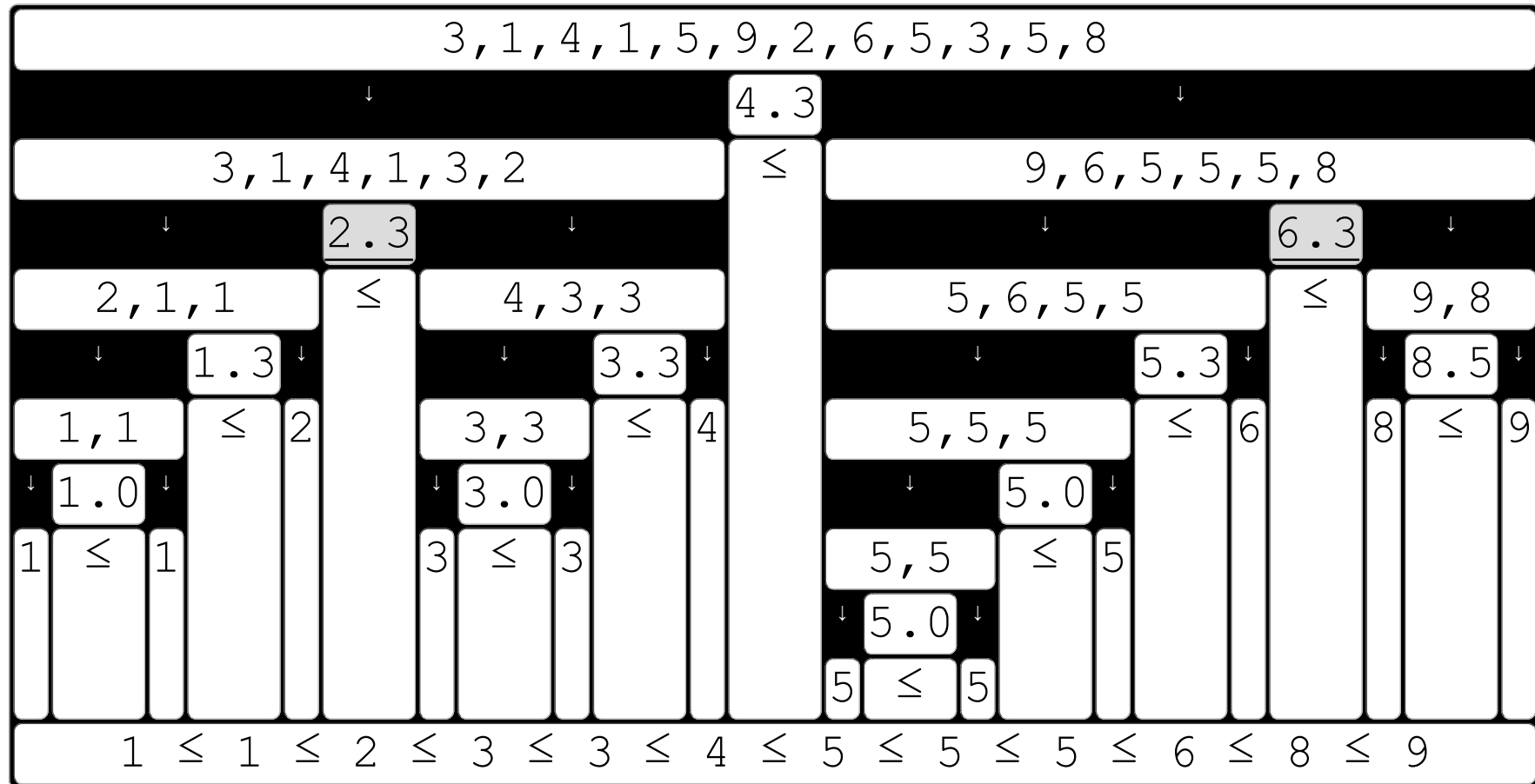


previous

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first

next

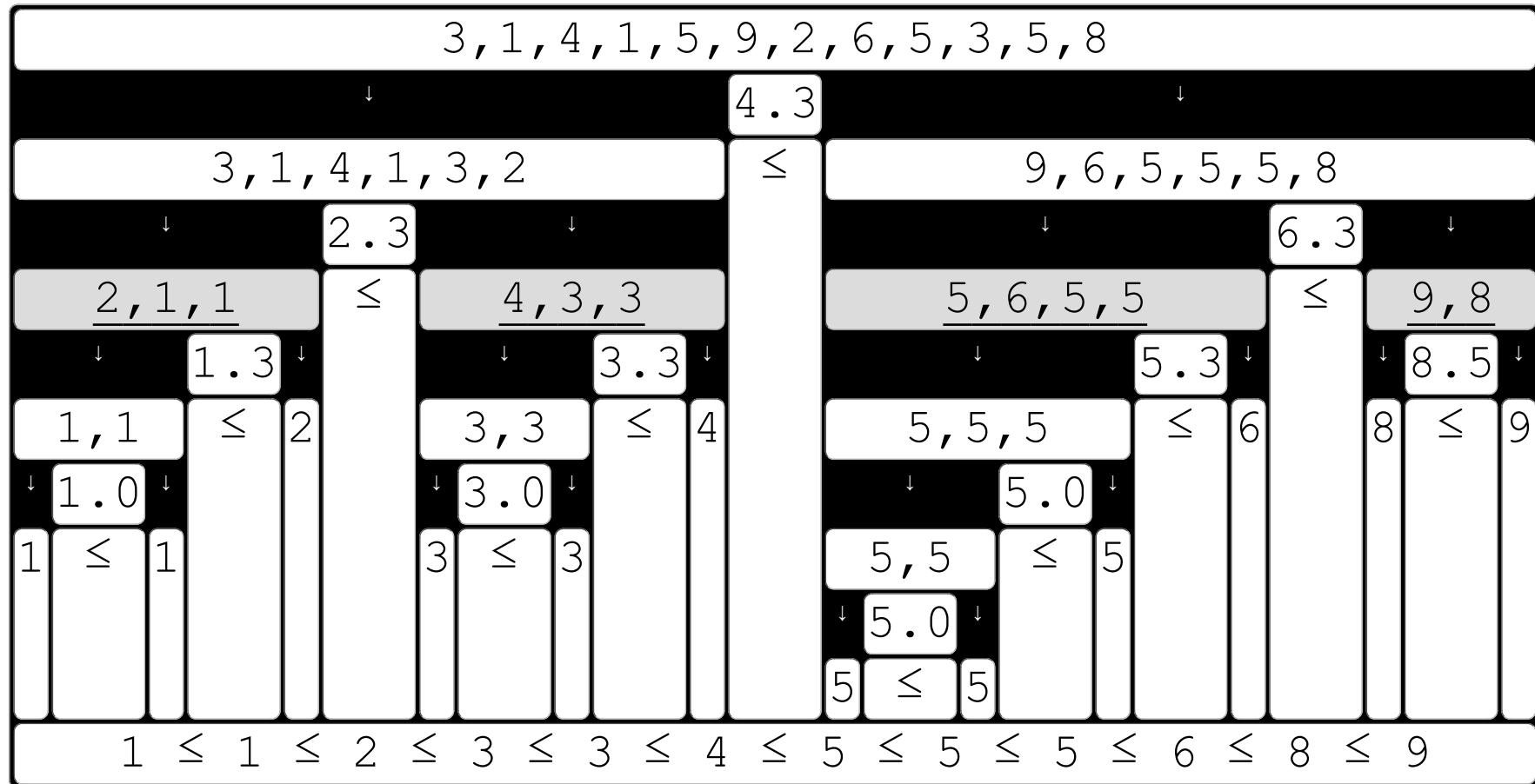


previous

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first

next

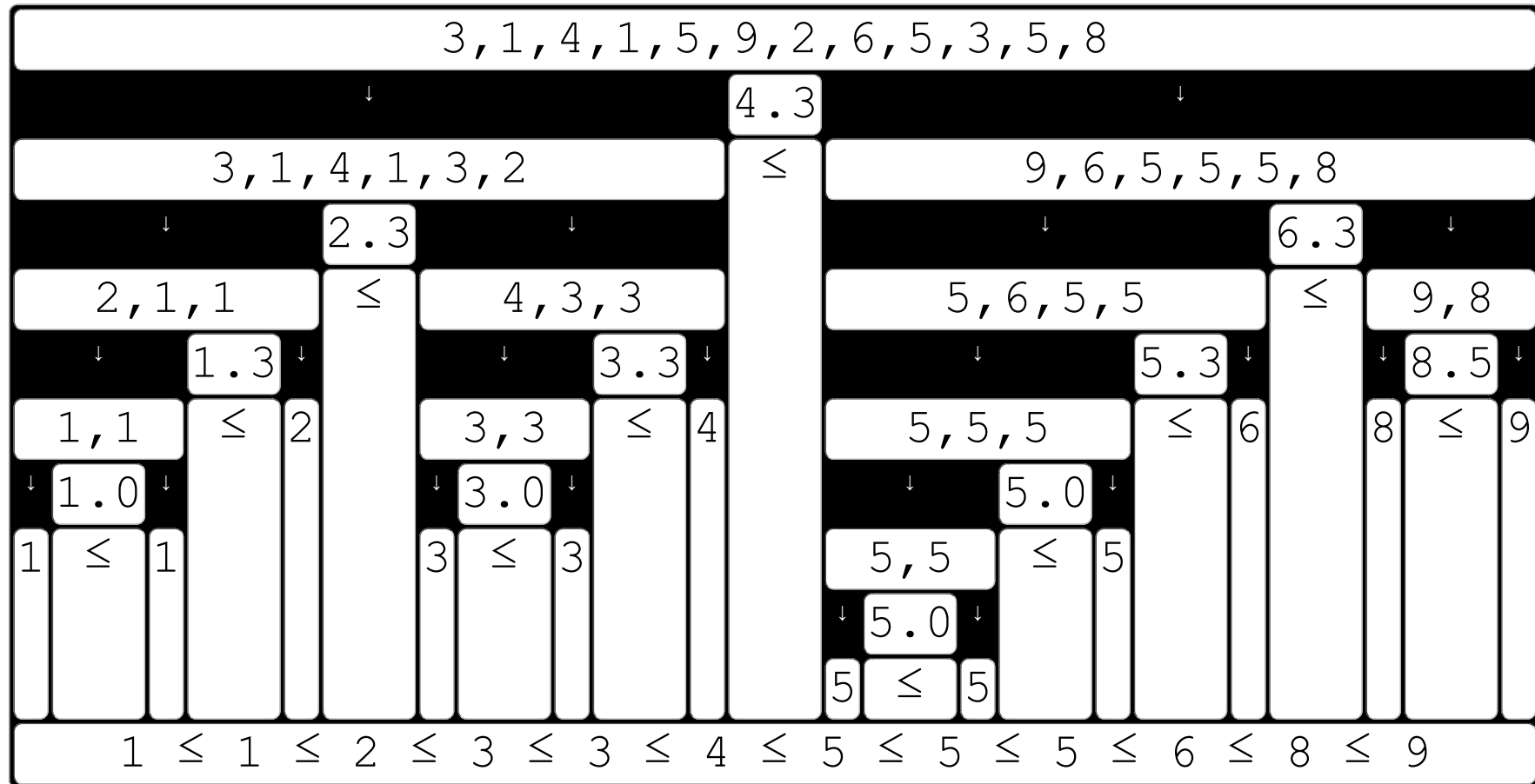


previous

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first

next

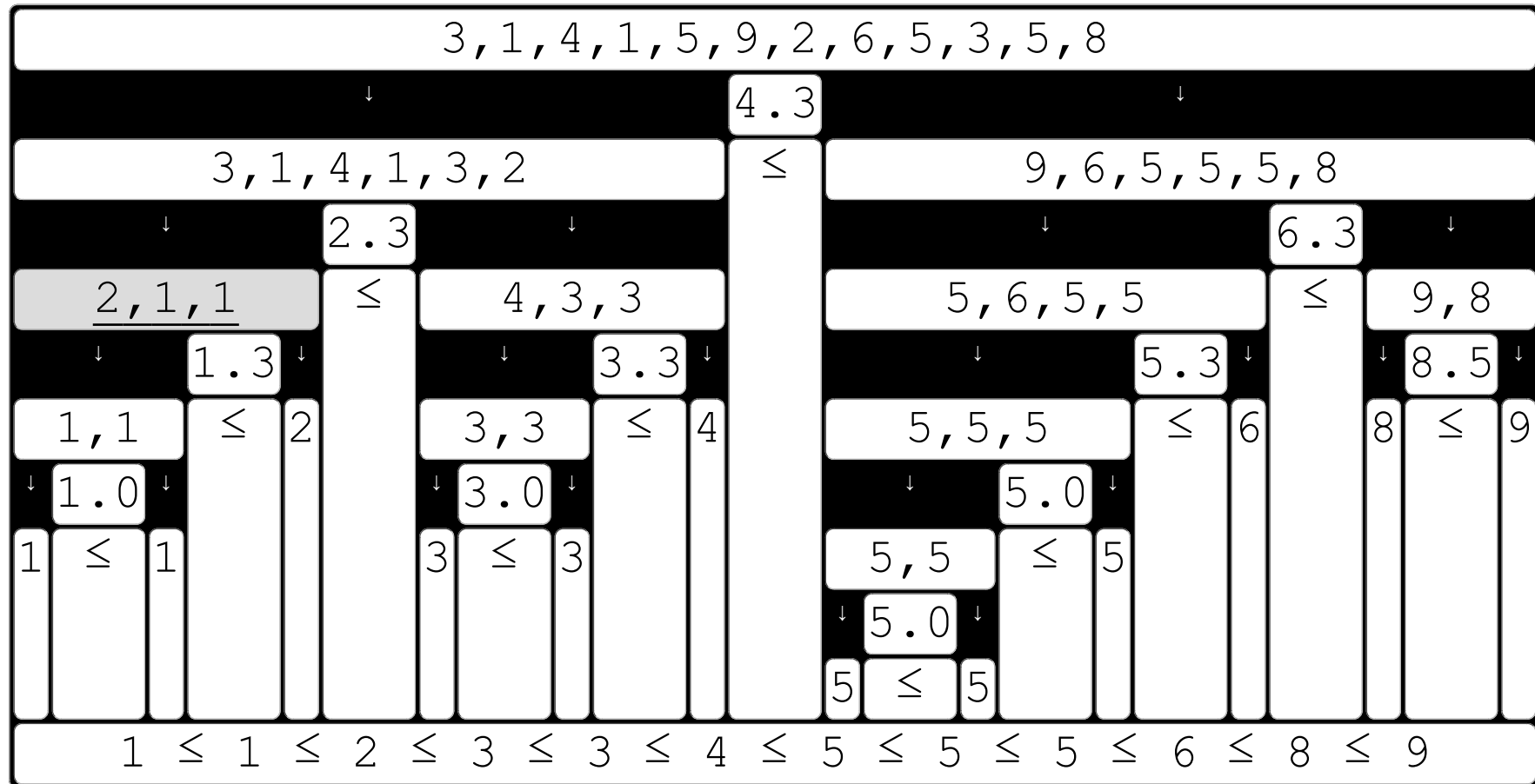


previous

first

The list at the left side has [the smallest values](#), the list at the right side has [the largest values](#), and [everything else](#) in between.

next

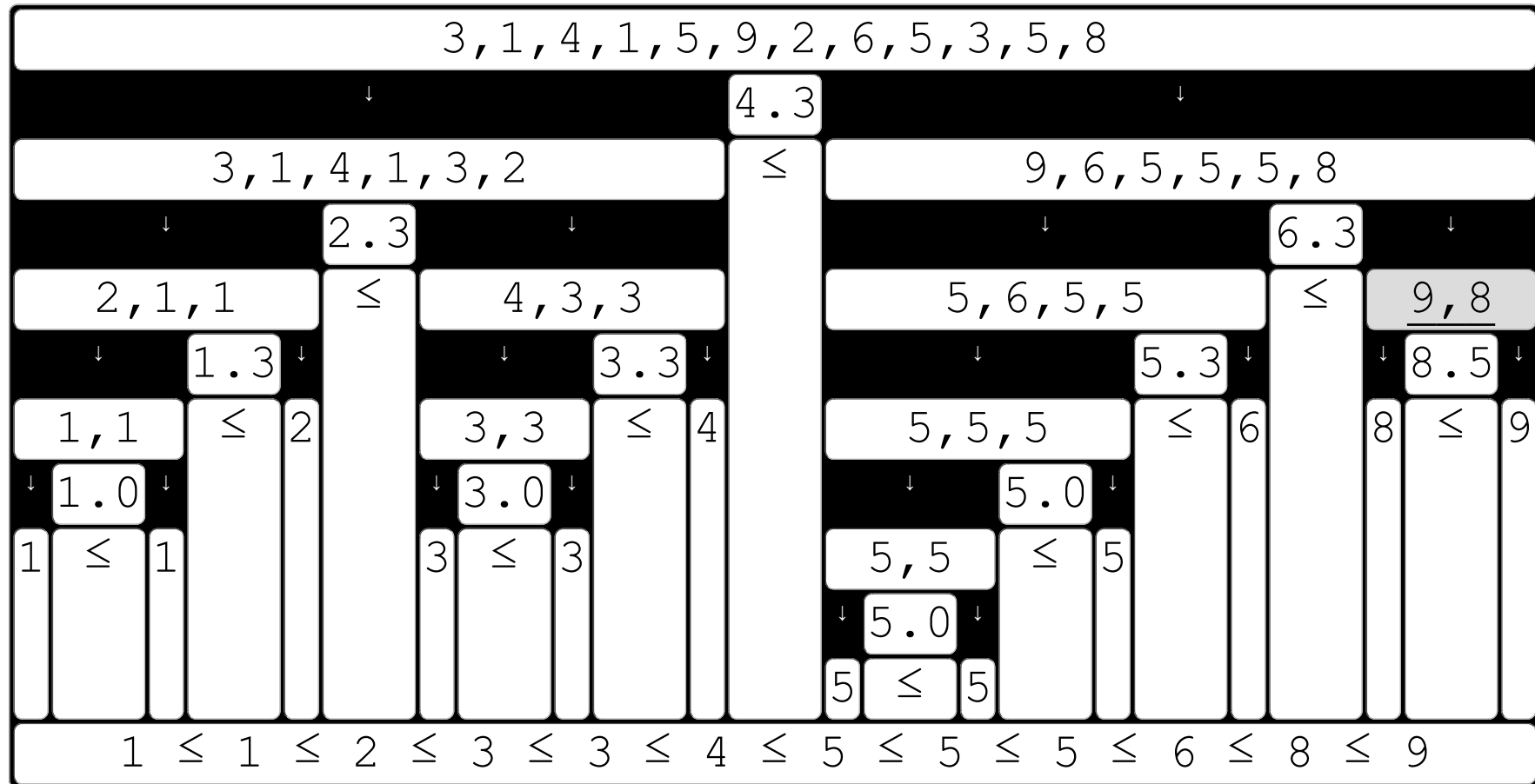


previous

first

The list at the left side has [the smallest values](#), the list at the right side has the largest values, and everything else in between.

next

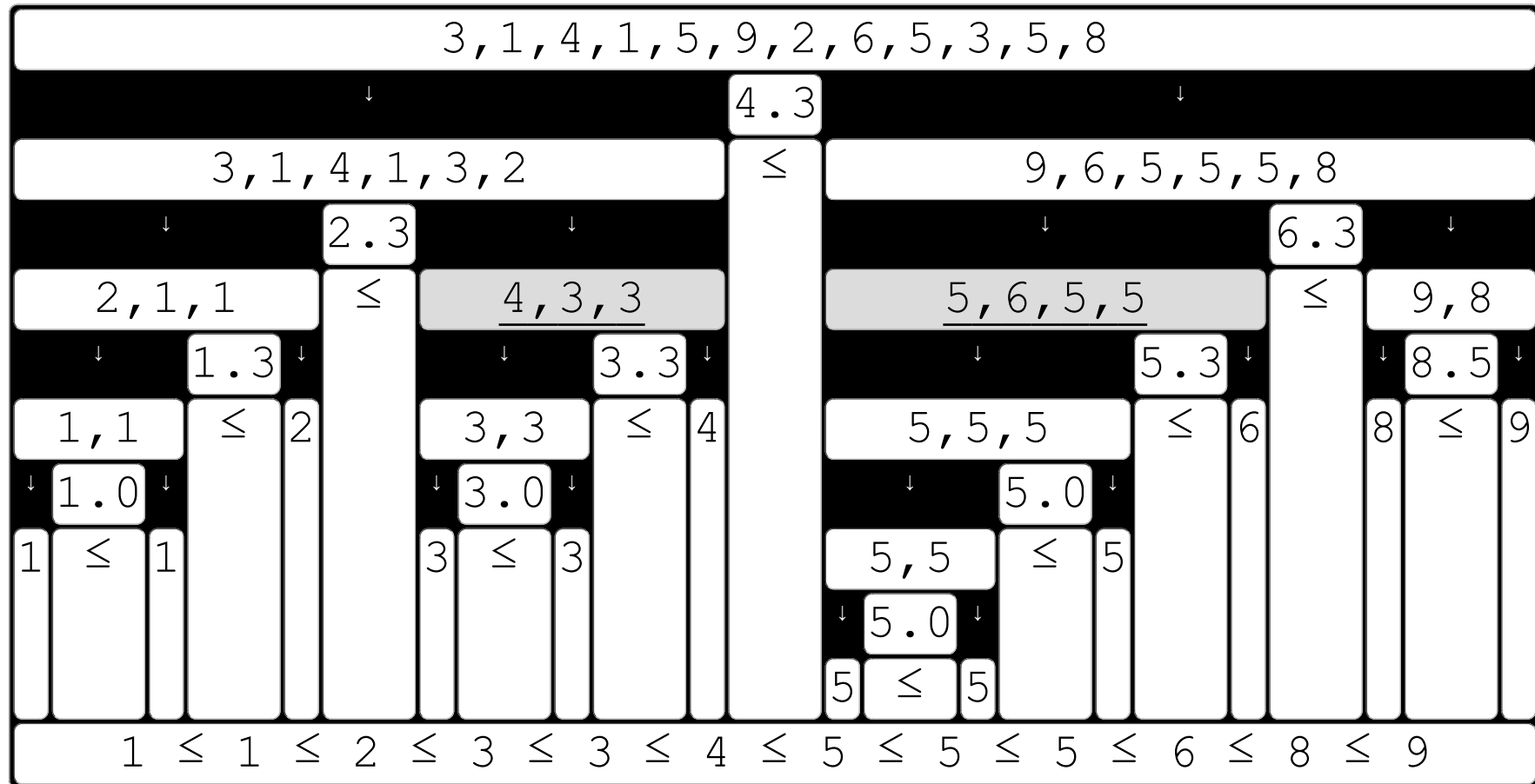


previous

first

The list at the left side has the smallest values, the list at the right side has [the largest values](#), and everything else in between.

next

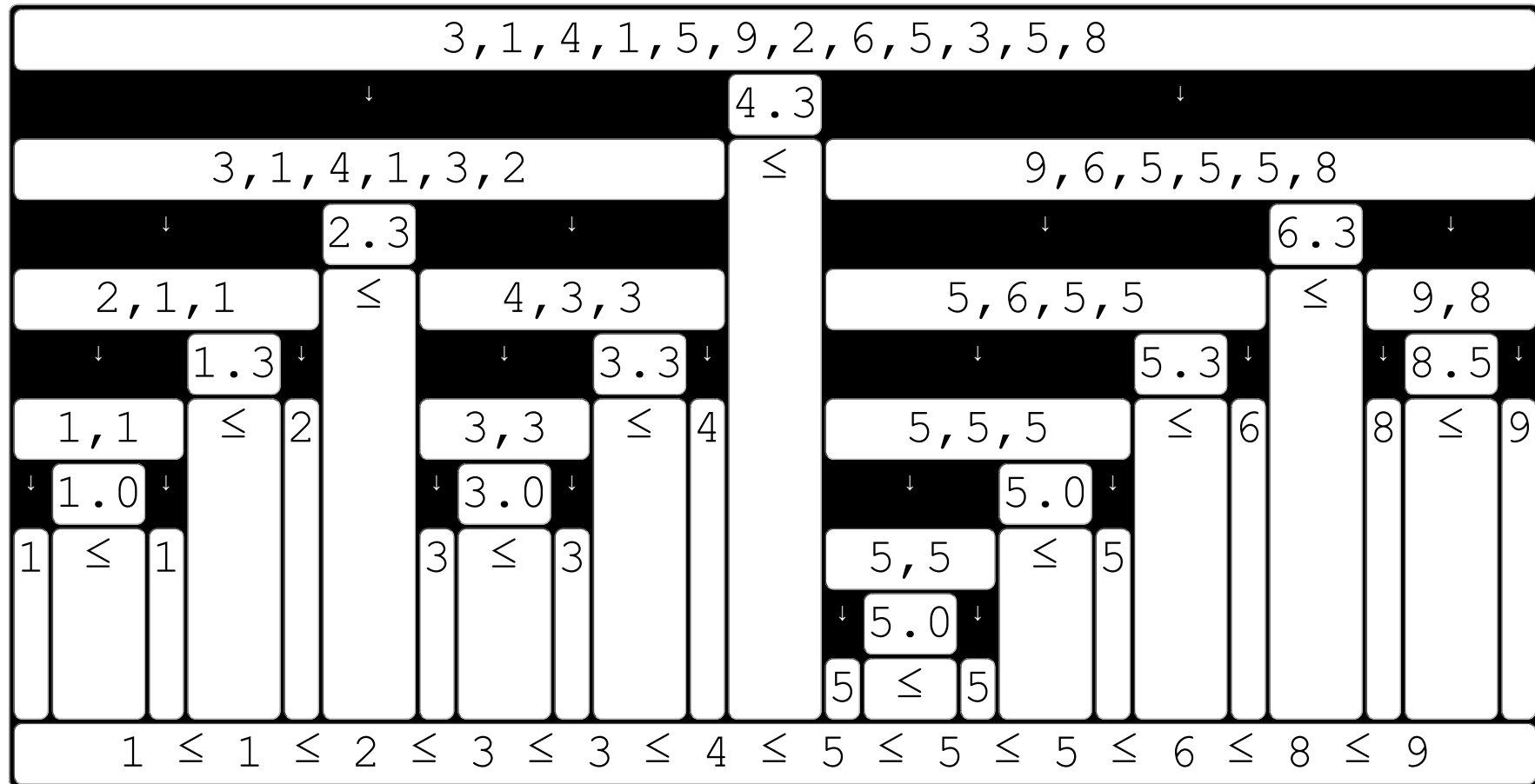


previous

first

The list at the left side has the smallest values, the list at the right side has the largest values, and [everything else](#) in between.

next

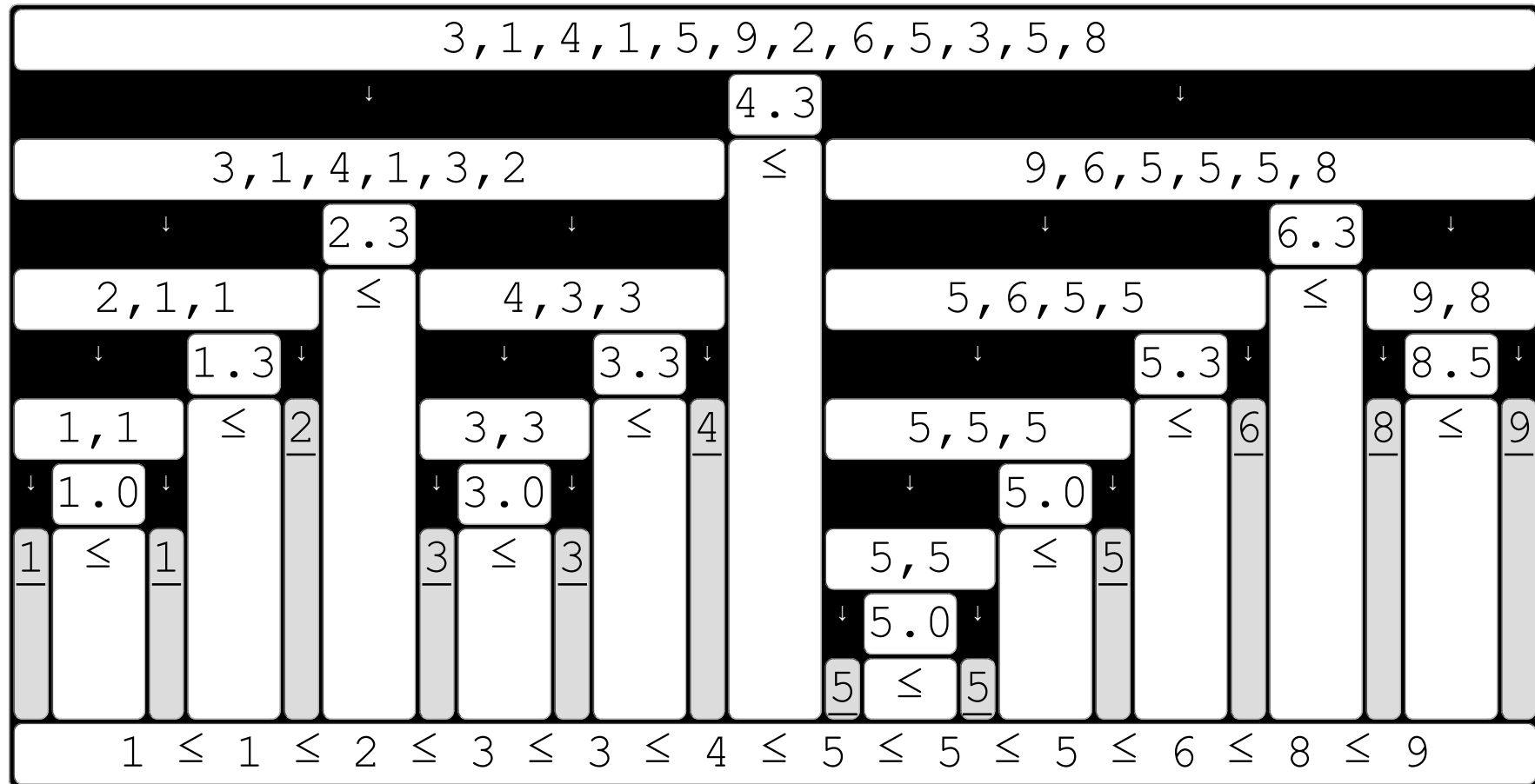


previous

The algorithm repeats this procedure until [each list has the size of only one value](#). Those values are [the result of the sort algorithm](#).

first

next

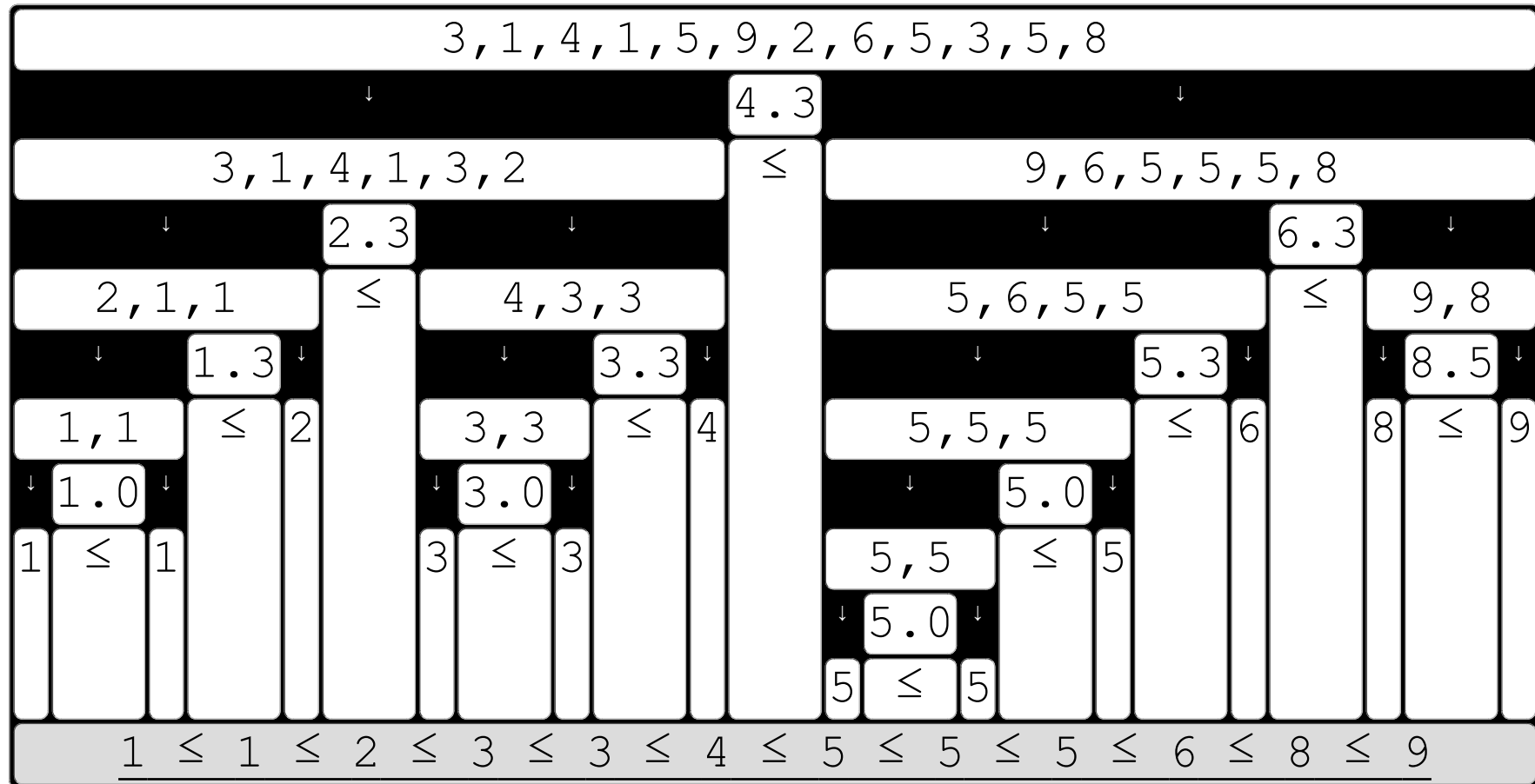


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first

next

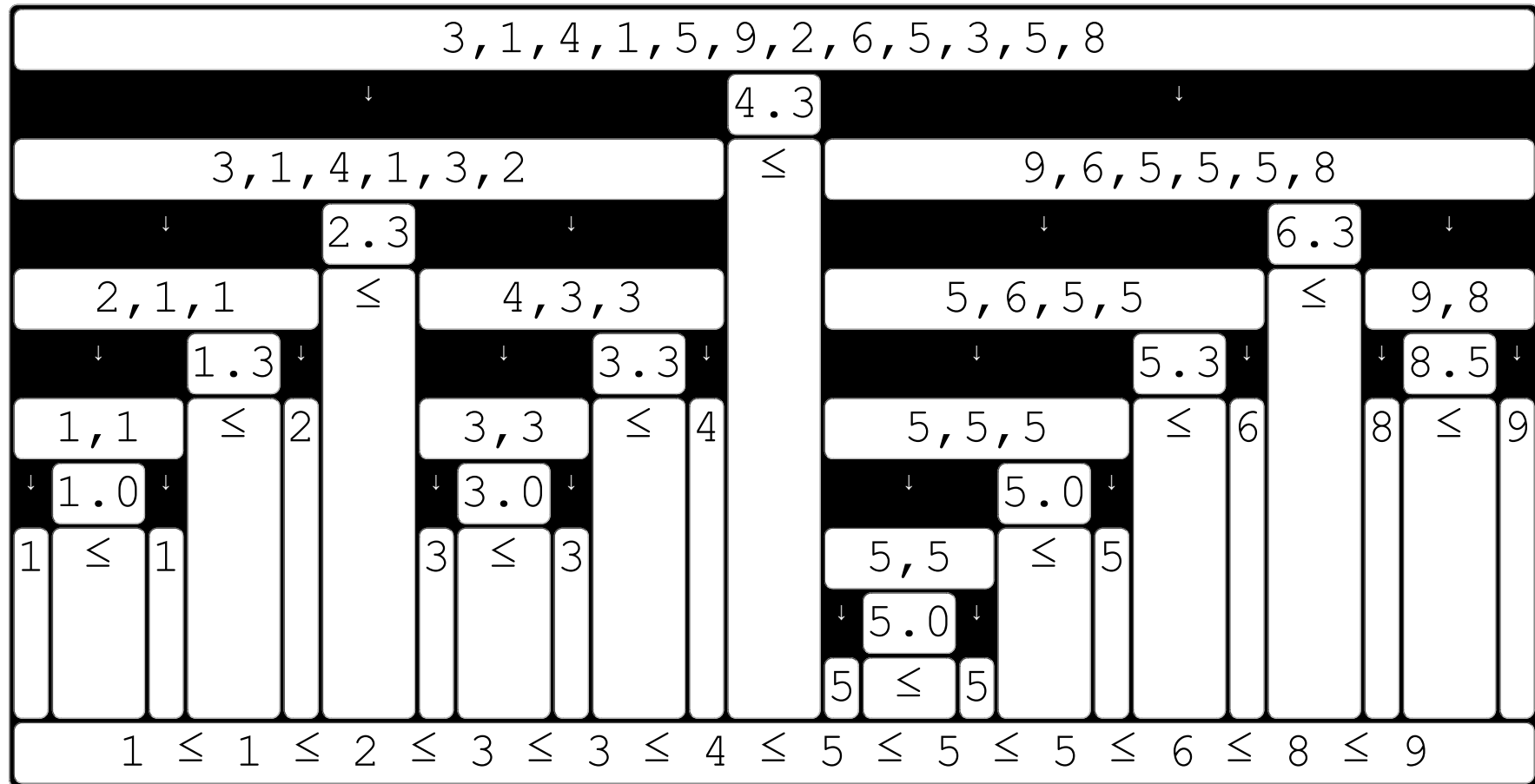


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first

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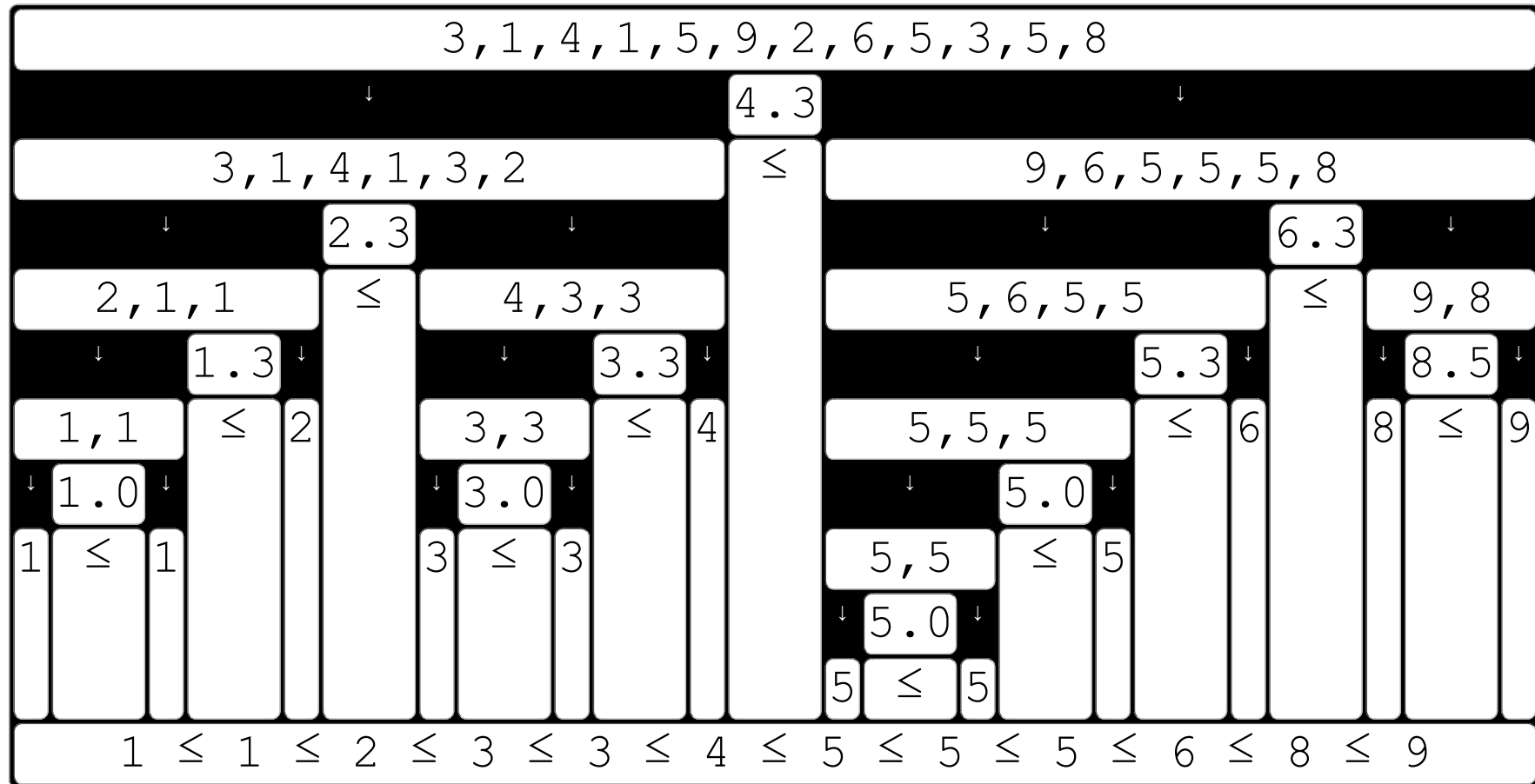


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first

"Recursive Algorithm" is the name of this kind of repetitive procedure.

next

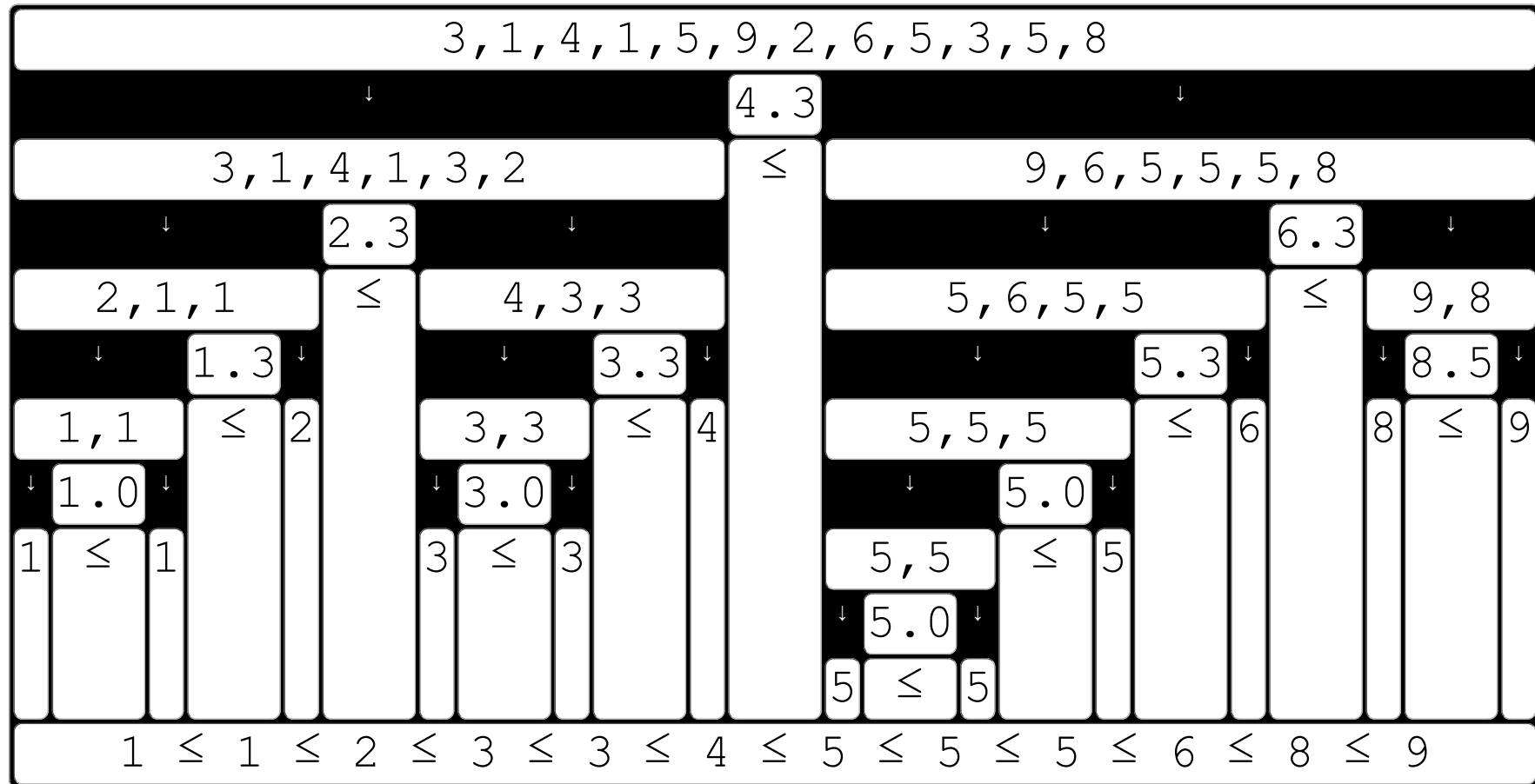


previous

first

(3) Advanced concepts: in-place.

next

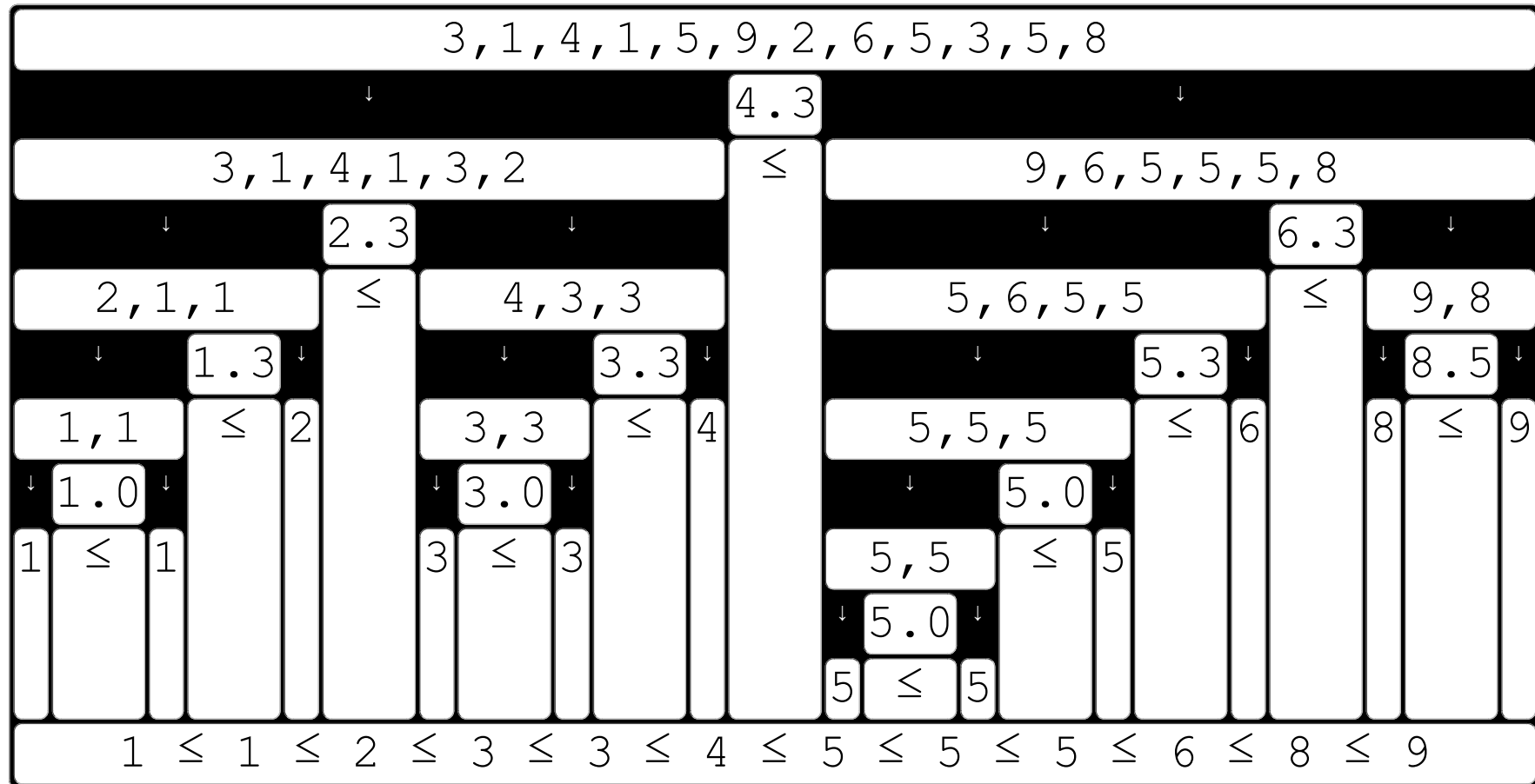


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first

Quicksort is an efficient sort algorithm, and one of the few in-place algorithms (in this category).

next

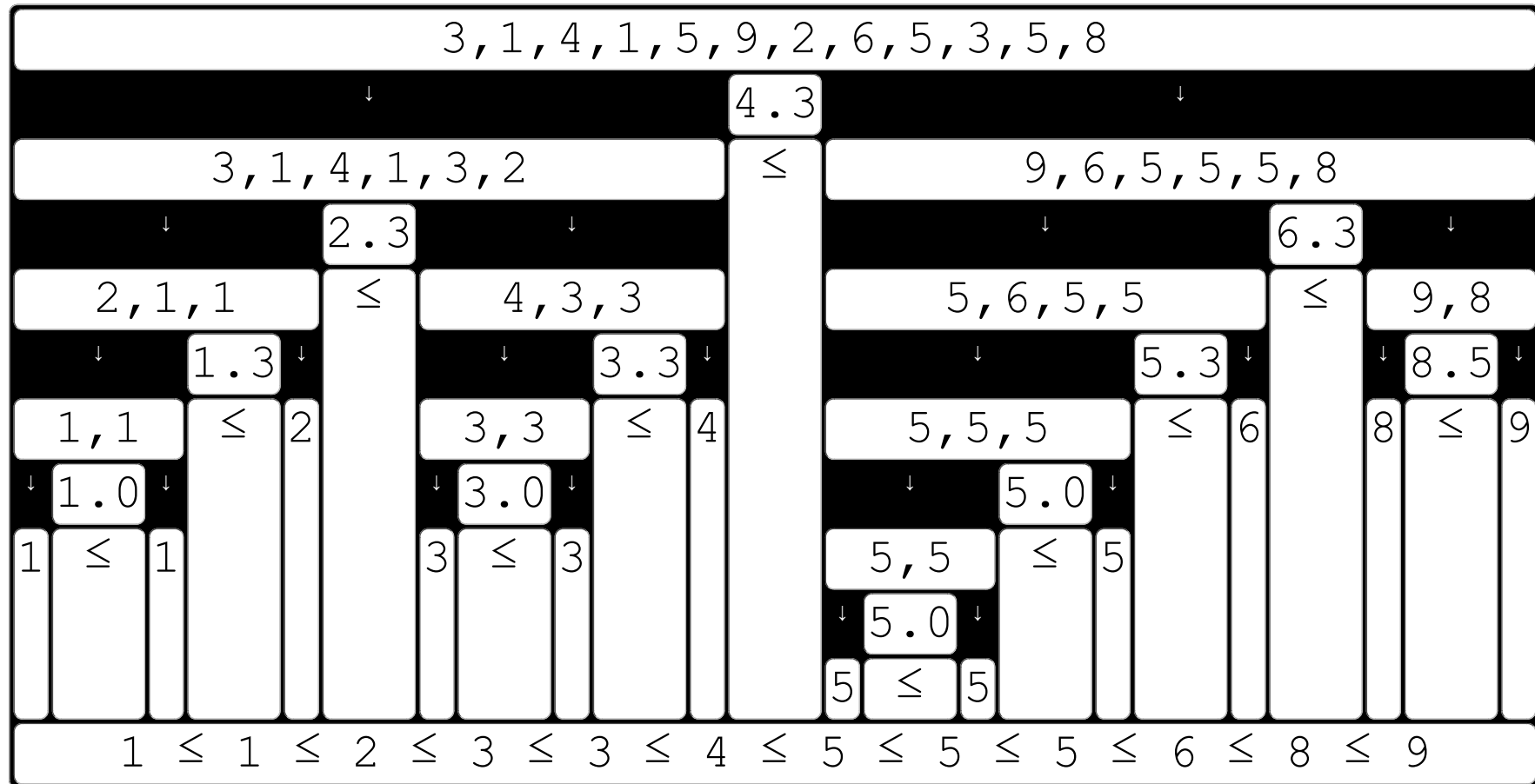


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first

In-place means: the algorithm does not require additional storage memory. Given two memory positions, Quicksort exchanges its values.

next

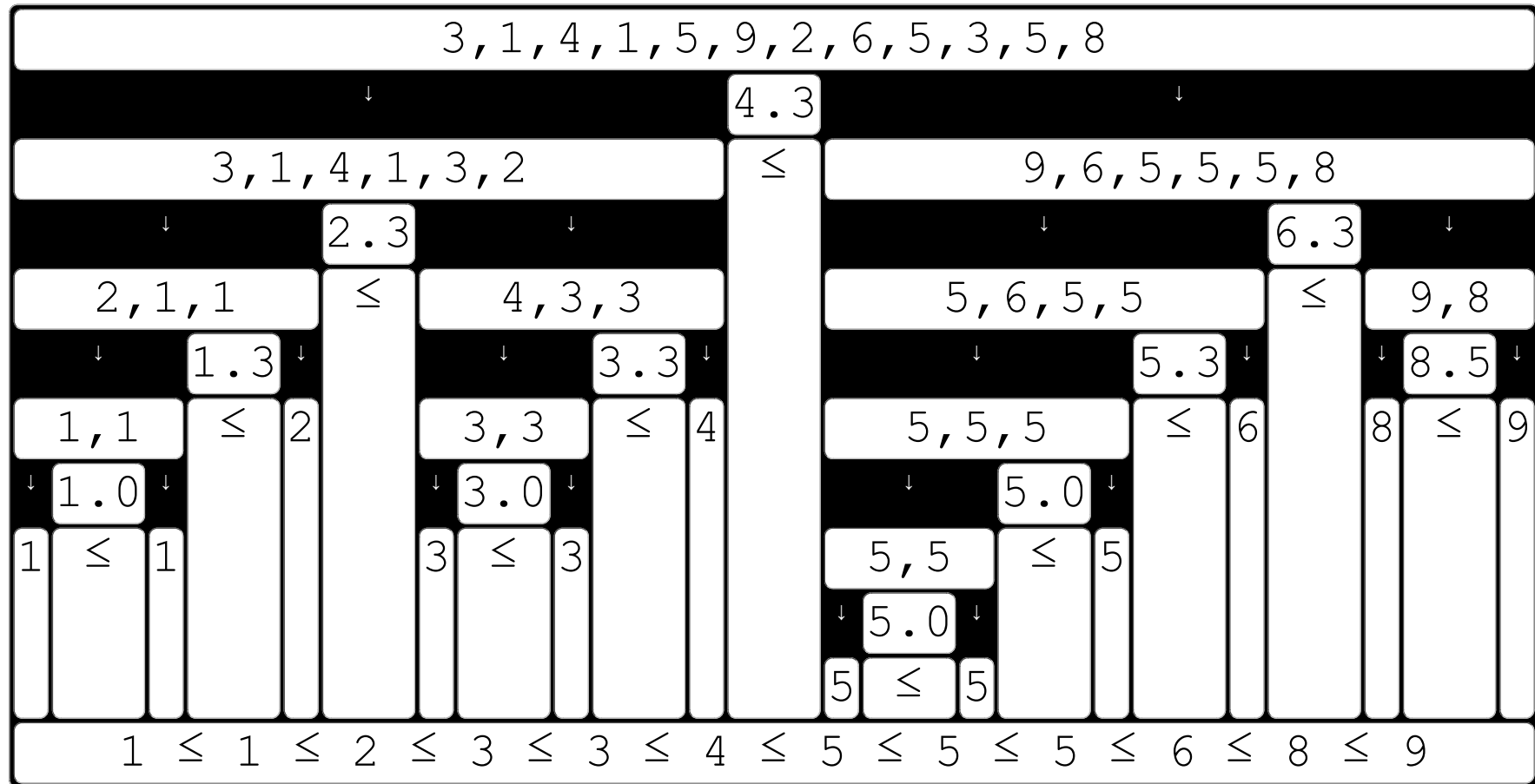


previous

first

A naive approach would create (at the split procedure) two lists with the same order of previous iteration.

next

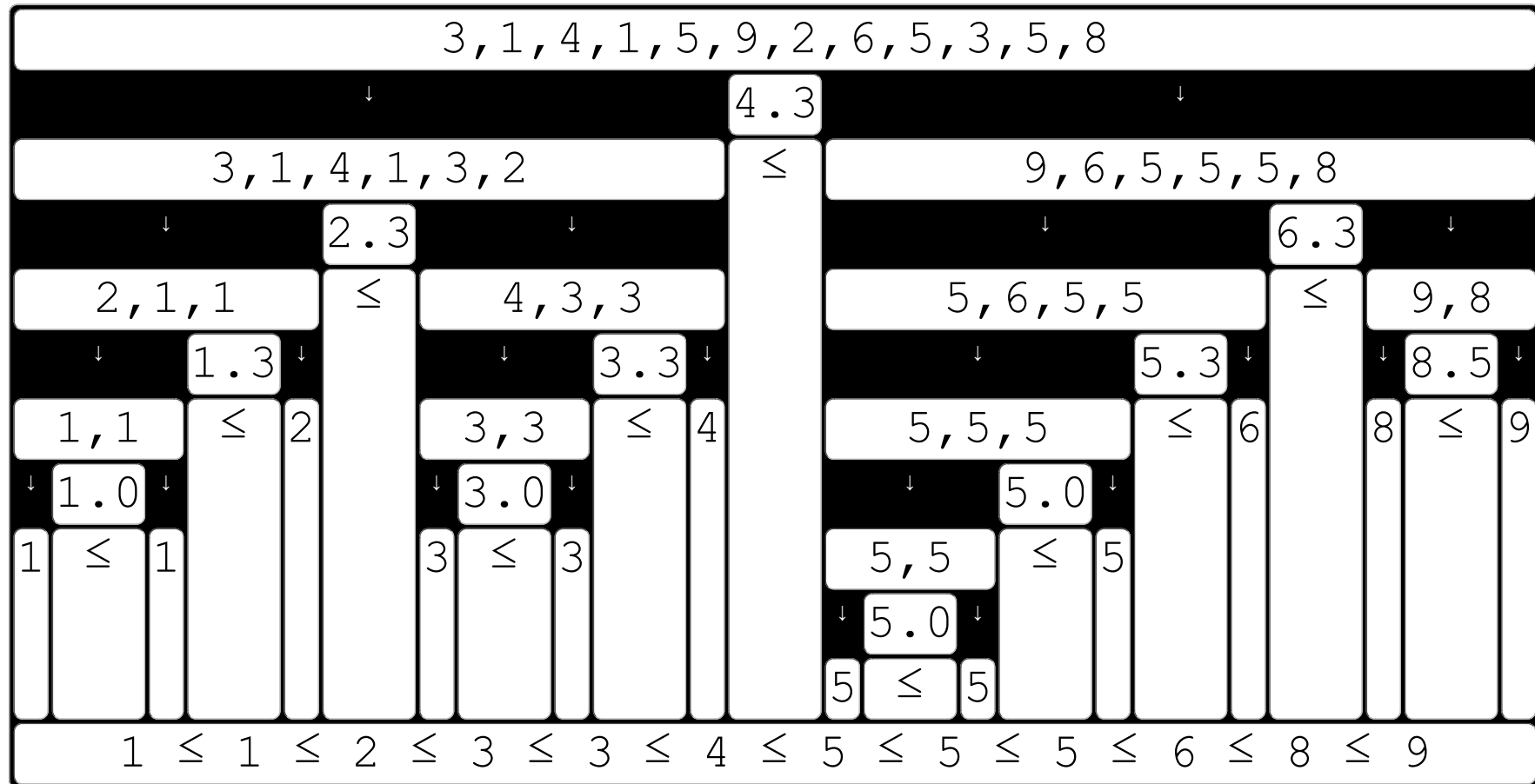


previous

As a matter of fact, QuickSort does not create any new list, it reuses the same storage space.

first

next

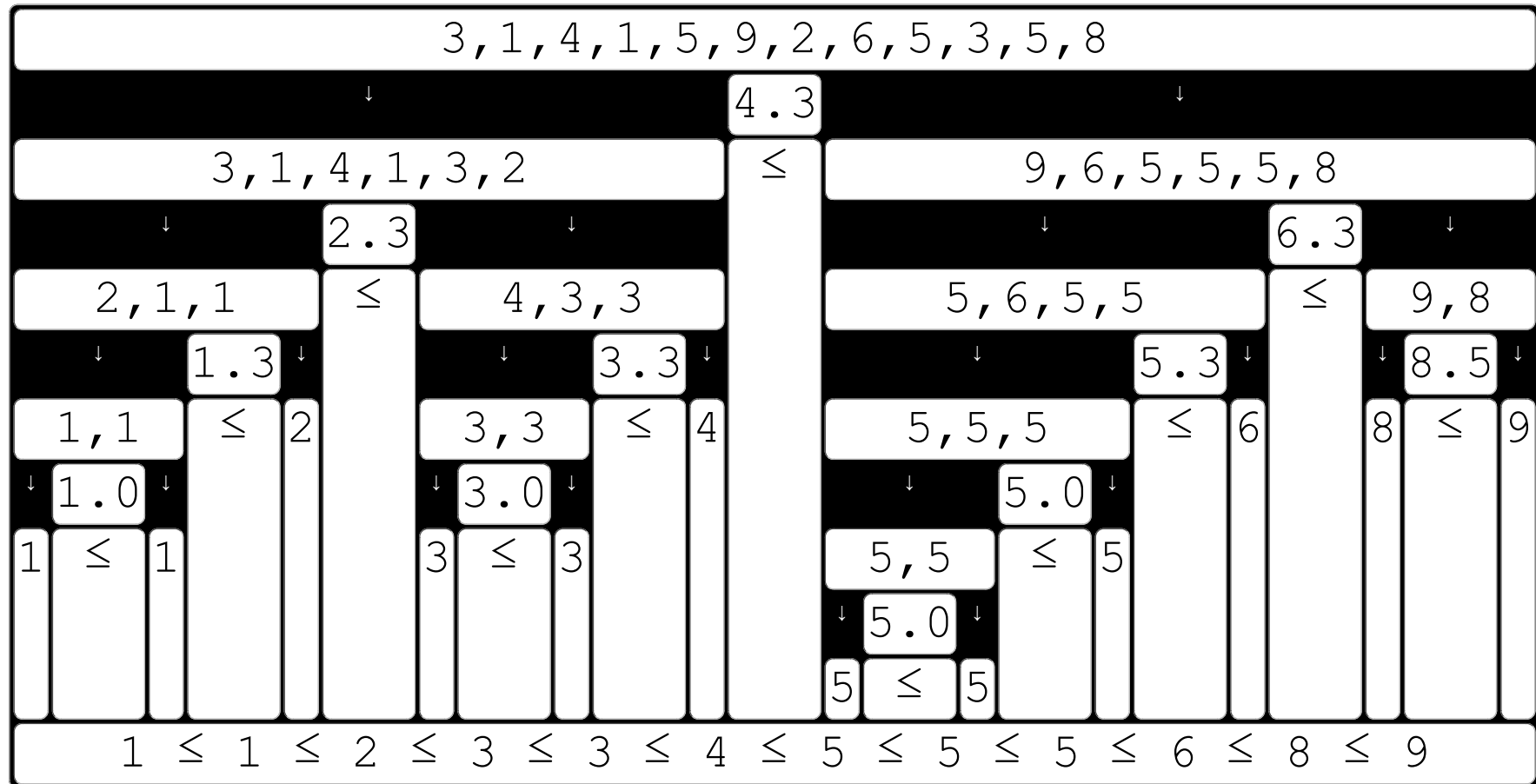


previous

first

So, usually (at the split procedure) the order of equal values is not the same as the previous iteration.

next

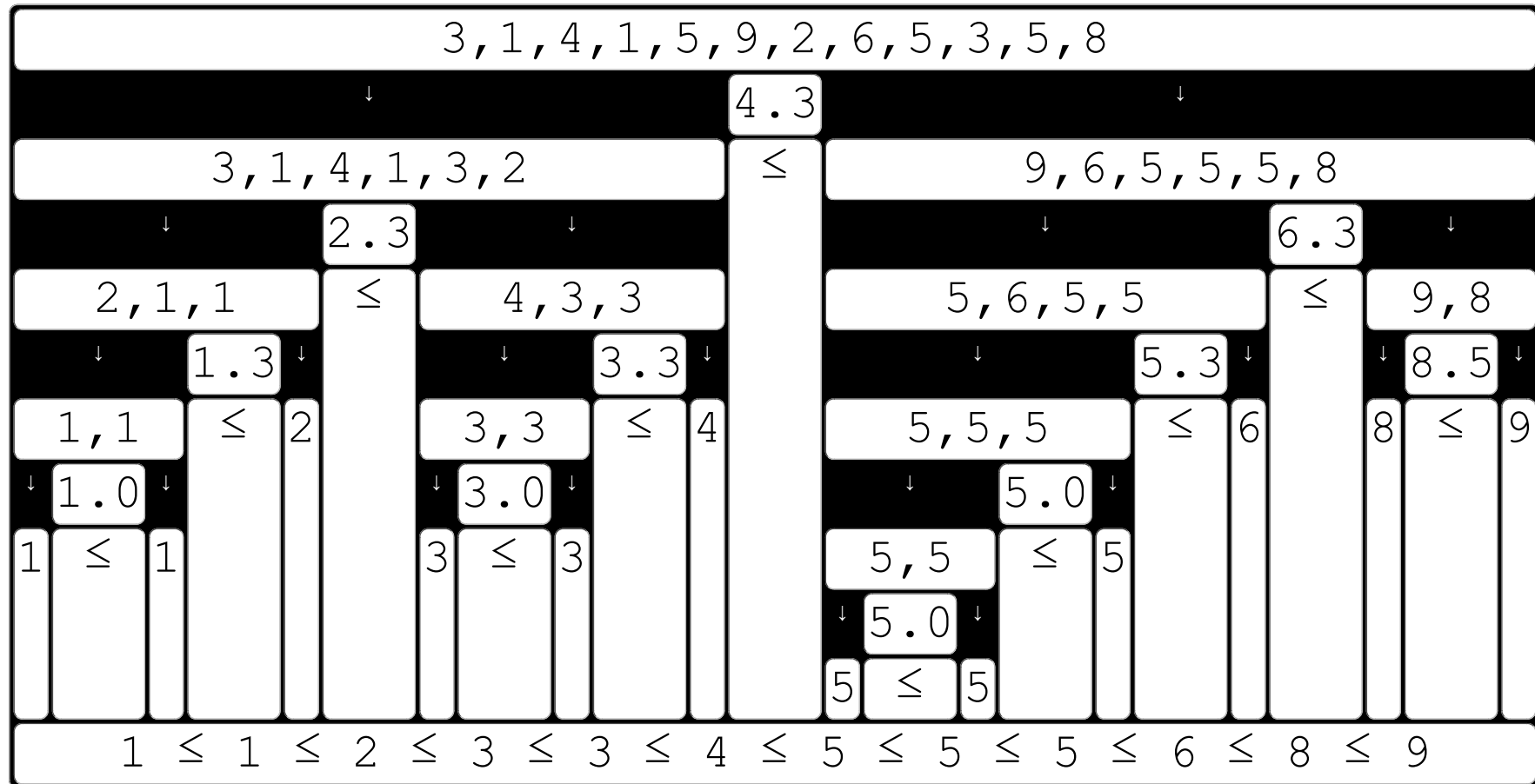


That makes Quicksort an unstable sort algorithm.

previous

first

next

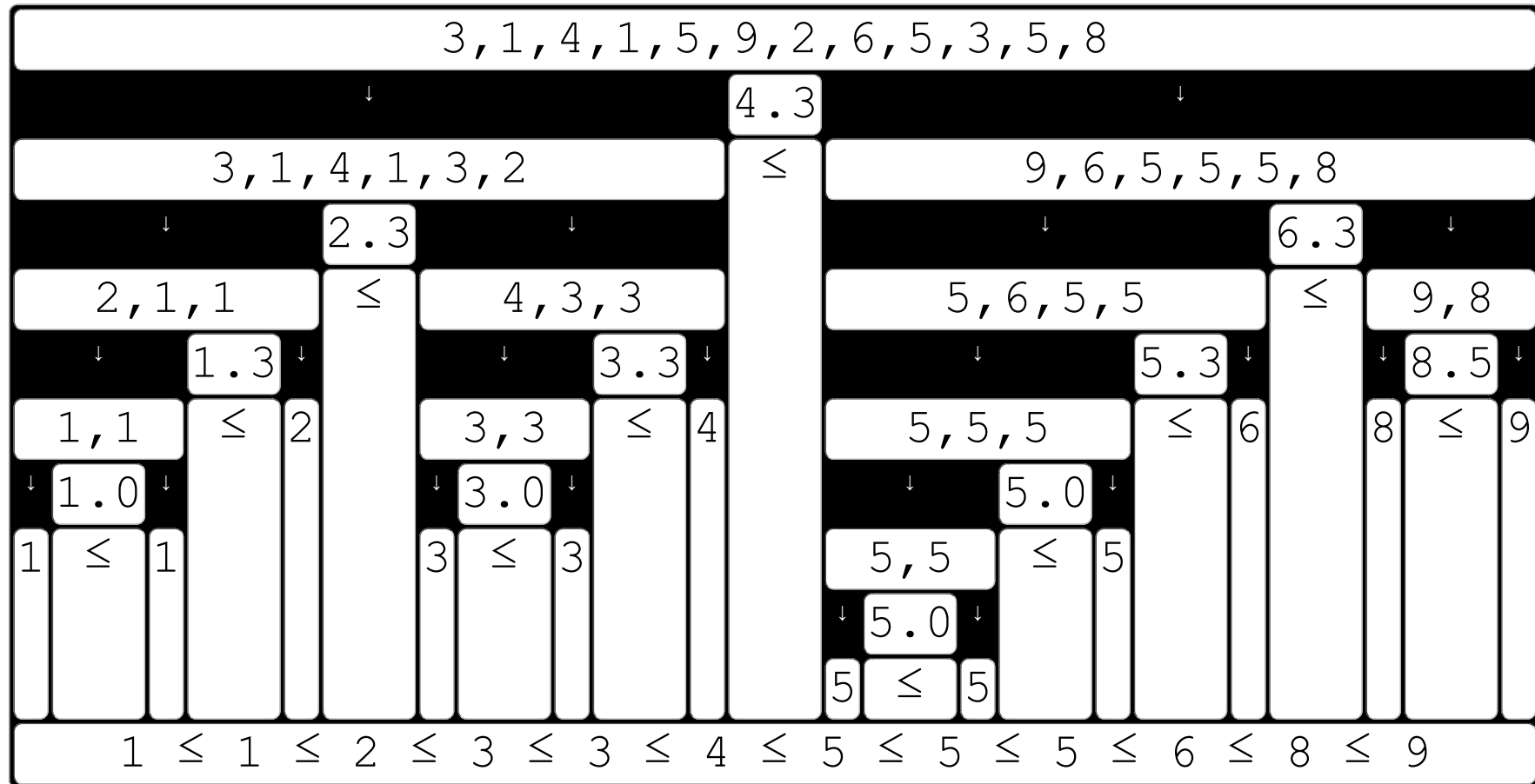


previous

first

(4) Advanced concepts: stable / unstable sort algorithm.

next

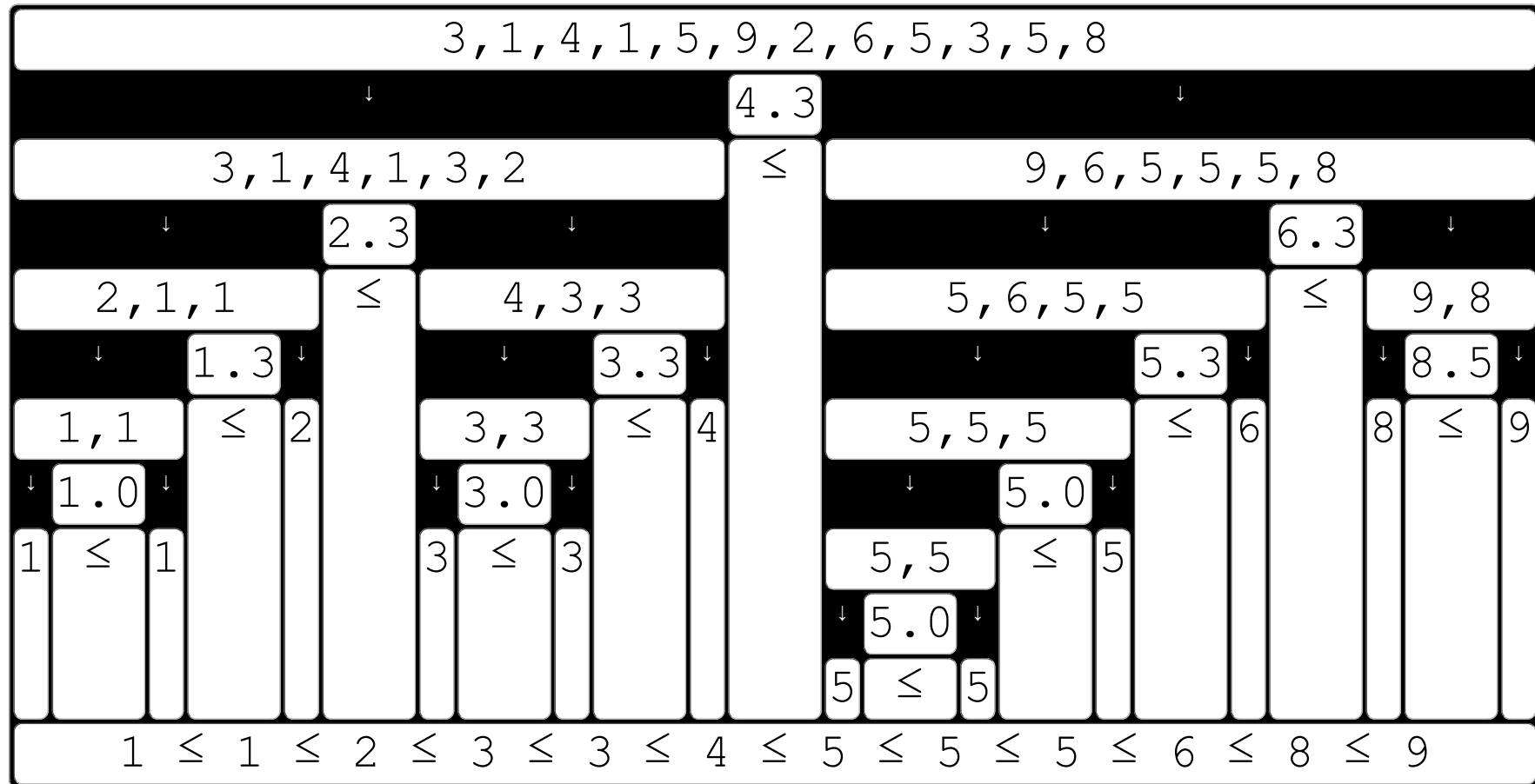


previous

first

When a list contains two or more values that are exactly the same, how do you compare those values?

next

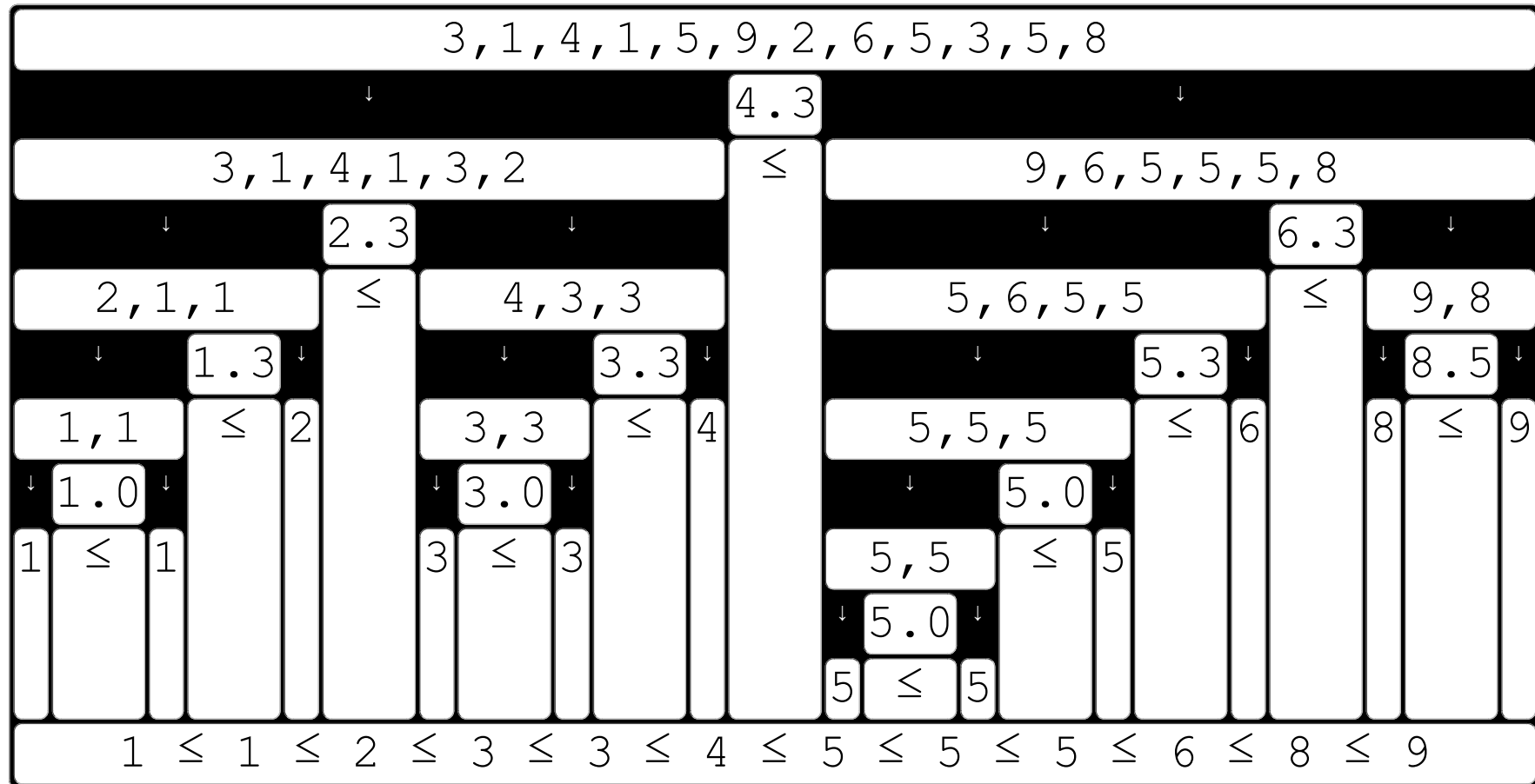


previous

first

In some applications, data with the same value may use additional information (what results in equal values not being the same).

next

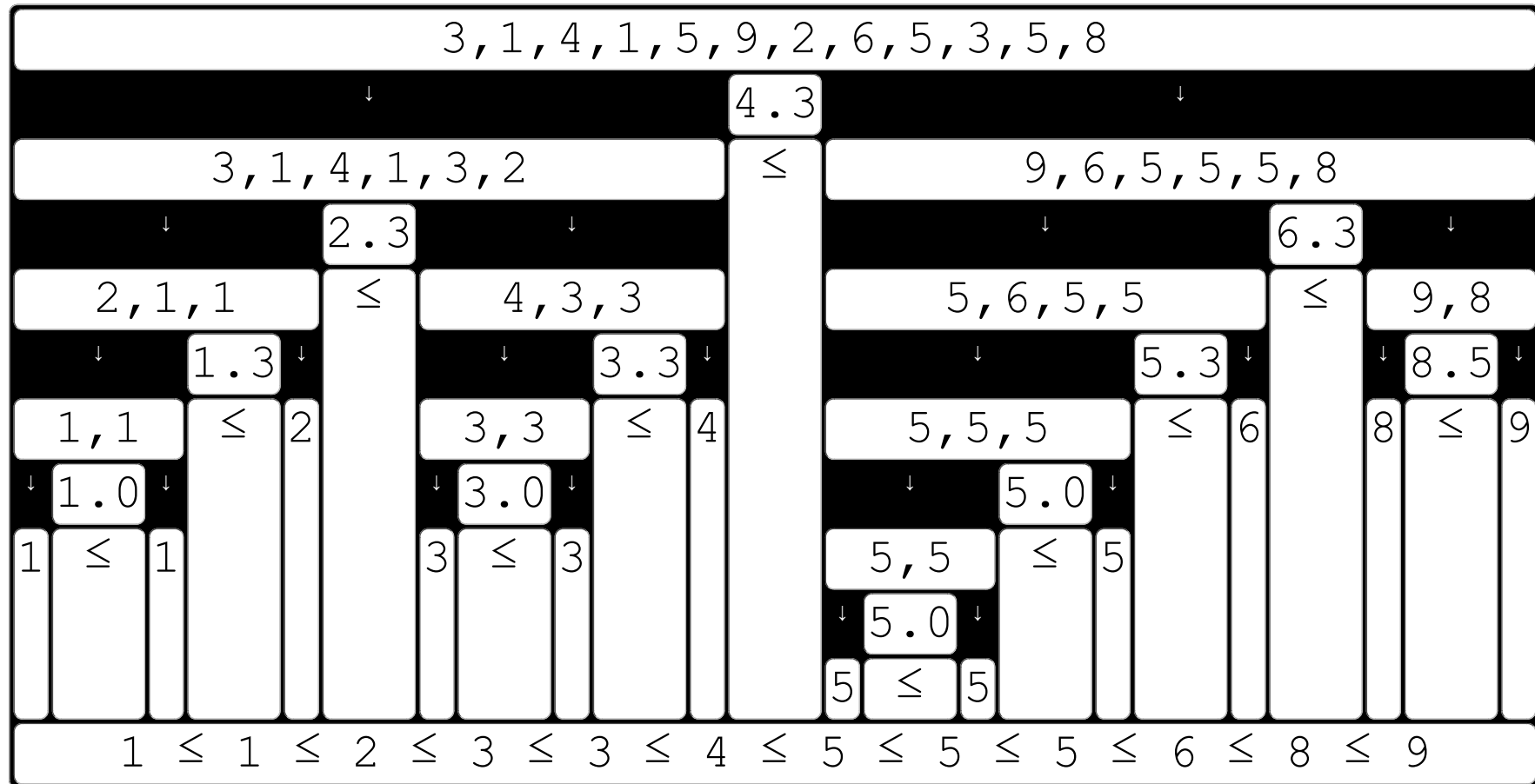


previous

So, a stable / unstable sort algorithm may be relevant, for some applications.

first

next

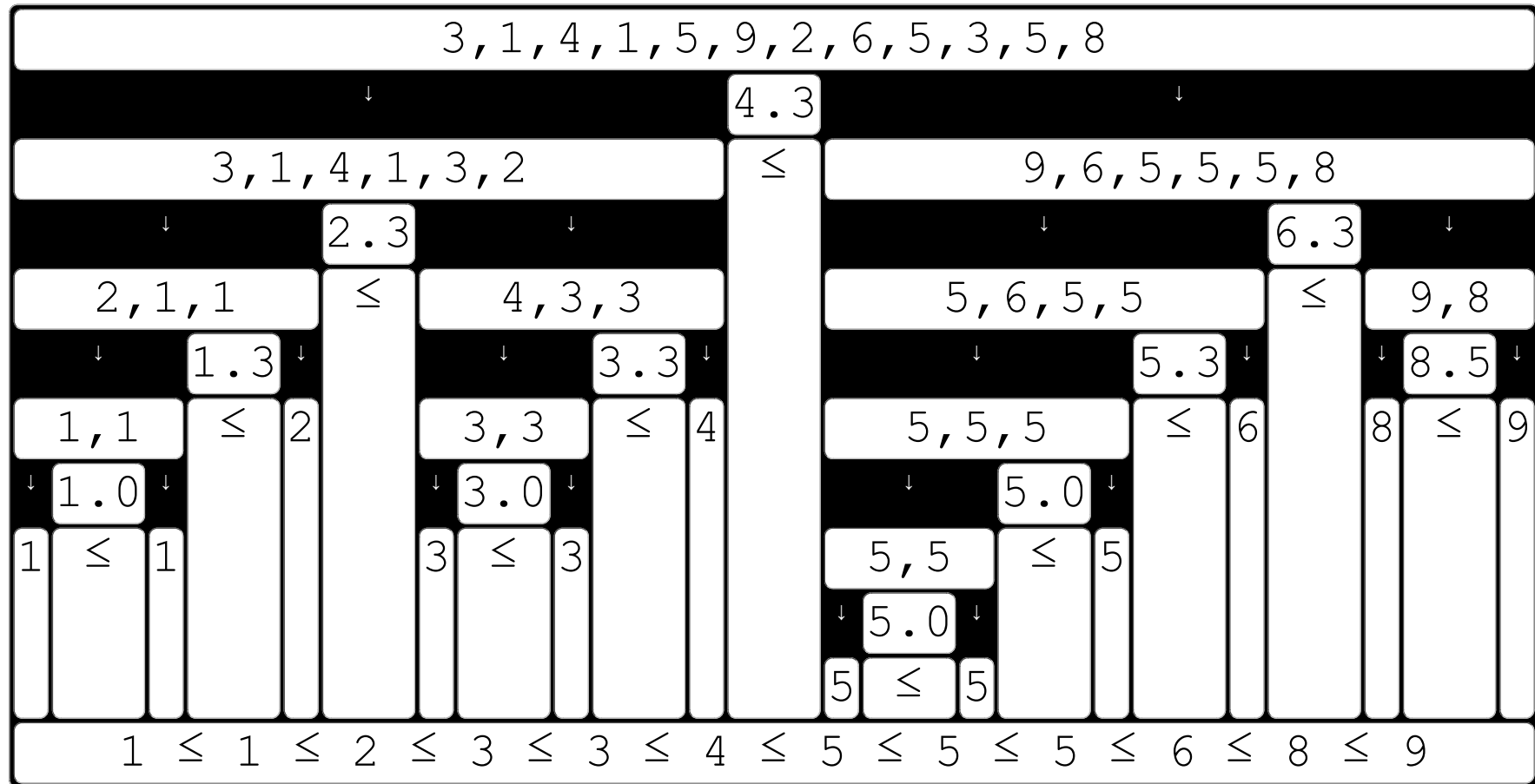


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Quicksort can be implemented as a stable sort algorithm, but with some additional storage space.

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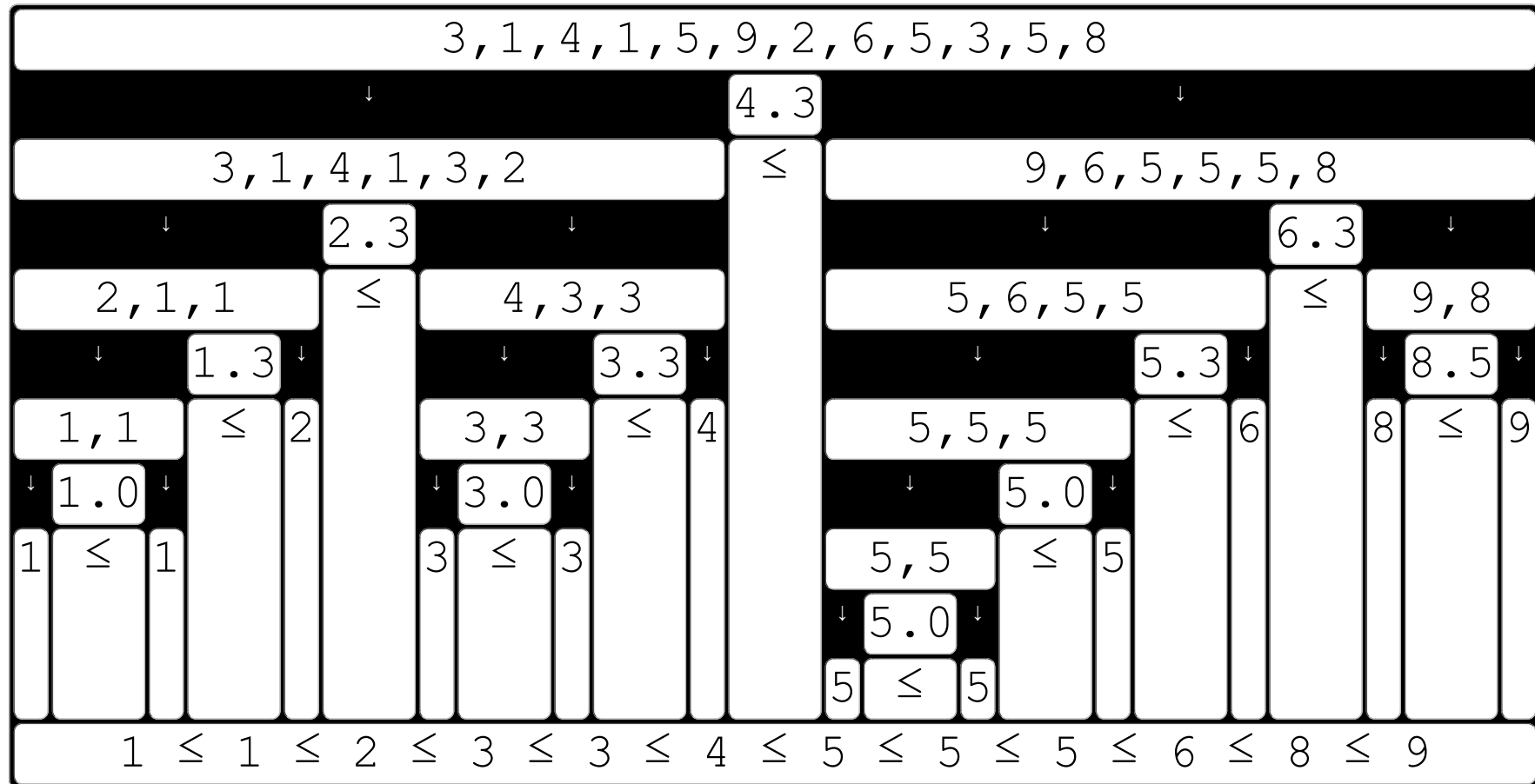


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first

(5) Advanced concepts: Median.

next

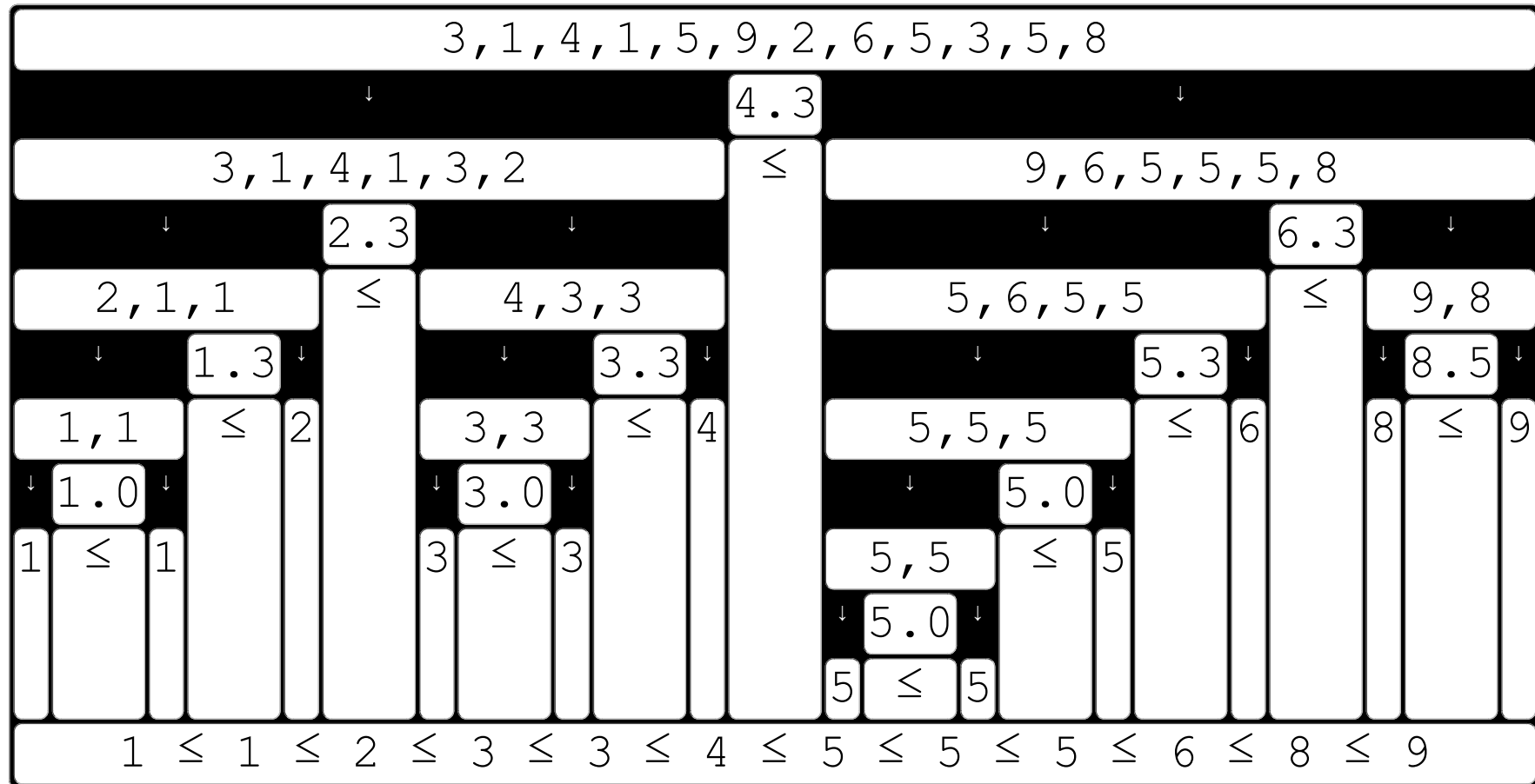


previous

This tutorial uses average as pivot. But the pivot could have another definition.

first

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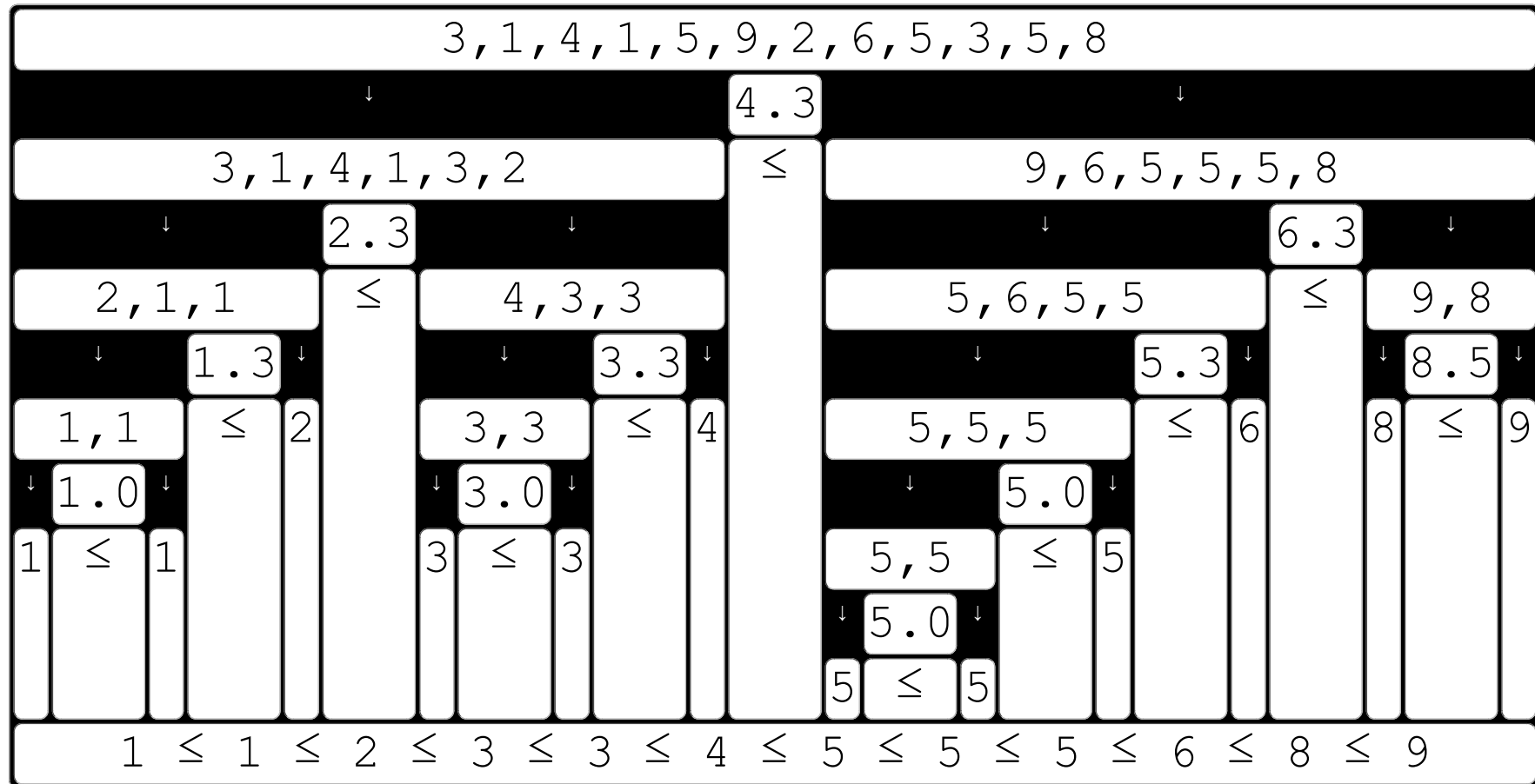


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The choice of pivot defines the size of each sublist. This can result in undesired unbalanced lists.

next

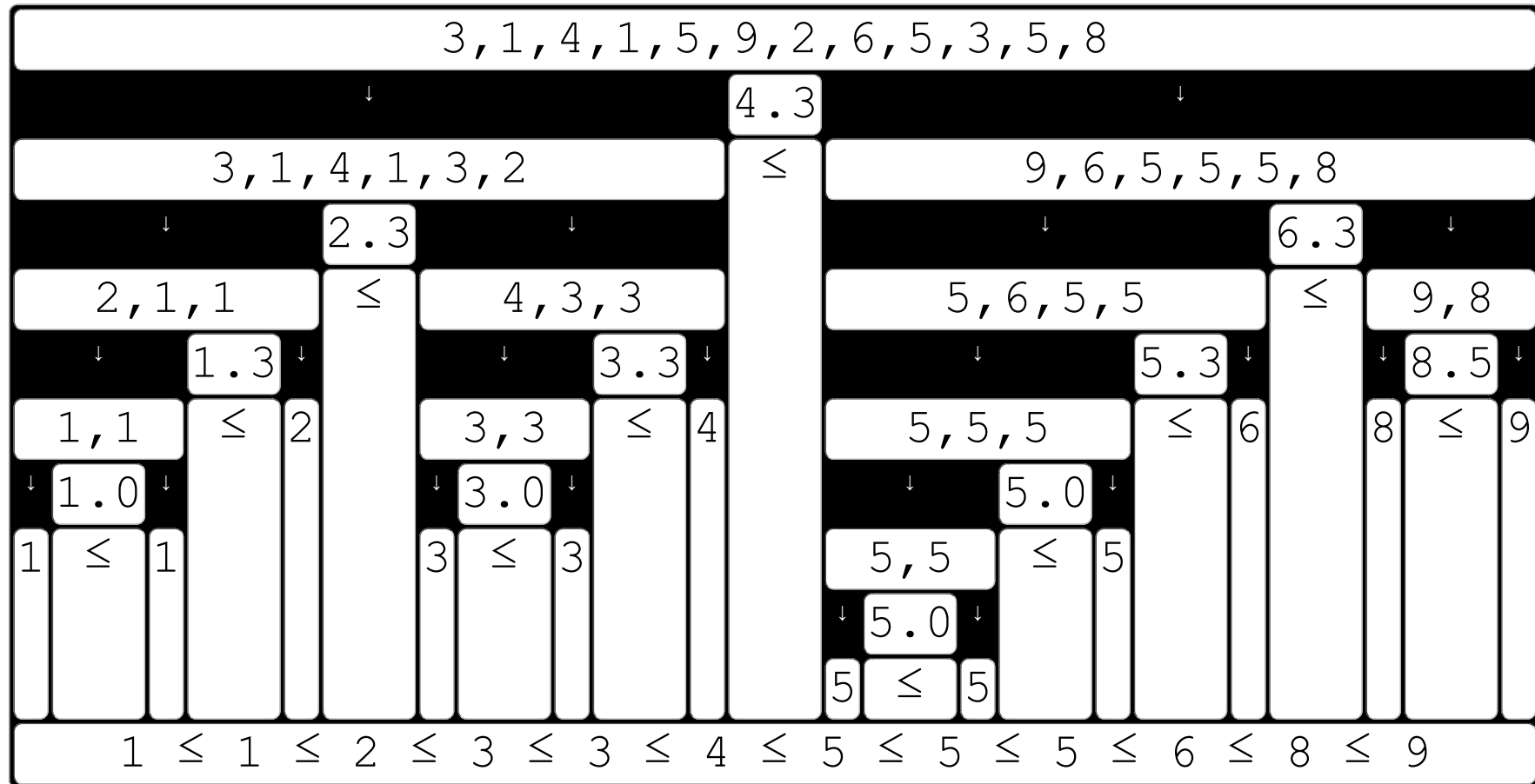


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first

The perfect pivot would be the median of a list.

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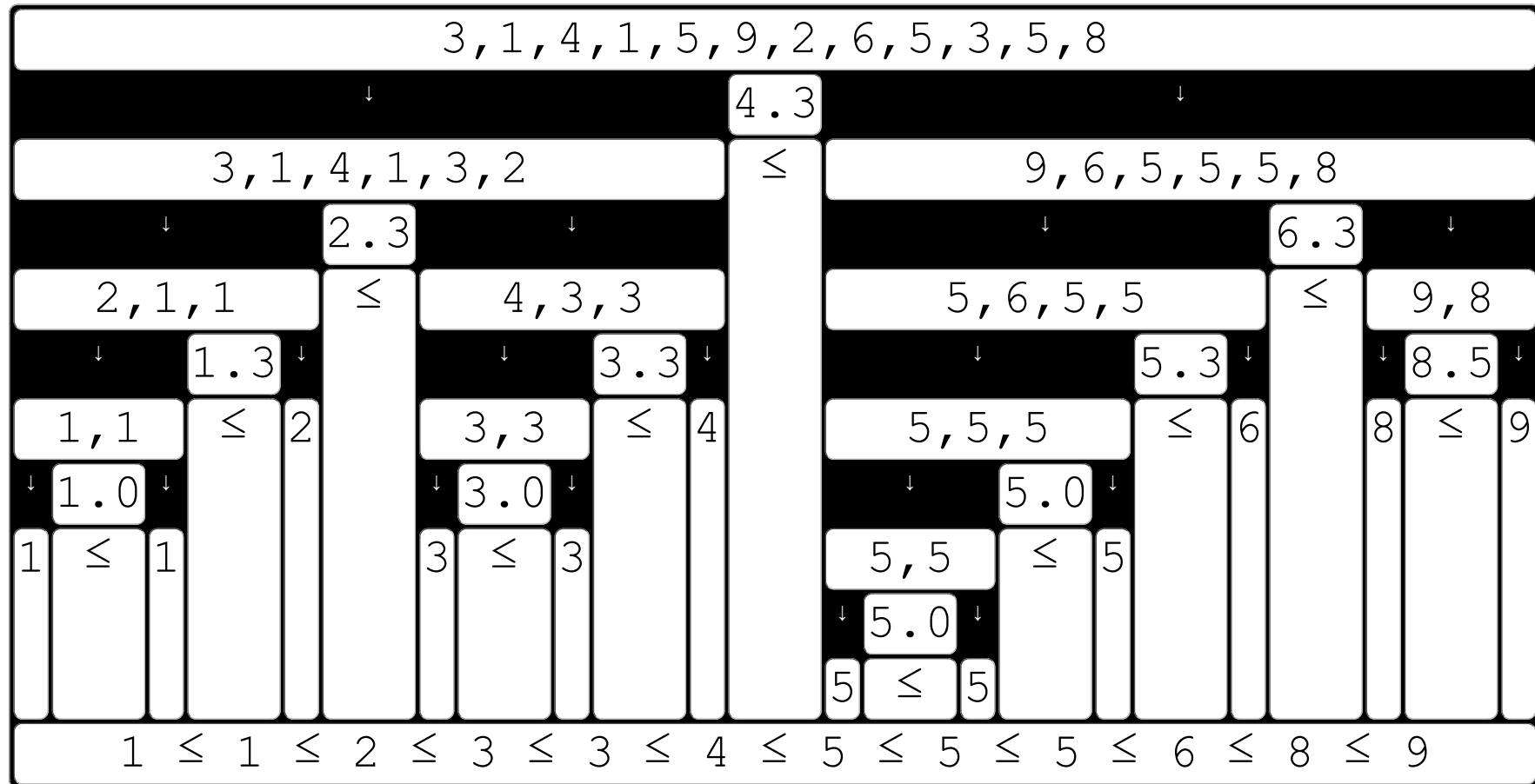


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first

But the effort to calculate the median would make Quicksort slower.

next

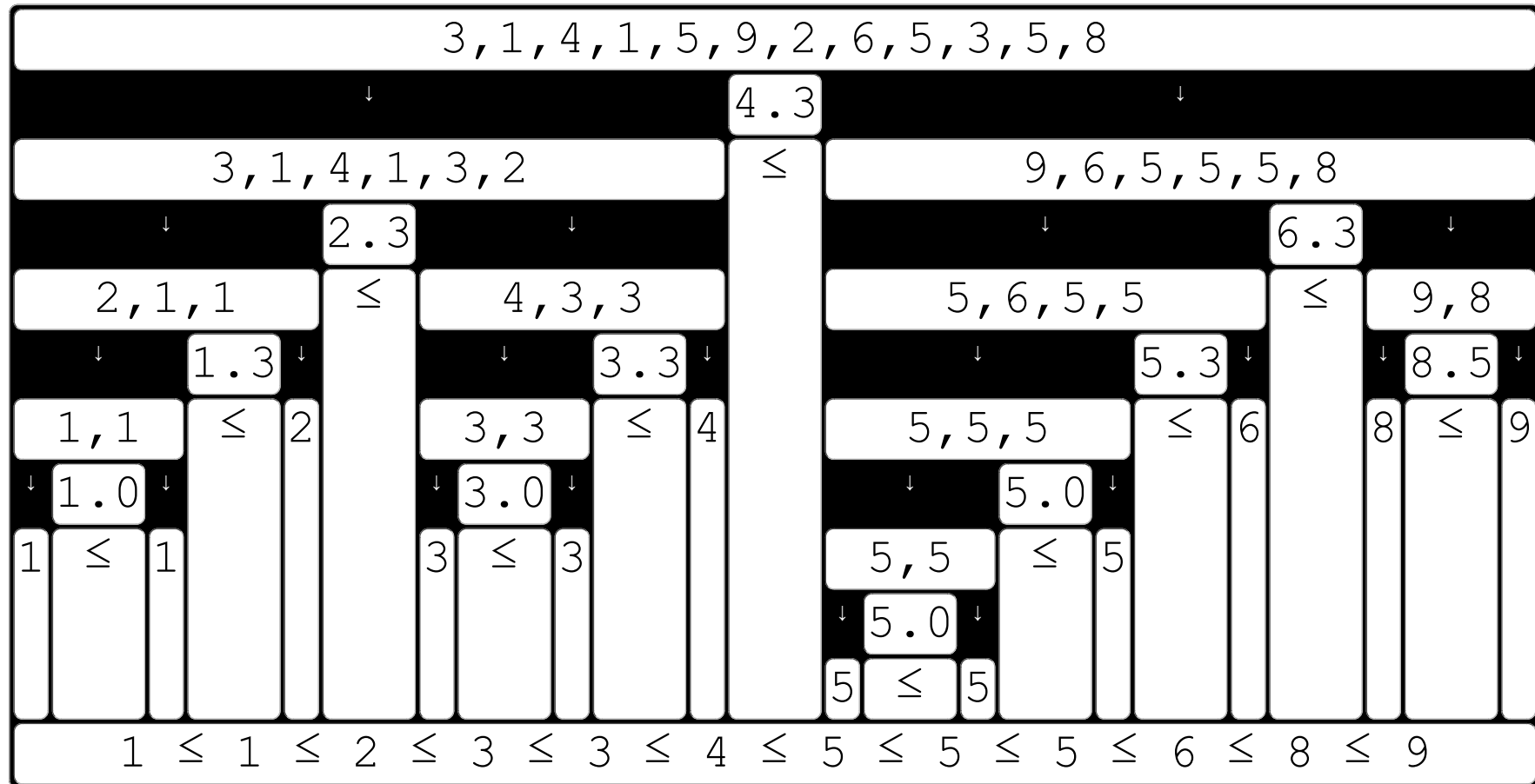


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first

So, in practice, some other value is defined as pivot: as near to the median as possible, and fast to calculate.

next

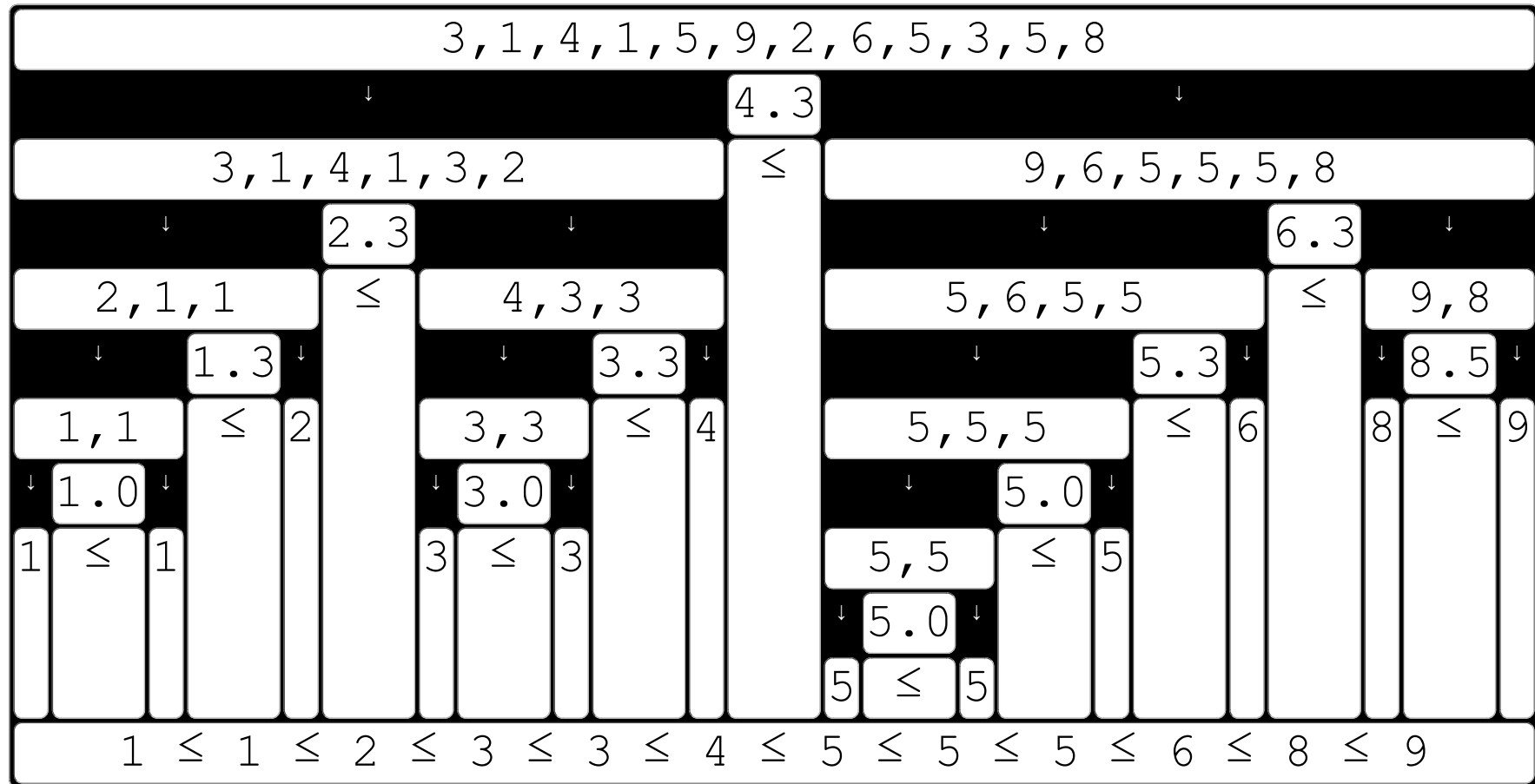


previous

Originally in 1959, [Tony Hoare](#) implemented Quicksort using the first (or last) element of a list as pivot. The algorithm usually works fine with this choice.

first

next



previous

first

The end.

Also, visit [Mergesort Interactive Tutorial](#).

next

Quicksort

Tutorial Highlight

<https://www.ime.usp.br/~otuyama/tutorial/sort/quicksort/Quicksort.html>