Dimension

# [6] Dimension

Key fact for this week: all bases for a vector space have the same size. We use this as the "basis" for answering many pending questions.

### Morphing Lemma

**Morphing Lemma:** Suppose S is a set of vectors, and B is a linearly independent set of vectors in Span S. Then  $|S| \ge |B|$ .

Before we prove it-what good is this lemma?

**Theorem:** Any basis for  $\mathcal{V}$  is a smallest generating set for  $\mathcal{V}$ .

**Proof:** Let S be a smallest generating set for  $\mathcal{V}$ . Let B be a basis for  $\mathcal{V}$ . Then B is a linearly independent set of vectors in Span S. By the Morphing Lemma, B is no bigger than S, so B is also a smallest generating set.

**Theorem:** All bases for a vector space  $\mathcal{V}$  have the same size.

**Proof:** They are all smallest generating sets.

**Morphing Lemma:** Suppose S is a set of vectors, and B is a linearly independent set of vectors in Span S. Then  $|S| \ge |B|$ .

Proof outline: modify S step by step, introducing vectors of B one by one, without increasing the size.

How? Using the Exchange Lemma....

**Exchange Lemma:** Suppose S is a set of vectors and A is a subset of S. Suppose z is a vector in Span S such that  $A \cup \{z\}$  is linearly independent. Then there is a vector  $\mathbf{w} \in S - A$  such that

Span 
$$S =$$
Span  $(S \cup \{z\} - \{w\})$ 

#### Proof of the Morphing Lemma

Let  $B = {\mathbf{b}_1, \dots, \mathbf{b}_n}$ . Define  $S_0 = S$ . Prove by induction on  $k \le n$  that there is a generating set  $S_k$  of Span S that contains  $\mathbf{b}_1, \dots, \mathbf{b}_k$  and has size |S|.

Base case: k = 0 is trivial.

To go from  $S_{k-1}$  to  $S_k$ : use the Exchange Lemma.

• 
$$A_k = \{ {f b}_1, \dots, {f b}_{k-1} \}$$
 and  ${f z} = {f b}_k$ 

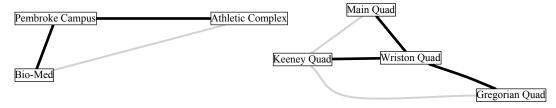
Exchange Lemma  $\Rightarrow$  there is a vector **w** in  $S_{k-1}$  such that

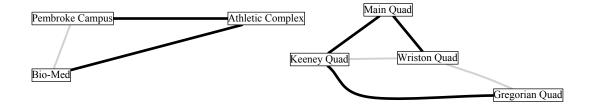
$$\mathsf{Span}\;(S_{k-1}\cup\{\mathbf{b}_k\}-\{\mathbf{w}\})=\mathsf{Span}\;S_{k-1}$$

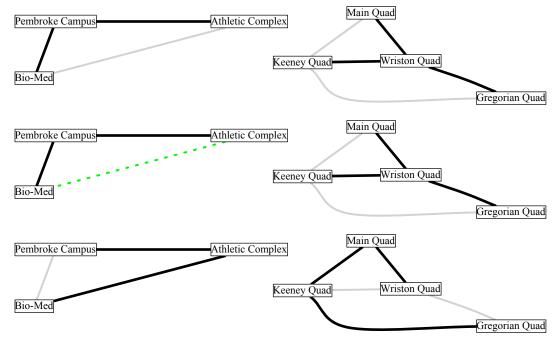
Set  $S_k = S_{k-1} \cup {\mathbf{b}_k} - {\mathbf{w}}.$ 

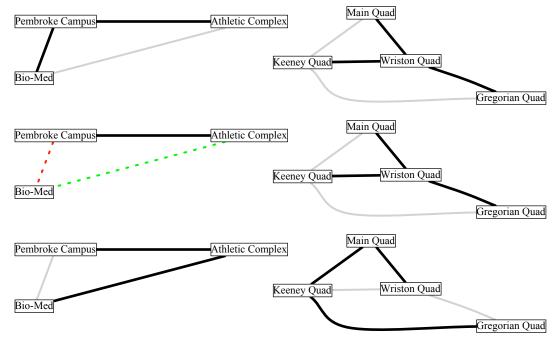
QED

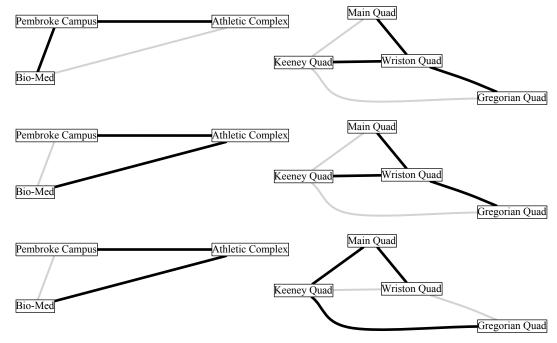
This induction proof is an algorithm.

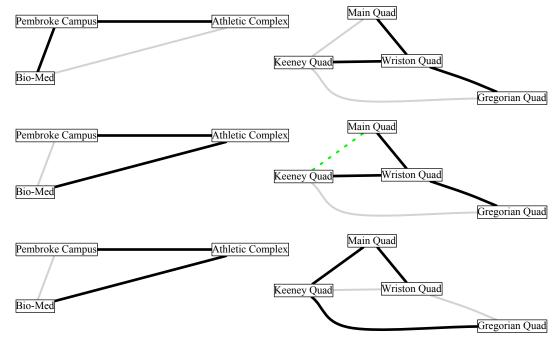


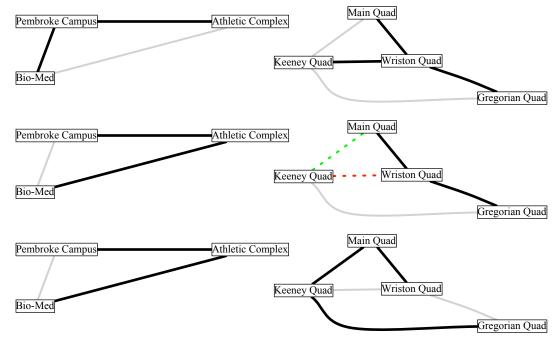


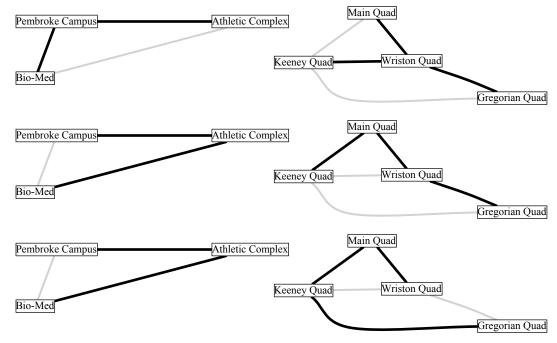


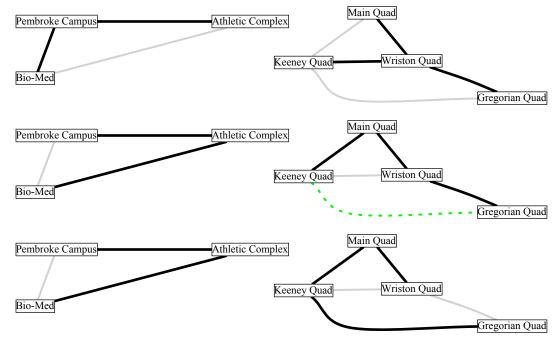


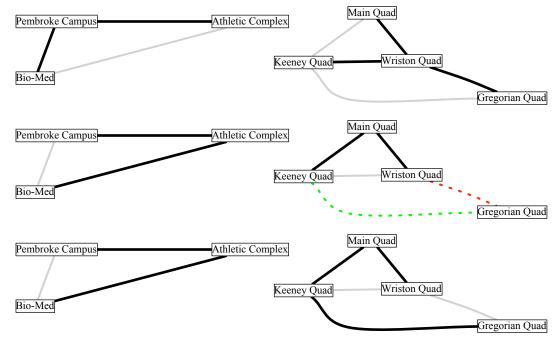


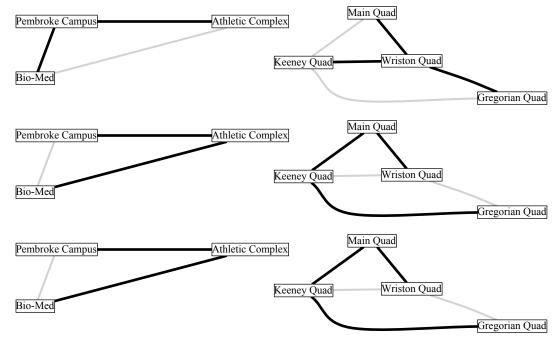












#### Dimension

**Definition:** We define the *dimension* of a vector space to be the size of a basis for that vector space. The dimension of a vector space  $\mathcal{V}$  is written dim  $\mathcal{V}$ .

**Definition:** We define the *rank* of a set S of vectors as the dimension of Span S. We write rank S.

**Example:** The vectors [1, 0, 0], [0, 2, 0], [2, 4, 0] are linearly dependent. Therefore their rank is less than three.

First two of these vectors form a basis for the span of all three, so the rank is two.

**Example:** The vector space Span  $\{[0,0,0]\}$  is spanned by an empty set of vectors. Therefore the rank of  $\{[0,0,0]\}$  is zero.

#### Row rank, column rank

**Definition:** For a matrix M, the row rank of M is the rank of its rows, and the column rank of M is the rank of its columns.

Equivalently, the row rank of M is the dimension of Row M, and the column rank of M is the dimension of Col M.

**Example:** Consider the matrix

$$M = \left[ \begin{array}{rrrr} 1 & 0 & 0 \\ 0 & 2 & 0 \\ 2 & 4 & 0 \end{array} \right]$$

whose rows are the vectors we saw before: [1, 0, 0], [0, 2, 0], [2, 4, 0]

The set of these vectors has rank two, so the row rank of M is two.

The columns of M are [1, 0, 2], [0, 2, 4], and [0, 0, 0].

Since the third vector is the zero vector, it is not needed for spanning the column space. Since each of the first two vectors has a nonzero where the other has a zero, these two are linearly independent, so the column rank is two.

#### Row rank, column rank

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**Example:** Consider the matrix

$$M = \left[ \begin{array}{rrrr} 1 & 0 & 0 & 5 \\ 0 & 2 & 0 & 7 \\ 0 & 0 & 3 & 9 \end{array} \right]$$

Each of the rows has a nonzero where the others have zeroes, so the three rows are linearly independent. Thus the row rank of M is three.

The columns of M are [1, 0, 0], [0, 2, 0], [0, 0, 3], and [5, 7, 9].

The first three columns are linearly independent, and the fourth can be written as a linear combination of the first three, so the column rank is three.

#### Row rank, column rank

**Definition:** For a matrix M, the row rank of M is the rank of its rows, and the column rank of M is the rank of its columns.

Equivalently, the row rank of M is the dimension of Row M, and the column rank of M is the dimension of Col M.

Does column rank always equal row rank? ©

## Geometry

We have asked:

**Fundamental Question:** How can we predict the dimensionality of the span of some vectors?



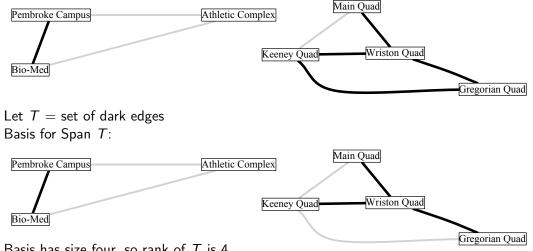
Now we can answer:

Compute the rank of the set of vectors.

#### **Examples:**

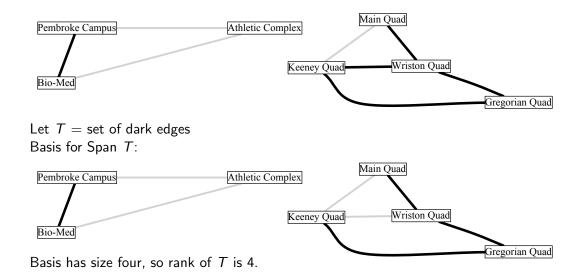
- Span  $\{[1,2,-2]\}$  is a line but Span  $\{[0,0,0]\}$  is a point. First vector space has dimension one, second has dimension zero.
- Span  $\{[1,2],[3,4]\}$  consists of all of  $\mathbb{R}^2$  but Span  $\{[1,3],[2,6]\}$  is a line The first has dimension two and the second has dimension one.
- Span  $\{[1,0,0], [0,1,0], [0,0,1]\}$  is  $\mathbb{R}^3$  but Span  $\{[1,0,0], [0,1,0], [1,1,0]\}$  is a plane. The first has dimension three and the second has dimension two.

### Dimension and rank in graphs



Basis has size four, so rank of T is 4.

### Dimension and rank in graphs



## Cardinality of a vector space over GF(2)

Recall checksum problem

Checksum function  $\textbf{x} \mapsto [\textbf{a}_1 \cdot \textbf{x}, \dots, \textbf{a}_{64} \cdot \textbf{x}]$ 

Original "file"  $\boldsymbol{p},$  transmission error  $\boldsymbol{e}$  so corrupted file is  $\boldsymbol{p}+\boldsymbol{e}.$ 

What is probability that corrupted file has the same checksum as original?

If error is chosen according to uniform distribution,

Probability  $(\mathbf{p} + \mathbf{e} \text{ has same checksum as } \mathbf{p})$ 

Probability (e is a solution to homogeneous linear system)
 number of solutions to homogeneous linear system

number of *n*-vectors number of solutions to homogeneous linear system

2<sup>n</sup>

raising Question

How to find number of solutions to a homogeneous linear system over GF(2)?

## Cardinality of a vector space over GF(2)

How to find number of solutions to a homogeneous linear system over GF(2)?

Solution set of a homogeneous linear system is a vector space. Question becomes

How to find out cardinality of a vector space  $\mathcal{V}$  over GF(2)?

- Suppose basis for  $\mathcal{V}$  is  $\mathbf{b}_1, \ldots, \mathbf{b}_n$ .
- Then  $\mathcal{V}$  is set of linear combinations

$$\beta_1 \mathbf{b}_1 + \cdots + \beta_n \mathbf{b}_n$$

- ▶ Number of linear combinations is 2<sup>n</sup>.
- By Unique-Representation Lemma, every linear combination gives a different vector of V.
- Thus cardinality is  $2^{\dim \mathcal{V}}$ .

## Cardinality of a vector space over GF(2)

Cardinality of a vector space  $\mathcal{V}$  over GF(2) is  $2^{\dim \mathcal{V}}$ .

How to find dimension of solution set of a homogeneous linear system?

Write linear system as  $A\mathbf{x} = \mathbf{0}$ .

How to find dimension of the null space of A?

Answers will come later.

#### Subset-Basis Lemma

**Lemma:** Every finite set T of vectors contains a subset S that is a basis for Span T. **Proof:** The Grow algorithm finds a basis for  $\mathcal{V}$  if it terminates.

Initialize  $S = \emptyset$ .

Repeat while possible: select a vector  $\mathbf{v}$  in  $\mathcal{V}$  that is not in Span S, and put it in S.

Revised version:

Initialize  $S = \emptyset$ Repeat while possible: select a vector **v** in *T* that is not in Span *S*, and put it in *S*.

Differs from original:

• This algorithm stops when Span S contains every vector in T.

► The original Grow algorithm stops only once Span S contains every vector in  $\mathcal{V}$ . However, that's okay: when Span S contains all the vectors in T, Span S also contains all linear combinations of vectors in T, so at this point Span  $S = \mathcal{V}$ . Shows that original Grow algorithm can be guided to make same choices as this algorithm, so result is a basis. QED

## Termination of Grow algorithm

 $\begin{array}{l} {\rm def \ GROW}(\mathcal{V}) \\ B=\emptyset \\ {\rm repeat \ while \ possible:} \\ {\rm find \ a \ vector \ V \ in \ } \mathcal{V} \ that \ is \ not \ in \ Span \ B, \ and \ put \ it \ in \ S. \end{array}$ 

**Grow-Algorithm-Termination Lemma:** If  $\mathcal{V}$  is a subspace of  $\mathbb{F}^D$  where D is finite then  $\operatorname{GROW}(\mathcal{V})$  terminates. **Proof:** By Grow-Algorithm Corollary, B is linearly independent throughout. Apply the Morphing Lemma with  $S = \{$ standard generators for  $\mathbb{F}^D \} \Rightarrow |B| \leq |S| = |D|$ .

QED

Since *B* grows in each iteration, there are at most |D| iterations.

## Every subspace of $\mathbb{F}^D$ contains a basis

**Grow-Algorithm-Termination Lemma:** If  $\mathcal{V}$  is a subspace of  $\mathbb{F}^D$  where D is finite then  $\operatorname{GROW}(\mathcal{V})$  terminates. **Theorem:** For finite D, every subspace of  $\mathbb{F}^D$  contains a basis. **Proof:** Let  $\mathcal{V}$  be a subspace of  $\mathbb{F}^D$ .

```
def GROW(\mathcal{V})

B = \emptyset

repeat while possible:

find a vector v in \mathcal{V} that is not in Span B, and put it in B.
```

Grow-Algorithm-Termination Lemma ensures algorithm terminates. Upon termination, every vector in  $\mathcal{V}$  is in Span B, so B is a set of generators for  $\mathcal{V}$ . By Grow-Algorithm Corollary, B is linearly independent. Therefore B is a basis for  $\mathcal{V}$ . QED

## Superset-Basis Lemma

**Grow-Algorithm-Termination Lemma:** If  $\mathcal{V}$  is a subspace of  $\mathbb{F}^D$  where D is finite then  $\operatorname{GROW}(\mathcal{V})$  terminates.

**Superset-Basis Lemma:** Let  $\mathcal{V}$  be a vector space consisting of *D*-vectors where *D* is finite. Let *C* be a linearly independent set of vectors belonging to  $\mathcal{V}$ . Then  $\mathcal{V}$  has a basis *B* containing all vectors in *C*.

**Proof:** Use version of Grow algorithm:

Initialize *B* to the empty set. Repeat while possible: select a vector  $\mathbf{v}$  in  $\mathcal{V}$  (preferably in *C*) that is not in Span *B*, and put it in *B*.

At first, *B* will consist of vectors in *C* until *B* contains all of *C*. Then more vectors will be added to *B* until Span  $B = \mathcal{V}$  By Grow-Algorithm Corollary, *B* is linearly independent throughout. Therefore, once algorithm terminates, *B* contains *C* and is a basis for  $\mathcal{U}$ .

Termination is implied by Grow Algorithm Termination Lemma.



#### Estimating dimension

 $\mathcal{T} = \{ [-0.6, -2.1, -3.5, -2.2], [-1.3, 1.5, -0.9, -0.5], [4.9, -3.7, 0.5, -0.3], [2.6, -3.5, -1.2, -2.0], [-1.5, -2.5, -3.5, 0.94] \}.$  What is the rank of  $\mathcal{T}$ ?

By Subset-Basis Lemma, T contains a basis.

```
Therefore dim Span T \leq |T|.
```

Therefore rank  $T \leq |T|$ .

**Proposition:** A set T of vectors has rank  $\leq |T|$ .

#### **Dimension Lemma**

**Dimension Lemma:** If  $\mathcal{U}$  is a subspace of  $\mathcal{W}$  then

- ▶ **D1:** dim  $\mathcal{U} \leq \dim \mathcal{W}$ , and
- ▶ **D2:** if dim U = dim W then U = W

**Proof:** Let  $\mathbf{u}_1, \ldots, \mathbf{u}_k$  be a basis for  $\mathcal{U}$ .

By Superset-Basis Lemma, there is a basis B for  $\mathcal{W}$  that contains  $\mathbf{u}_1, \ldots, \mathbf{u}_k$ .

- $\triangleright B = {\mathbf{u}_1, \dots, \mathbf{u}_k, \mathbf{b}_1, \dots, \mathbf{b}_r}$
- Thus  $k \leq |B|$ , and

• If 
$$k = |B|$$
 then  $\{\mathbf{u}_1, \ldots, \mathbf{u}_k\} = B$ 

**Example:** Suppose  $\mathcal{V} = \text{Span} \{[1, 2], [2, 1]\}$ . Clearly  $\mathcal{V}$  is a subspace of  $\mathbb{R}^2$ . However, the set  $\{[1, 2], [2, 1]\}$  is linearly independent, so dim  $\mathcal{V} = 2$ . Since dim  $\mathbb{R}^2 = 2$ , D2 shows that  $\mathcal{V} = \mathbb{R}^2$ . **Example:**  $S = \{[-0.6, -2.1, -3.5, -2.2], [-1.3, 1.5, -0.9, -0.5], [4.9, -3.7, 0.5, -0.3], [2.6, -3.5, -1.2, -2.0], [-1.5, -2.5, -3.5, 0.94]\}$ Since every vector in S is a 4-vector, Span S is a subspace of  $\mathbb{R}^4$ . Since dim  $\mathbb{R}^4 = 4$ , D1 shows dim Span  $S \leq 4$ .

QED

**Proposition:** Any set of *D*-vectors has rank at most |D|.

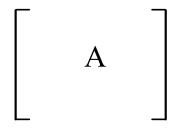
**Rank Theorem:** For every matrix M, row rank equals column rank.

**Lemma:** For any matrix A, row rank of  $A \leq$  column rank of A To show theorem:

- ▶ Apply lemma to  $M \Rightarrow$  row rank of  $M \le$  column rank of M
- Apply lemma to M<sup>T</sup> ⇒ row rank of M<sup>T</sup> ≤ column rank of M<sup>T</sup> ⇒ column rank of M ≤ row rank of M

Combine  $\Rightarrow$  row rank of M = column rank of M

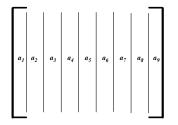
Proof of lemma: For any matrix A, row rank of  $A \leq$  column rank of A



Think of A as columns  $\mathbf{a}_1, \dots, \mathbf{a}_n$ . Let  $\mathbf{b}_1, \dots, \mathbf{b}_r$  be basis for column space (so column rank = r). Write each column of A in terms of basis:  $\begin{bmatrix} \mathbf{a}_j \end{bmatrix} = \begin{bmatrix} \mathbf{b}_1 & \cdots & \mathbf{b}_r \end{bmatrix} \begin{bmatrix} \mathbf{u}_j \end{bmatrix}$ Use matrix-vector definition of matrix-matrix multiplication to rewrite as matrix-matrix equation A = BU.

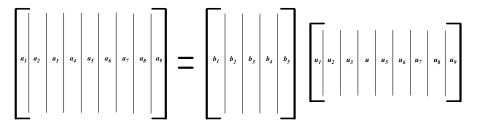
B has r columns and U has r rows.

Write A and B in terms of rows: row i of A equals row i of B times U. Write U in terms of rows: row i of A is a linear combination of rows of U. Each row of A is in span of the r rows of U. **Thus row rank of** A **is at most** r. Proof of lemma: For any matrix A, row rank of  $A \leq$  column rank of A



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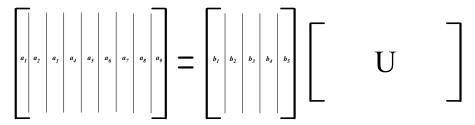


Think of *A* as columns  $\mathbf{a}_1, \dots, \mathbf{a}_n$ . Let  $\mathbf{b}_1, \dots, \mathbf{b}_r$  be basis for column space (so column rank = *r*). Write each column of *A* in terms of basis:  $\begin{bmatrix} \mathbf{a}_j \end{bmatrix} = \begin{bmatrix} \mathbf{b}_1 & \cdots & \mathbf{b}_r \end{bmatrix} \begin{bmatrix} \mathbf{u}_j \end{bmatrix}$ 

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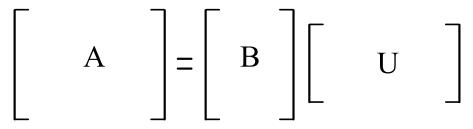
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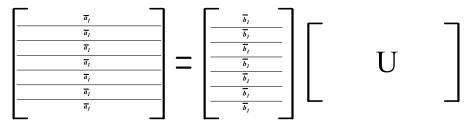
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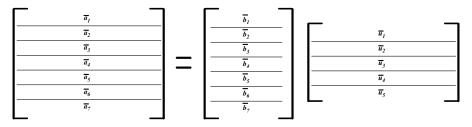
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# Simple authentication revisited

- Password is an *n*-vector  $\hat{\mathbf{x}}$  over GF(2)
- **Challenge:** Computer sends random *n*-vector **a**
- Response: Human sends back  $\mathbf{a} \cdot \hat{\mathbf{x}}$ . Repeated until Computer is convinced that Human knows password  $\hat{\mathbf{x}}$ .

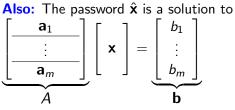
Eve eavesdrops on communication, learns m pairs

$$\mathbf{a_1}, b_1$$
  
 $\vdots$   
 $\mathbf{a_m}, b_m$   
such that  $b_i$  is right response to challenge  
 $\mathbf{a}_i$ 

Then Eve can calculate right response to any challenge in Span  $\{a_1, \ldots, a_m\}$ :

Suppose  $\mathbf{a} = \alpha_1 \mathbf{a}_1 + \dots + \alpha_m \mathbf{a}_m$ Then right response is  $\alpha_1 b_1 + \dots + \alpha_m b_m$  **Fact:** Probably rank  $[\mathbf{a}_1, \ldots, \mathbf{a}_m]$  is not much less than min $\{m, n\}$ .

Once m > n, probably Span  $\{\mathbf{a}_1, \ldots, \mathbf{a}_m\}$ is all of  $GF(2)^n$ so Eve can respond to any challenge.



Solution set of  $A\mathbf{x} = \mathbf{b}$  is  $\hat{\mathbf{x}} + \text{Null } A$ 

Once rank A reaches n, columns of A are linearly independent so Null A is trivial, so only solution is the password  $\hat{\mathbf{x}}$ , so Eve can compute the password using solver.

## Direct Sum

Let  $\mathcal{U}$  and  $\mathcal{V}$  be two vector spaces consisting of D-vectors over a field  $\mathbb{F}$ .

**Definition:** If  $\mathcal{U}$  and  $\mathcal{V}$  share only the zero vector then we define the *direct sum* of  $\mathcal{U}$  and  $\mathcal{V}$  to be the set

 $\{\mathbf{u} + \mathbf{v} : \mathbf{u} \in \mathcal{U}, \mathbf{v} \in \mathcal{V}\}$ 

written  $\mathcal{U}\oplus\mathcal{V}$ 

That is,  $\mathcal{U} \oplus \mathcal{V}$  is the set of all sums of a vector in  $\mathcal{U}$  and a vector in  $\mathcal{V}$ .

In Python, [u+v for u in U for v in V]

(But generally  $\mathcal{U}$  and  $\mathcal{V}$  are infinite so the Python is just suggestive.)

## Direct Sum: Example

Vectors over GF(2):

**Example:** Let  $\mathcal{U} = \text{Span} \{1000, 0100\}$  and let  $\mathcal{V} = \text{Span} \{0010\}$ .

- Every nonzero vector in U has a one in the first or second position (or both) and nowhere else.
- $\blacktriangleright$  Every nonzero vector in  ${\cal V}$  has a one in the third position and nowhere else.

Therefore the only vector in both  $\mathcal{U}$  and  $\mathcal{V}$  is the zero vector.

Therefore  $\mathcal{U} \oplus \mathcal{V}$  is defined.

 $\mathcal{U} \oplus \mathcal{V} = \{0000 + 0000, 1000 + 0000, 0100 + 0000, 1100 + 0000, 0000 + 0010, 1000 + 0010, 1000 + 0010, 1100 + 0010\}$ 

which is equal to  $\{0000, 1000, 0100, 1100, 0010, 1010, 0110, 1110\}$ .

## Direct Sum: Example

Vectors over  $\mathbb{R}$ :

**Example:** Let  $\mathcal{U} = \text{Span} \{ [1, 2, 1, 2], [3, 0, 0, 4] \}$  and let  $\mathcal{V}$  be the null space of  $\begin{bmatrix} 0 & 1 & -1 & 0 \\ 1 & 0 & 0 & -1 \end{bmatrix}$ .

- ▶ The vector [2, -2, -1, 2] is in  $\mathcal{U}$  because it is [3, 0, 0, 4] [1, 2, 1, 2]
- It is also in  $\mathcal{V}$  because

$$\left[\begin{array}{rrrr} 0 & 1 & -1 & 0 \\ 1 & 0 & 0 & -1 \end{array}\right] \left[\begin{array}{r} 2 \\ -2 \\ -1 \\ 2 \end{array}\right] = \left[\begin{array}{r} 0 \\ 0 \end{array}\right]$$

Therefore we cannot form  $\mathcal{U} \oplus \mathcal{V}$ .

## Direct Sum: Example

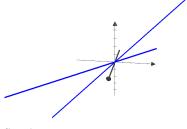
Vectors over  $\mathbb{R}:$ 

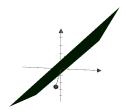
**Example:** 

- ▶ Let  $U = \text{Span} \{ [4, -1, 1] \}.$
- Let  $\mathcal{V} = \text{Span } \{[0, 1, 1]\}.$

The only intersection is at the origin, so  $\mathcal{U}\oplus\mathcal{V}$  is defined.

- $\mathcal{U} \oplus \mathcal{V}$  is the set of vectors  $\mathbf{u} + \mathbf{v}$ where  $\mathbf{u} \in \mathcal{U}$  and  $\mathbf{v} \in \mathcal{V}$ .
- This is just Span  $\{[4, -1, 1], [0, 1, 1]\}$
- Plane containing the two lines





## Properties of direct sum

**Lemma:**  $\mathcal{U} \oplus \mathcal{V}$  is a vector space.

(Prove using Properties V1, V2, V3.)

Lemma: The union of

- $\blacktriangleright$  a set of generators of  $\mathcal U,$  and
- $\blacktriangleright$  a set of generators of  ${\cal V}$
- is a set of generators for  $\mathcal{U}\oplus\mathcal{V}.$

**Proof:** Suppose  $\mathcal{U} = \text{Span} \{\mathbf{u}_1, \dots, \mathbf{u}_m\}$  and  $\mathcal{V} = \text{Span} \{\mathbf{v}_1, \dots, \mathbf{v}_n\}$ . Then

- every vector in  $\mathcal{U}$  can be written as  $\alpha_1 \mathbf{u}_1 + \cdots + \alpha_m \mathbf{u}_m$ , and
- every vector in  $\mathcal{V}$  can be written as  $\beta_1 \mathbf{v}_1 + \cdots + \beta_n \mathbf{v}_n$

so every vector in  $\mathcal{U}\oplus\mathcal{V}$  can be written as

$$\alpha_1 \mathbf{u}_1 + \dots + \alpha_m \mathbf{u}_m + \beta_1 \mathbf{v}_1 + \dots + \beta_n \mathbf{v}_n$$

## Properties of direct sum

**Direct Sum Basis Lemma:** 

Union of a basis of  $\mathcal{U}$  and a basis of  $\mathcal{V}$  is a basis of  $\mathcal{U} \oplus \mathcal{V}$ .

**Proof:** Clearly

- $\blacktriangleright$  a basis of  ${\cal U}$  is a set of generators for  ${\cal U},$  and
- a basis of  $\mathcal{V}$  is a set of generators for  $\mathcal{V}$ .

Therefore the previous lemma shows that

• the union of a basis for  $\mathcal{U}$  and a basis for  $\mathcal{V}$  is a generating set for  $\mathcal{U} \oplus \mathcal{V}$ .

We just need to show that the union is linearly independent.

## Properties of direct sum

#### **Direct Sum Basis Lemma:**

Union of a basis of  $\mathcal{U}$  and a basis of  $\mathcal{V}$  is a basis of  $\mathcal{U} \oplus \mathcal{V}$ .

**Proof, cont'd:** Let  $\{\mathbf{u}_1, \ldots, \mathbf{u}_m\}$  be a basis for  $\mathcal{U}$ . Let  $\{\mathbf{v}_1, \ldots, \mathbf{v}_n\}$  be a basis for  $\mathcal{V}$ . We need to show that  $\{\mathbf{u}_1, \ldots, \mathbf{u}_m, \mathbf{v}_1, \ldots, \mathbf{v}_n\}$  is independent. Suppose

$$\mathbf{0} = \alpha_1 \, \mathbf{u}_1 + \cdots + \alpha_m \mathbf{u}_m + \beta_1 \, \mathbf{v}_1 + \cdots + \beta_n \, \mathbf{v}_n.$$

Then

$$\underbrace{\alpha_1 \ \mathbf{u}_1 + \dots + \alpha_m \ \mathbf{u}_m}_{\text{in } \mathcal{U}} = \underbrace{(-\beta_1) \ \mathbf{v}_1 + \dots + (-\beta_n) \ \mathbf{v}_n}_{\text{in } \mathcal{V}}$$

Left-hand side is a vector in  $\mathcal{U}$ , and right-hand side is a vector in  $\mathcal{V}$ .

By definition of  $\mathcal{U} \oplus \mathcal{V}$ , the only vector in both  $\mathcal{U}$  and  $\mathcal{V}$  is the zero vector. This shows:

$$\mathbf{0} = \alpha_1 \, \mathbf{u}_1 + \dots + \alpha_m \, \mathbf{u}_m$$

and

$$\mathbf{0} = (-\beta_1) \mathbf{v}_1 + \cdots + (-\beta_n) \mathbf{v}_n$$

By linear independence, the linear combinations must be trivial.



## Direct Sum

Direct Sum Basis Lemma:

Union of a basis of  $\mathcal{U}$  and a basis of  $\mathcal{V}$  is a basis of  $\mathcal{U} \oplus \mathcal{V}$ .

**Direct Sum Dimension Corollary:** dim  $\mathcal{U}$  + dim  $\mathcal{V}$  = dim  $\mathcal{U} \oplus \mathcal{V}$ 

**Proof:** A basis for  $\mathcal{U}$  together with a basis for  $\mathcal{V}$  forms a basis for  $\mathcal{U} \oplus \mathcal{V}$ . QED

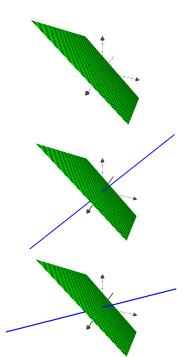
## Complementary subspace

If  $\mathcal{U} \oplus \mathcal{V} = \mathcal{W}$ , we say  $\mathcal{U}$  and  $\mathcal{V}$  are *complementary* subspaces of  $\mathcal{W}$ .

**Example:** Suppose  $\mathcal{U}$  is a plane in  $\mathbb{R}^3$ .

Then any line through the origin that does not lie in  ${\cal U}$  is complementary subspace with respect to  $\mathbb{R}^3$ 

Example illustrates that, for a given subspace  $\mathcal{U}$  of  $\mathcal{W}$ , there can be many different subspaces  $\mathcal{V}$  such that  $\mathcal{U}$  and  $\mathcal{V}$  are complementary.



## Complementary subspace

**Proposition:** For any finite-dimensional vector space  $\mathcal{W}$  and any subspace  $\mathcal{U}$ , there is a subspace  $\mathcal{V}$  such that  $\mathcal{U}$  and  $\mathcal{V}$  are complementary.

**Proof:** Let  $\mathbf{u}_1, \ldots, \mathbf{u}_k$  be a basis for  $\mathcal{U}$ . By Superset-Basis Lemma, there is a basis for  $\mathcal{W}$  that includes  $\mathbf{u}_1, \ldots, \mathbf{u}_k$ :

$$\mathsf{B} = \{\mathsf{u}_1, \ldots, \mathsf{u}_k, \mathsf{v}_1, \ldots, \mathsf{v}_r\}$$

Let  $\mathcal{V} = \text{Span } \{\mathbf{v}_1, \dots, \mathbf{v}_r\}.$ 

Any vector in  $\mathcal{W}$  can be written in terms of its basis:

$$\mathbf{w} = \underbrace{\alpha_1 \, \mathbf{u}_1 + \dots + \alpha_k \, \mathbf{u}_k}_{\text{in } \mathcal{U}} + \underbrace{\beta_1 \, \mathbf{v}_1 + \dots + \beta_r \, \mathbf{v}_r}_{\text{in } \mathcal{V}}$$

If some vector  $\mathbf{v}$  is in Span  $\{\mathbf{u}_1, \dots, \mathbf{u}_k\}$  and in Span  $\{\mathbf{v}_1, \dots, \mathbf{v}_r\}$ then  $\mathbf{v} = \alpha_1 \mathbf{u}_1 + \dots + \alpha_k \mathbf{u}_k$  and  $\mathbf{v} = \beta_1 \mathbf{v}_1 + \dots + \beta_r \mathbf{v}_r$ so

$$\alpha_1 \mathbf{u}_1 + \dots + \alpha_k \mathbf{u}_k = \beta_1 \mathbf{v}_1 + \dots + \beta_r \mathbf{v}_r$$

$$\mathbf{0} = \alpha_1 \, \mathbf{u}_1 + \dots + \alpha_k \, \mathbf{u}_k - \beta_1 \, \mathbf{v}_1 - \dots - \beta_r \, \mathbf{v}_r$$
  
so  $\alpha_1 = \dots = \alpha_k = \beta_1 = \dots = \beta_r = 0$  so  $\mathbf{v} = \mathbf{0}$ . QED

## Linear function invertibility

How to tell if a linear function  $f : \mathcal{V} \longrightarrow \mathcal{W}$  is invertible?

- One-to-one? f is one-to-one if its kernel is trivial. Equivalent: if its kernel has dimension zero.
- Onto? f is onto if its image equals its co-domain

Recall that the image of a function f with domain  $\mathcal{V}$  is  $\{f(\mathbf{v}) : \mathbf{v} \in \mathcal{V}\}$ .

```
Lemma: The image of f is a subspace of \mathcal{W}.
```

How can we tell if the image of f equals W?

```
Dimension Lemma: If \mathcal{U} is a subspace of \mathcal{W} then
```

```
Property D1: dim \mathcal{U} \leq \dim \mathcal{W}, and
```

Property D2: if dim  $\mathcal{U} = \dim \mathcal{W}$  then  $\mathcal{U} = \mathcal{W}$ 

```
Use Property D2 with U = \text{Im } f.
Shows that the function f is onto iff dim Im f = \dim W
```

We conclude:

f is invertible dim Ker f = 0 and dim Im  $f = \dim \mathcal{W}$ 

## Linear function invertibility

f is one-to-one if dim Ker f = 0 and dim Im  $f = \dim \mathcal{W}$ 

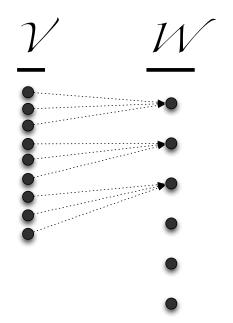
How does this relate to dimension of the domain?

**Conjecture:** For *f* to be invertible, need dim  $\mathcal{V} = \dim \mathcal{W}$ .

Starting with a linear function f we will extract a largest possible subfunction that is invertible.

Make it onto by setting co-domain to be image of f.

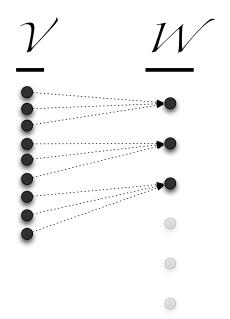
Make it one-to-one by getting rid of extra domain elements sharing same image.



Starting with a linear function f we will extract a largest possible subfunction that is invertible.

Make it onto by setting co-domain to be image of f.

Make it one-to-one by getting rid of extra domain elements sharing same image.

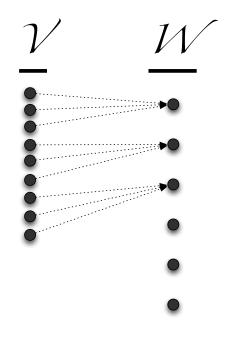


Start with linear function  $f : \mathcal{V} \longrightarrow \mathcal{W}$ Step 1: Choose smaller co-domain  $\mathcal{W}^*$ 

Step 2: Choose smaller domain  $\mathcal{V}^*$ 

Step 3: Define function 
$$f^* : \mathcal{V}^* \longrightarrow \mathcal{W}^*$$
 by  $f^*(\mathbf{x}) = f(\mathbf{x})$ 

In fact, we will end up selecting a *basis* of  $\mathcal{W}^*$  and a basis of  $\mathcal{V}^*$ .

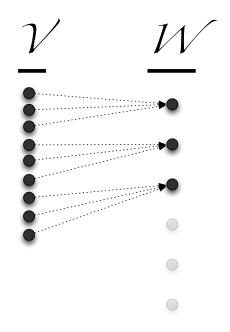


Start with linear function  $f : \mathcal{V} \longrightarrow \mathcal{W}$ Step 1: Choose smaller co-domain  $\mathcal{W}^*$ 

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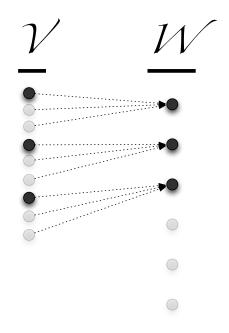


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In fact, we will end up selecting a *basis* of  $\mathcal{W}^*$  and a basis of  $\mathcal{V}^*$ .



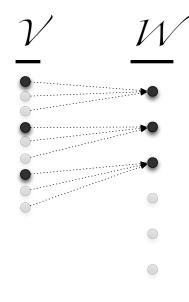
Choose smaller co-domain W\*
 Let W\* be image of f

Let  $\mathbf{w}_1, \ldots, \mathbf{w}_r$  be a basis of  $\mathcal{W}^*$ 

- ► Choose smaller domain V\* Let v<sub>1</sub>,..., v<sub>r</sub> be pre-images of w<sub>1</sub>,..., w<sub>r</sub> That is, f(v<sub>1</sub>) = w<sub>1</sub>,..., f(v<sub>r</sub>) = w<sub>r</sub> Let V\* = Span {v<sub>1</sub>,..., v<sub>r</sub>}
- ▶ Define function f\* : V\* → W\* by f\*(x) = f(x)

We will show:

- ► *f*<sup>\*</sup> is onto
- ▶ *f*<sup>\*</sup> is one-to-one (kernel is trivial)
- Bonus:  $\mathbf{v}_1, \ldots, \mathbf{v}_r$  form a basis for  $\mathcal{V}^*$



Choose smaller co-domain W\*
 Let W\* be image of f

Let  $\mathbf{w}_1, \ldots, \mathbf{w}_r$  be a basis of  $\mathcal{W}^*$ 

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#### **Onto:**

Let **w** be any vector in co-domain  $\mathcal{W}^*$ . There are scalars  $\alpha_1, \ldots, \alpha_r$  such that

 $\mathbf{w} = \alpha_1 \, \mathbf{w}_1 + \dots + \alpha_r \, \mathbf{w}_r$ Because f is linear,  $f(\alpha_1 \, \mathbf{v}_1 + \dots + \alpha_r \, \mathbf{v}_r)$  $= \alpha_1 \, f(\mathbf{v}_1) + \dots + \alpha_r \, f(\mathbf{v}_r)$  $= \alpha_1 \, \mathbf{w}_1 + \dots + \alpha_r \, \mathbf{w}_r$ so w is image of  $\alpha_1 \, \mathbf{v}_1 + \dots + \alpha_r \, \mathbf{v}_r \in \mathcal{V}^*$ QED

Choose smaller co-domain W\*
 Let W\* be image of f

Let  $\mathbf{w}_1, \ldots, \mathbf{w}_r$  be a basis of  $\mathcal{W}^*$ 

► Choose smaller domain V\* Let v<sub>1</sub>,..., v<sub>r</sub> be pre-images of w<sub>1</sub>,..., w<sub>r</sub> That is, f(v<sub>1</sub>) = w<sub>1</sub>,..., f(v<sub>r</sub>) = w<sub>r</sub> Let V\* = Span {v<sub>1</sub>,..., v<sub>r</sub>}

We will show:

► *f*<sup>\*</sup> is onto

- ► *f*<sup>\*</sup> is one-to-one (kernel is trivial)
- Bonus:  $\mathbf{v}_1, \ldots, \mathbf{v}_r$  form a basis for  $\mathcal{V}^*$

#### **One-to-one:**

By One-to-One Lemma, need only show kernel is trivial.

Suppose  $\mathbf{v}^*$  is in  $\mathcal{V}^*$  and  $f(\mathbf{v}^*) = \mathbf{0}$ 

Because  $\mathcal{V}^* = \text{Span } \{\mathbf{v}_1, \dots, \mathbf{v}_r\}$ , there are scalars  $\alpha_1, \dots, \alpha_r$  such that

$$\mathbf{v}^* = \alpha_1 \, \mathbf{v}_1 + \dots + \alpha_r \, \mathbf{v}_r$$

Applying f to both sides,

$$\mathbf{0} = f(\alpha_1 \, \mathbf{v}_1 + \dots + \alpha_r \, \mathbf{v}_r)$$
$$= \alpha_1 \, \mathbf{w}_1 + \dots + \alpha_r \, \mathbf{w}_r$$

QED

Because  $\mathbf{w}_1, \dots, \mathbf{w}_r$  are linearly independent,  $\alpha_1 = \dots = \alpha_r = 0$ so  $\mathbf{v}^* = \mathbf{0}$ 

Choose smaller co-domain W\*
 Let W\* be image of f

Let  $\mathbf{w}_1, \ldots, \mathbf{w}_r$  be a basis of  $\mathcal{W}^*$ 

► Choose smaller domain V\* Let v<sub>1</sub>,..., v<sub>r</sub> be pre-images of w<sub>1</sub>,..., w<sub>r</sub> That is, f(v<sub>1</sub>) = w<sub>1</sub>,..., f(v<sub>r</sub>) = w<sub>r</sub> Let V\* = Span {v<sub>1</sub>,..., v<sub>r</sub>}

We will show:

- ► *f*<sup>\*</sup> is onto
- ▶ *f*<sup>\*</sup> is one-to-one (kernel is trivial)
- Bonus:  $\mathbf{v}_1, \ldots, \mathbf{v}_r$  form a basis for  $\mathcal{V}^*$

**Bonus:**  $\mathbf{v}_1, \ldots, \mathbf{v}_r$  form a basis for  $\mathcal{V}^*$ Need only show linear independence Suppose  $\mathbf{0} = \alpha_1 \mathbf{v}_1 + \cdots + \alpha_r \mathbf{v}_r$ 

Applying f to both sides,

$$\mathbf{0} = f(\alpha_1 \, \mathbf{v}_1 + \dots + \alpha_r \, \mathbf{v}_r)$$
$$= \alpha_1 \, \mathbf{w}_1 + \dots + \alpha_r \, \mathbf{w}_r$$

Because  $\mathbf{w}_1, \ldots, \mathbf{w}_r$  are linearly independent,  $\alpha_1 = \cdots = \alpha_r = 0$ . QED

Choose smaller co-domain W\*
 Let W\* be image of f

Let  $\mathbf{w}_1, \ldots, \mathbf{w}_r$  be a basis of  $\mathcal{W}^*$ 

- Choose smaller domain V<sup>\*</sup> Let v<sub>1</sub>,..., v<sub>r</sub> be pre-images of w<sub>1</sub>,..., w<sub>r</sub> That is, f(v<sub>1</sub>) = w<sub>1</sub>,..., f(v<sub>r</sub>) = w<sub>r</sub> Let V<sup>\*</sup> = Span {v<sub>1</sub>,..., v<sub>r</sub>}
- ▶ Define function f\* : V\* → W\* by f\*(x) = f(x)

We will show:

- ► *f*<sup>\*</sup> is onto
- ▶ *f*<sup>\*</sup> is one-to-one (kernel is trivial)
- Bonus:  $\mathbf{v}_1, \ldots, \mathbf{v}_r$  form a basis for  $\mathcal{V}^*$

#### Example:

Let  $A = \begin{bmatrix} 1 & 2 & 1 \\ 2 & 1 & 1 \\ 1 & 2 & 1 \end{bmatrix}$ , and define  $\mathbf{f} : \mathbb{R}^3 \longrightarrow \mathbb{R}^3$  by  $f(\mathbf{x}) = A\mathbf{x}$ .

Define  $W^* = \text{Im } f = \text{Col } A =$ Span {[1,2,1],[2,1,2],[1,1,1]}.

One basis for  $\mathcal{W}^*$  is  $\label{eq:w1} \boldsymbol{w}_1 = [0,1,0]\text{, } \boldsymbol{w}_2 = [1,0,1]$ 

Pre-images for  $\mathbf{w}_1$  and  $\mathbf{w}_2$ :  $\mathbf{v}_1 = [\frac{1}{2}, -\frac{1}{2}, \frac{1}{2}]$  and  $\mathbf{v}_2 = [-\frac{1}{2}, \frac{1}{2}, \frac{1}{2}]$ , for then  $A\mathbf{v}_1 = \mathbf{w}_1$  and  $A\mathbf{v}_2 = \mathbf{w}_2$ .

Let 
$$\mathcal{V}^*=\mathsf{Span}\ \{\textbf{v}_1,\textbf{v}_2\}$$

Then  $f^* : \mathcal{V}^* \longrightarrow \text{Im } f$  is onto and one-to-one.

► Choose smaller co-domain W\* Let W\* be image of f

Let  $\mathbf{w}_1, \ldots, \mathbf{w}_r$  be a basis of  $\mathcal{W}^*$ 

- Choose smaller domain V<sup>\*</sup> Let v<sub>1</sub>,..., v<sub>r</sub> be pre-images of w<sub>1</sub>,..., w<sub>r</sub> That is, f(v<sub>1</sub>) = w<sub>1</sub>,..., f(v<sub>r</sub>) = w<sub>r</sub> Let V<sup>\*</sup> = Span {v<sub>1</sub>,..., v<sub>r</sub>}
- ▶ Define function f\* : V\* → W\* by f\*(x) = f(x)

We will show:

- ► *f*<sup>\*</sup> is onto
- ▶ *f*<sup>\*</sup> is one-to-one (kernel is trivial)
- Bonus:  $\mathbf{v}_1, \ldots, \mathbf{v}_r$  form a basis for  $\mathcal{V}^*$

To show about original function f: original domain  $\mathcal{V} = \text{Ker } f \oplus \mathcal{V}^*$ Must prove two things:

- 1. Ker f and  $\mathcal{V}^*$  share only zero vector
- 2. every vector in  $\mathcal{V}$  is the sum of a vector in Ker f and a vector in  $\mathcal{V}^*$

We already showed kernel of  $f^*$  is trivial. This shows only vector of Ker f in  $\mathcal{V}^*$  is zero vector. —thing 1 is proved.

Let  $\mathbf{v}$  be any vector in  $\mathcal{V}$ , and let  $\mathbf{w} = f(\mathbf{v})$ . Since  $f^*$  is onto, its domain  $\mathcal{V}^*$  contains a vector  $\mathbf{v}^*$  such that  $f(\mathbf{v}^*) = \mathbf{w}$ Therefore  $f(\mathbf{v}) = f(\mathbf{v}^*)$  so  $f(\mathbf{v}) - f(\mathbf{v}^*) = \mathbf{0}$  so  $f(\mathbf{v} - \mathbf{v}^*) = \mathbf{0}$ Thus  $\mathbf{u} = \mathbf{v} - \mathbf{v}^*$  is in Ker fand  $\mathbf{v} = \mathbf{u} + \mathbf{v}^*$ —thing 2 is proved.

Choose smaller co-domain W\*
 Let W\* be image of f

Let  $\mathbf{w}_1, \ldots, \mathbf{w}_r$  be a basis of  $\mathcal{W}^*$ 

 Choose smaller domain V<sup>\*</sup> Let v<sub>1</sub>,..., v<sub>r</sub> be pre-images of w<sub>1</sub>,..., w<sub>r</sub> That is, f(v<sub>1</sub>) = w<sub>1</sub>,..., f(v<sub>r</sub>) = w<sub>r</sub> Let V<sup>\*</sup> = Span {v<sub>1</sub>,..., v<sub>r</sub>}

We will show:

- ▶ f\* is onto
- ► *f*<sup>\*</sup> is one-to-one (kernel is trivial)
- Bonus:  $\mathbf{v}_1, \ldots, \mathbf{v}_r$  form a basis for  $\mathcal{V}^*$

original domain  $\mathcal{V} = \text{Ker } f \oplus \mathcal{V}^*$ **Example:** Let  $A = \begin{bmatrix} 1 & 2 & 1 \\ 2 & 1 & 1 \\ 1 & 2 & 1 \end{bmatrix}$ , and define  $\mathbf{f} : \mathbb{R}^3 \longrightarrow \mathbb{R}^3$  by  $f(\mathbf{x}) = A\mathbf{x}$ .  $\mathbf{v}_1 = [\frac{1}{2}, -\frac{1}{2}, \frac{1}{2}]$  and  $\mathbf{v}_2 = [-\frac{1}{2}, \frac{1}{2}, \frac{1}{2}]$  $\mathcal{V}^* = \text{Span} \{\mathbf{v}_1, \mathbf{v}_2\}$ Ker  $f = \text{Span} \{ [1, 1, -3] \}$ Therefore  $\mathcal{V} = (\text{Span} \{[1, 1, -3]\}) \oplus (\text{Span} \{\mathbf{v}_1, \mathbf{v}_2\})$ 

Choose smaller co-domain W\*
 Let W\* be image of f

Let  $\mathbf{w}_1, \ldots, \mathbf{w}_r$  be a basis of  $\mathcal{W}^*$ 

► Choose smaller domain V\* Let v<sub>1</sub>,..., v<sub>r</sub> be pre-images of w<sub>1</sub>,..., w<sub>r</sub> That is, f(v<sub>1</sub>) = w<sub>1</sub>,..., f(v<sub>r</sub>) = w<sub>r</sub> Let V\* = Span {v<sub>1</sub>,..., v<sub>r</sub>}

We will show:

- ► *f*<sup>\*</sup> is onto
- ▶ *f*<sup>\*</sup> is one-to-one (kernel is trivial)
- Bonus:  $\mathbf{v}_1, \ldots, \mathbf{v}_r$  form a basis for  $\mathcal{V}^*$

original domain  $\mathcal{V} = \text{Ker } f \oplus \mathcal{V}^*$ By Direct-Sum Dimension Corollary,  $\dim \mathcal{V} = \dim \text{Ker } f + \dim \mathcal{V}^*$ 

Since  $\mathbf{v}_1, \ldots, \mathbf{v}_r$  form a basis for  $\mathcal{V}^*$ , dim  $\mathcal{V}^* = r = \dim \operatorname{Im} f$ 

We have proved...

Kernel-Image Theorem: For any linear function  $f : \mathcal{V} \to W$ , dim Ker  $f + \dim \operatorname{Im} f = \dim \mathcal{V}$ 

## Linear function invertibility, revisited

**Kernel-Image Theorem:** For any linear function  $f : \mathcal{V} \to W$ ,

 $\dim \operatorname{Ker} \, f + \dim \operatorname{Im} \, f = \dim \mathcal{V}$ 

**Linear-Function Invertibility Theorem:** Let  $f : \mathcal{V} \longrightarrow \mathcal{W}$  be a linear function. Then f is invertible iff dim Ker f = 0 and dim  $\mathcal{V} = \dim \mathcal{W}$ .

**Proof:** We saw before that *f* 

- is one-to-one iff dim Ker f = 0
- is onto if dim Im  $f = \dim W$

Therefore f is invertible if dim Ker f = 0 and dim Im  $f = \dim W$ .

Kernel-Image Theorem states dim Ker  $f + \dim \operatorname{Im} f = \dim \mathcal{V}$ 

Therefore

dim Ker f = 0 and dim Im  $f = \dim \mathcal{W}$ iff dim Ker f = 0 and dim  $\mathcal{V} = \dim \mathcal{W}$ 

## Rank-Nullity Theorem

Kernel-Image Theorem:

For any linear function  $f: \mathcal{V} \to W$ ,

 $\dim \operatorname{Ker} f + \dim \operatorname{Im} f = \dim \mathcal{V}$ 

Apply Kernel-Image Theorem to the function  $f(\mathbf{x}) = A\mathbf{x}$ :

- Ker f = Null A
- dim Im  $f = \dim \text{Col } A = \text{rank } A$

**Definition:** The *nullity* of matrix A is dim Null A

**Rank-Nullity Theorem:** For any *n*-column matrix *A*,

nullity  $A + \operatorname{rank} A = n$ 

## Checksum problem revisited

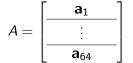
Checksum function maps *n*-vectors over *GF*(2) to 64-vectors over *GF*(2):  $\mathbf{x} \mapsto [\mathbf{a}_1 \cdot \mathbf{x}, \dots, \mathbf{a}_{64} \cdot \mathbf{x}]$ 

Original "file"  $\mathbf{p}$ , transmission error  $\mathbf{e}$  so corrupted file is  $\mathbf{p} + \mathbf{e}$ .

If error is chosen according to uniform distribution, Probability  $(\mathbf{p} + \mathbf{e} \text{ has same checksum as } \mathbf{p})$ 

$$=\frac{2^{\dim \mathcal{V}}}{2^n}$$

where  $\mathcal{V}$  is the null space of the matrix



Fact: Can easily choose  $\mathbf{a}_1, \ldots, \mathbf{a}_{64}$  so that rank A = 64

(Randomly chosen vectors will probably work.)

**Rank-Nullity Theorem**  $\Rightarrow$ rank A + nullity A = n64 + dim  $\mathcal{V}$  = ndim  $\mathcal{V}$  = n-64Therefore Probability =  $\frac{2^{n-64}}{2^n} = \frac{1}{2^{64}}$ 

**very** tiny chance that the change is undetected

Rank-Nullity Theorem: For any n-column matrix A,

nullity  $A + \operatorname{rank} A = n$ 

**Corollary:** Let A be an  $R \times C$  matrix. Then A is invertible if and only if |R| = |C| and the columns of A are linearly independent.

**Proof:** Let  $\mathbb{F}$  be the field. Define  $f : \mathbb{F}^C \longrightarrow \mathbb{F}^R$  by  $f(\mathbf{x}) = A\mathbf{x}$ . Then A is an invertible matrix if and only if f is an invertible function.

The function f is invertible iff dim Ker f = 0 and dim  $\mathbb{F}^C = \dim \mathbb{F}^R$ iff nullity A = 0 and |C| = |R|.

nullity A = 0 iff dim Null A = 0

- iff Null  $A = \{0\}$
- iff the only vector  $\mathbf{x}$  such that  $A\mathbf{x} = \mathbf{0}$  is  $\mathbf{x} = \mathbf{0}$
- iff the columns of A are linearly independent. QED

## Matrix invertibility examples

```
\begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{bmatrix} is not square so cannot be invertible.
\begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix} is square and its columns are linearly independent so it is invertible.
\begin{bmatrix} 1 & 1 & 2 \\ 2 & 1 & 3 \\ 3 & 1 & 4 \end{bmatrix} is square but columns not linearly independent so it is not invertible.
```

## Transpose of invertible matrix is invertible

Theorem: The transpose of an invertible matrix is invertible.

$$A = \left[ \begin{array}{c|c} \mathbf{v}_1 \\ \cdots \\ \mathbf{v}_n \end{array} \right] = \left[ \begin{array}{c|c} \mathbf{a}_1 \\ \vdots \\ \hline \mathbf{a}_n \end{array} \right] \qquad \qquad A^T = \left[ \begin{array}{c|c} \mathbf{a}_1 \\ \cdots \\ \mathbf{a}_n \end{array} \right]$$

**Proof:** Suppose A is invertible. Then A is square and its columns are linearly independent. Let n be the number of columns. Then rank A = n.

Because A is square, it has n rows. By Rank Theorem, rows are linearly independent. Columns of transpose  $A^T$  are rows of A, so columns of  $A^T$  are linearly independent. Since  $A^T$  is square and columns are linearly independent,  $A^T$  is invertible. QED

## More matrix invertibility

Earlier we proved: If A has an inverse  $A^{-1}$  then  $AA^{-1}$  is identity matrix **Converse:** If BA is identity matrix then A and B are inverses? **Not always true.** 

**Theorem:** Suppose A and B are square matrices such that BA is an identity matrix 1. Then A and B are inverses of each other.

**Proof:** To show that A is invertible, need to show its columns are linearly independent.

Let **u** be any vector such that  $A\mathbf{u} = \mathbf{0}$ . Then  $B(A\mathbf{u}) = B\mathbf{0} = \mathbf{0}$ . On the other hand,  $(BA)\mathbf{u} = \mathbb{1}\mathbf{u} = \mathbf{u}$ , so  $\mathbf{u} = \mathbf{0}$ .

This shows A has an inverse  $A^{-1}$ . Now must show  $B = A^{-1}$ . We know  $AA^{-1}$  is an identity matrix.

$$BA = 1$$

$$(BA)A^{-1} = 1A^{-1}$$
by multiplying on the right by  $B^{-1}$ 

$$(BA)A^{-1} = A^{-1}$$

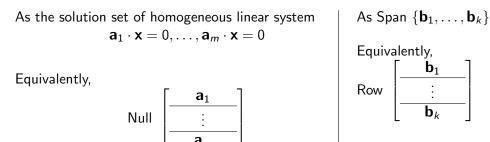
$$B(AA^{-1}) = A^{-1}$$
by associativity of matrix-matrix mult
$$B 1 = A^{-1}$$

$$B = A^{-1}$$

$$QED$$

## Representations of vector spaces

Two important ways to represent a vector space:



How to transform between these two representations?

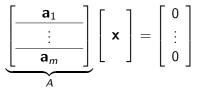
**From left to right:** Given homogeneous linear system  $\mathbf{a}_1 \cdot \mathbf{x} = 0, \dots, \mathbf{a}_m \cdot \mathbf{x} = 0$ , find generators  $\mathbf{b}_1, \dots, \mathbf{b}_k$  for solution set

#### From right to left:

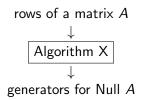
Given generators  $\mathbf{b}_1, \ldots, \mathbf{b}_k$ , find homogeneous linear system  $\mathbf{a}_1 \cdot \mathbf{x} = 0, \ldots, \mathbf{a}_m \cdot \mathbf{x} = 0$  whose solution set equals Span  $\{\mathbf{b}_1, \ldots, \mathbf{b}_k\}$ 

**From left to right:** Given system  $\mathbf{a}_1 \cdot \mathbf{x} = 0, \dots, \mathbf{a}_m \cdot \mathbf{x} = 0$ , find generators  $\mathbf{b}_1, \dots, \mathbf{b}_k$  for solution set

Solution set is the set of vectors  $\mathbf{u}$  such that  $\mathbf{a}_1 \cdot \mathbf{u} = 0, \dots, \mathbf{a}_m \cdot \mathbf{u} = 0$ 



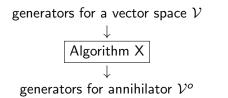
**Equivalent:** Given rows of a matrix *A*, find generators for Null *A* 



If **u** is such a vector then  $\mathbf{u} \cdot (\alpha_1 \mathbf{a}_1 + \dots + \alpha_m \mathbf{a}_m) = 0$ for any coefficients  $\alpha_1, \dots, \alpha_m$ .

**Definition:** The set of vectors  $\mathbf{u}$  such that  $\mathbf{u} \cdot \mathbf{v} = 0$  for every vector  $\mathbf{v}$  in  $\mathcal{V}$  is called the *annihilator* of  $\mathcal{V}$ . Written as  $\mathcal{V}^o$ .

**Example:** The annihilator of Span  $\{\mathbf{a}_1, \ldots, \mathbf{a}_m\}$  is the solution set for  $\mathbf{a}_1 \cdot \mathbf{x} = 0, \ldots, \mathbf{a}_m \cdot \mathbf{x} = 0$ 



**Definition:** For a subspace  $\mathcal{V}$  of  $\mathbb{F}^n$ , the *annihilator* of  $\mathcal{V}$ , written  $\mathcal{V}^o$ , is

$$\mathcal{V}^{o} = \{ \mathbf{u} \in \mathbb{F}^{n} : \mathbf{u} \cdot \mathbf{v} = 0 \text{ for every vector } \mathbf{v} \in \mathcal{V} \}$$

**Example over**  $\mathbb{R}$ : Let  $\mathcal{V} = \text{Span} \{[1,0,1], [0,1,0]\}$ . Then  $\mathcal{V}^o = \text{Span} \{[1,0,-1]\}$ :

- ▶ Note that  $[1, 0, -1] \cdot [1, 0, 1] = 0$  and  $[1, 0, -1] \cdot [0, 1, 0] = 0$ . Therefore  $[1, 0, -1] \cdot \mathbf{v} = 0$  for every vector  $\mathbf{v}$  in Span  $\{[1, 0, 1], [0, 1, 0]\}$ .
- For any scalar  $\beta$ ,

$$\beta \left[ 1,0,-1
ight] \cdot \mathbf{v}=eta \left( \left[ 1,0,-1
ight] \cdot \mathbf{v}
ight) =0$$

for every vector **v** in Span  $\{[1, 0, 1], [0, 1, 0]\}$ .

Which vectors u satisfy u · v = 0 for every vector v in Span {[1,0,1],[0,1,0]}? Only scalar multiples of [1,0,−1].

**Example over** *GF*(2): Let  $\mathcal{V} = \text{Span} \{[1,0,1], [0,1,0]\}$ . Then  $\mathcal{V}^o = \text{Span} \{[1,0,1]\}$ :

- Note that [1,0,1] · [1,0,1] = 0 (remember GF(2) addition) and [1,0,1] · [0,1,0] = 0.
- Therefore  $[1,0,1] \cdot \mathbf{v} = 0$  for every vector  $\mathbf{v}$  in Span  $\{[1,0,1], [0,1,0]\}$ .
- Of course  $[0,0,0] \cdot \mathbf{v} = 0$  for every vector  $\mathbf{v}$  in Span  $\{[1,0,1], [0,1,0]\}$ .
- [1,0,1] and [0,0,0] are the only such vectors.

**Example over**  $\mathbb{R}$ : Let  $\mathcal{V} = \text{Span} \{[1,0,1], [0,1,0]\}$ . Then  $\mathcal{V}^o = \text{Span} \{[1,0,-1]\}$ dim  $\mathcal{V} + \text{dim} \mathcal{V}^o = 3$ 

**Example over** *GF*(2): Let  $\mathcal{V} = \text{Span} \{[1, 0, 1], [0, 1, 0]\}$ . Then  $\mathcal{V}^o = \text{Span} \{[1, 0, 1]\}$ . dim  $\mathcal{V} + \dim \mathcal{V}^o = 3$ 

**Example over**  $\mathbb{R}$ : Let  $\mathcal{V} = \text{Span} \{[1, 0, 1, 0], [0, 1, 0, 1]\}.$ Then  $\mathcal{V}^o = \text{Span} \{[1, 0, -1, 0], [0, 1, 0, -1]\}.$ dim  $\mathcal{V} + \text{dim} \mathcal{V}^o = 4$ 

**Annihilator Dimension Theorem:** dim  $\mathcal{V}$  + dim  $\mathcal{V}^o$  = n

**Proof:** Let  $\mathbf{a}_1, \ldots, \mathbf{a}_m$  be generators for  $\mathcal{V}$ .

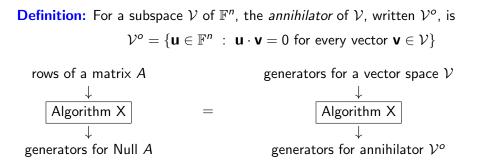
Let 
$$A = \begin{bmatrix} \mathbf{a}_1 \\ \vdots \\ \mathbf{a}_m \end{bmatrix}$$

Then  $\mathcal{V}^{o} = \text{Null } A$ .

Rank-Nullity Theorem states that

 $\begin{array}{rrr} {\rm rank} \ A & + & {\rm nullity} \ A & = & n \\ {\rm dim} \ \mathcal{V} & + & {\rm dim} \ \mathcal{V}^o & = & n \end{array}$ 

QED

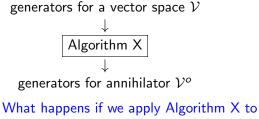


**From left to right:** Given system  $\mathbf{a}_1 \cdot \mathbf{x} = 0, \dots, \mathbf{a}_m \cdot \mathbf{x} = 0$ , find generators  $\mathbf{b}_1, \dots, \mathbf{b}_k$  for solution set

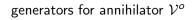
Algorithm X solves left-to-right problem....

what about right-to-left problem?

From left to right: Given system  $\mathbf{a}_1 \cdot \mathbf{x} = 0, \dots, \mathbf{a}_m \cdot \mathbf{x} = 0$ , find generators  $\mathbf{b}_1, \dots, \mathbf{b}_k$  for solution set



generators for annihilator  $\mathcal{V}^o$ ?



#### ↓ Algorithm X

generators for annihilator of annihilator  $(\mathcal{V}^o)^o$ 

From right to left: Given generators  $\mathbf{b}_1, \ldots, \mathbf{b}_k$ , find system  $\mathbf{a}_1 \cdot \mathbf{x} = 0, \ldots, \mathbf{a}_m \cdot \mathbf{x} = 0$  whose solution set equals Span  $\{\mathbf{b}_1, \ldots, \mathbf{b}_k\}$ 

generators for annihilator  $\mathcal{V}^o$   $\downarrow$ Algorithm Y  $\downarrow$ generators for original space  $\mathcal{V}$ 

**Theorem:**  $(\mathcal{V}^o)^o = \mathcal{V}$  (The annihilator of the annihilator is the original space.)

Theorem shows:

Algorithm X = Algorithm Y

We still must prove the Theorem ...

Annihilator

**Theorem:**  $(\mathcal{V}^o)^o = \mathcal{V}$  (The annihilator of the annihilator is the original space.)

#### **Proof:**

Let  $\mathbf{a}_1, \ldots, \mathbf{a}_m$  be a basis for  $\mathcal{V}$ . Let  $\mathbf{b}_1, \ldots, \mathbf{b}_k$  be a basis for  $\mathcal{V}^o$ . Since  $\mathbf{b}_1 \cdot \mathbf{v} = 0$  for every vector  $\mathbf{v}$  in  $\mathcal{V}$ ,

$$\mathbf{b}_1 \cdot \mathbf{a}_1 = 0, \mathbf{b}_1 \cdot \mathbf{a}_2 = 0, \dots, \mathbf{b}_1 \cdot \mathbf{a}_m = 0$$

Similarly  $\mathbf{b}_i \cdot \mathbf{a}_1 = 0, \mathbf{b}_i \cdot \mathbf{a}_2 = 0, \dots, \mathbf{b}_i \cdot \mathbf{a}_m = 0$  for  $i = 1, 2, \dots, k$ . Reorganizing,

$$\mathbf{a}_1 \cdot \mathbf{b}_1 = 0, \mathbf{a}_1 \cdot \mathbf{b}_2 = 0, \dots, \mathbf{a}_1 \cdot \mathbf{b}_k = 0$$
  
which implies that  $\mathbf{a}_1 \cdot \mathbf{u} = 0$  for every vector  $\mathbf{u}$  in  $\underbrace{\text{Span } \{\mathbf{b}_1, \dots, \mathbf{b}_k\}}_{\mathcal{V}^o}$ 

This shows  $\mathbf{a}_1$  is in  $(\mathcal{V}^o)^o$ . Similarly  $\mathbf{a}_2$  is in  $(\mathcal{V}^o)^o$ ,  $\mathbf{a}_3$  is in  $(\mathcal{V}^o)^o$ , ...,  $\mathbf{a}_m$  is in  $(\mathcal{V}^o)^o$ . Therefore every vector in Span  $\{\mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_m\}$  is in  $(\mathcal{V}^o)^o$ . Thus  $\underbrace{\text{Span } \{\mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_m\}}_{\mathcal{V}}$  is a subspace of  $(\mathcal{V}^o)^o$ .

To show that these are equal, we must show that dim  $\mathcal{V} = \dim(\mathcal{V}^o)^o$ .

 $\mathbf{D}_{1}$ ,  $\mathbf{A}_{1}$  is the interval of  $\mathbf{D}_{1}$  is a second of  $\mathbf{D}_{1}$  is the interval of  $\mathbf{D}_{1}$  is the interval of  $\mathbf{D}_{2}$  is the interval

Annihilator

**Theorem:**  $(\mathcal{V}^o)^o = \mathcal{V}$  (The annihilator of the annihilator is the original space.) **Proof:** Reorganizing,

$$\mathbf{a}_1 \cdot \mathbf{b}_1 = 0, \mathbf{a}_1 \cdot \mathbf{b}_2 = 0, \dots, \mathbf{a}_1 \cdot \mathbf{b}_k = 0$$
  
which implies that  $\mathbf{a}_1 \cdot \mathbf{u} = 0$  for every vector  $\mathbf{u}$  in  $\underbrace{\text{Span } \{\mathbf{b}_1, \dots, \mathbf{b}_k\}}_{\mathcal{V}^o}$ 

This shows  $\mathbf{a}_1$  is in  $(\mathcal{V}^o)^o$ . Similarly  $\mathbf{a}_2$  is in  $(\mathcal{V}^o)^o$ ,  $\mathbf{a}_3$  is in  $(\mathcal{V}^o)^o$ , ...,  $\mathbf{a}_m$  is in  $(\mathcal{V}^o)^o$ . Therefore every vector in Span  $\{\mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_m\}$  is in  $(\mathcal{V}^o)^o$ . Thus Span  $\{\mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_m\}$  is a subspace of  $(\mathcal{V}^o)^o$ .

 $\mathcal{V}$ To show that these are equal, we must show that dim  $\mathcal{V} = \dim(\mathcal{V}^o)^o$ . By Annihilator Dimension Theorem, dim  $\mathcal{V} + \dim \mathcal{V}^o = n$ .

By Annihilator Dimension Theorem applied to  $\mathcal{V}^o$ , dim  $\mathcal{V}^o + \dim(\mathcal{V}^o)^o = n$ . Together these equations show dim  $\mathcal{V} = \dim(\mathcal{V}^o)^o$ .

QED